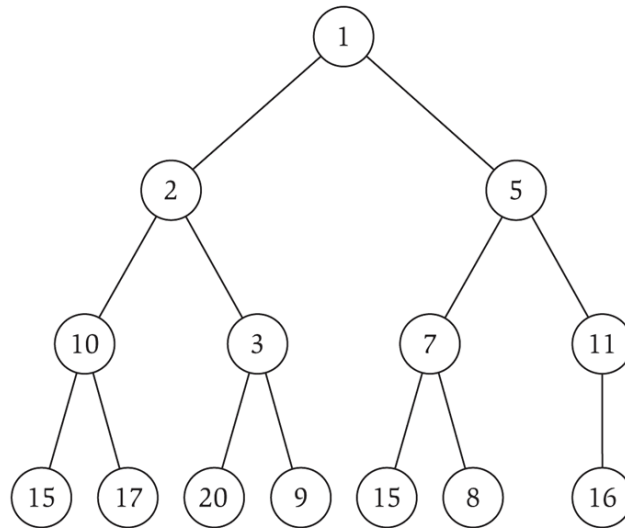


3.1 Graphs: Basic Definitions and Applications

From trees to graphs

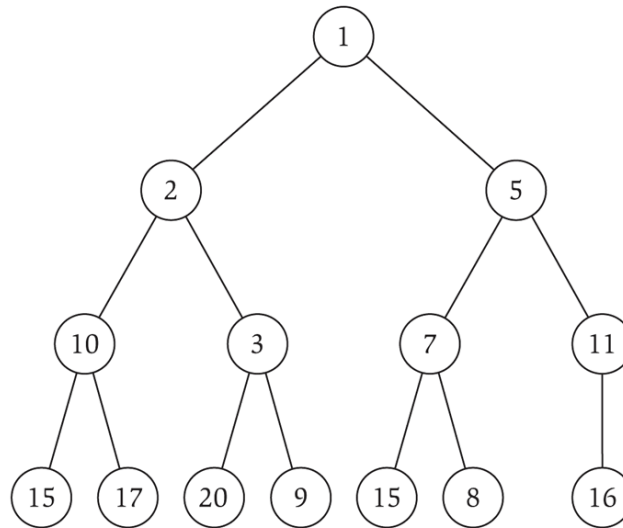
Q. How many edges does a tree with n nodes (internal+leaves) have?



From trees to graphs

Q. How many edges does a tree with n nodes (internal+leaves) have?

A. $n-1$

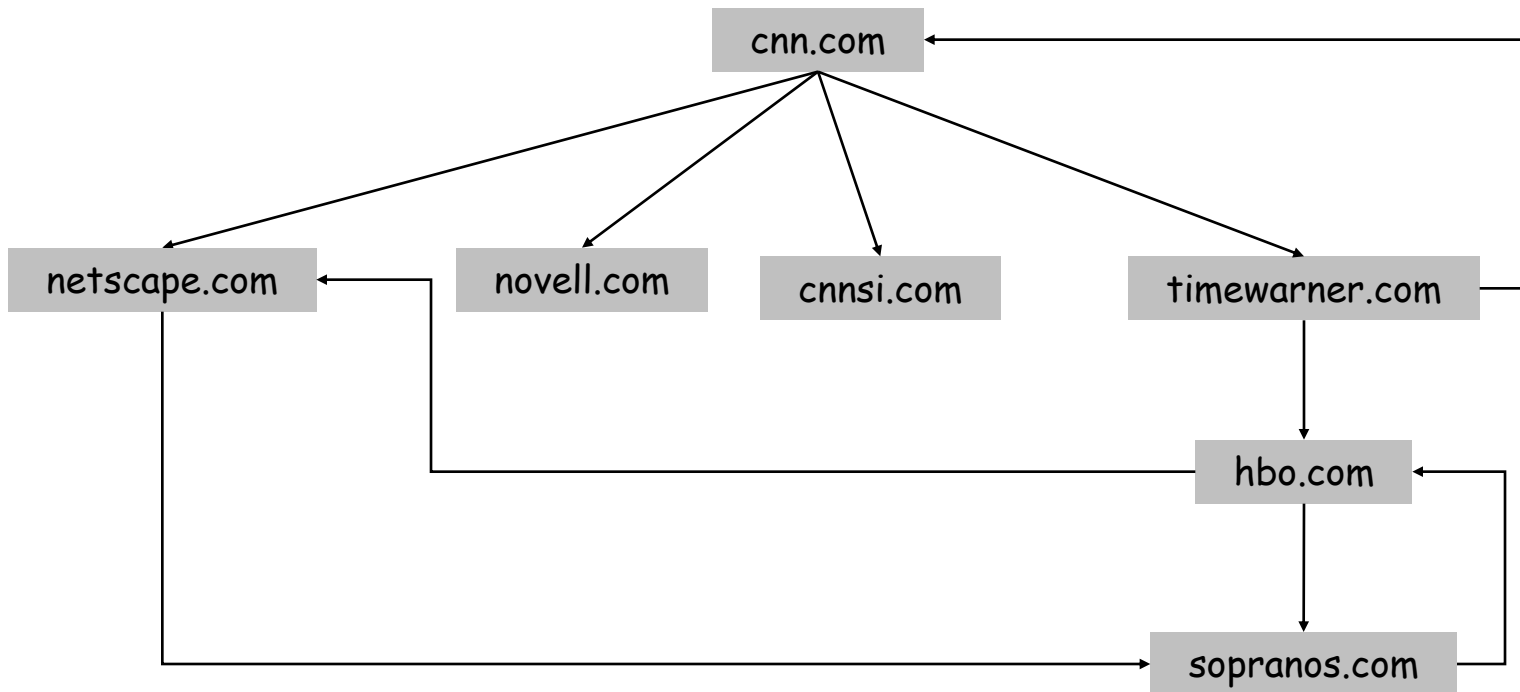


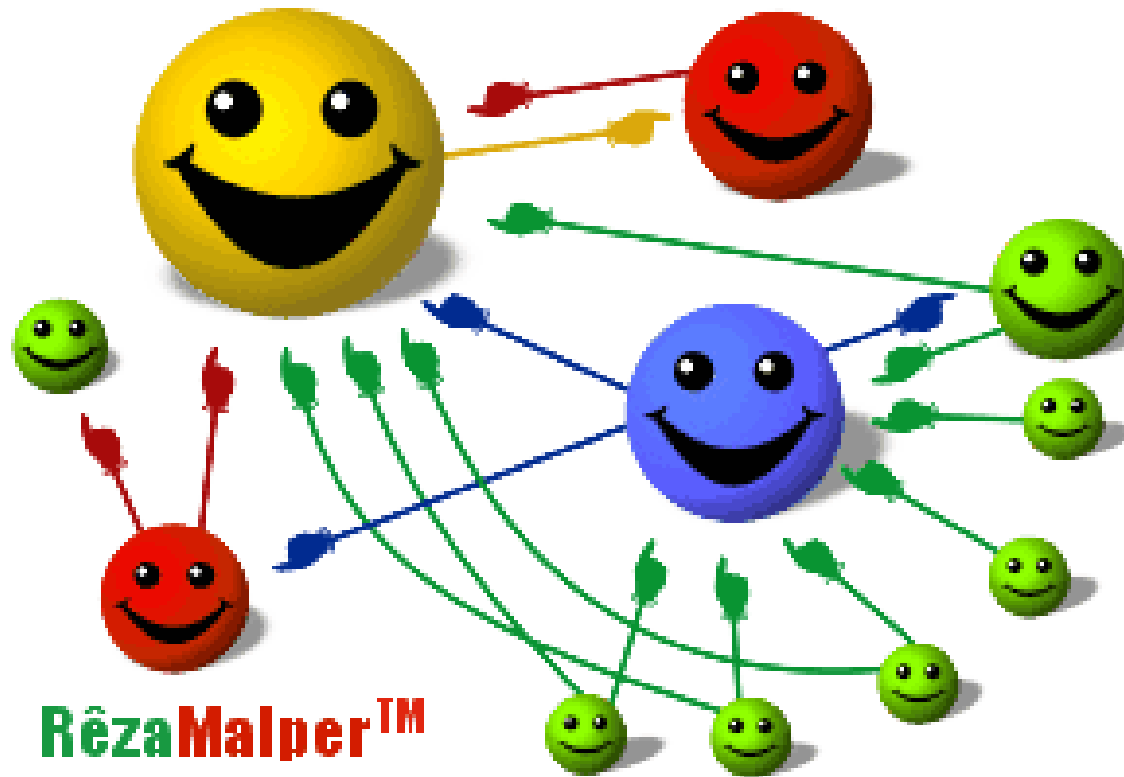


World Wide Web

Web graph.

- Node: web page.
- Edge: hyperlink from one page to another.





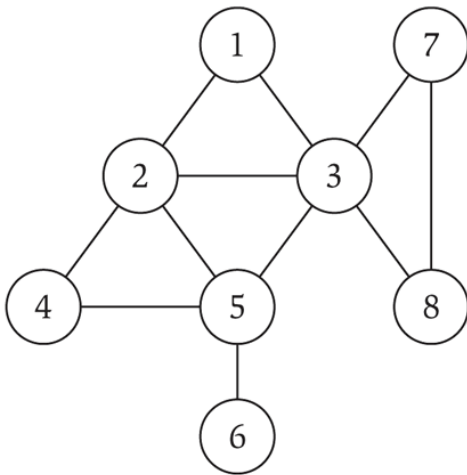
Why Graphs?

<i>Graph</i>	<i>Nodes</i>	<i>Edges</i>
transportation	street intersections	highways
communication	computers	fiber optic cables
World Wide Web	web pages	hyperlinks
social	people	relationships
food web	species	predator-prey
software systems	functions	function calls
scheduling	tasks	precedence constraints
circuits	gates	wires

Undirected Graphs

Undirected graph. $G = (V, E)$

- V = set of nodes (=vertices). ("knopen")
- E = set of edges between pairs of nodes. ("kanten")
- Captures pairwise relationship between objects.
- Graph size parameters: $n = |V|$, $m = |E|$.



$V = \{ 1, 2, 3, 4, 5, 6, 7, 8 \}$

$E = \{ 1-2, 1-3, 2-3, 2-4, 2-5, 3-5, 3-7, 3-8, 7-8, 4-5, 5-6 \}$

$n = 8$

$m = 11$

Q. How to implement a graph? Which datastructure to use? (1 min.)

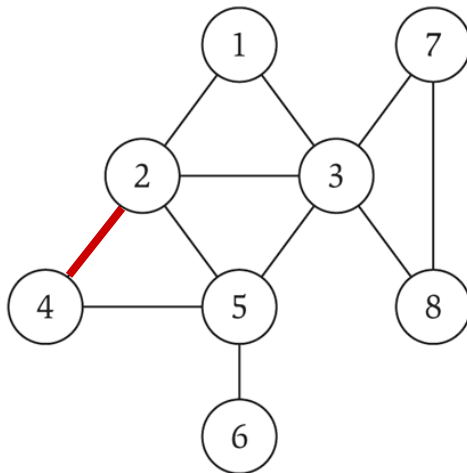
- How much space do you need?
- How much time to check whether node 2 and 4 neighbors?
- How much time to list all edges?

Graph Representation: Adjacency Matrix

Adjacency matrix. n -by- n matrix with $A_{uv} = 1$ if (u, v) is an edge.

- Two representations of each edge.
- Space proportional to n^2 .
- Checking if (u, v) is an edge takes $\Theta(1)$ time.
- Identifying all edges takes $\Theta(n^2)$ time.

Q. How is **this edge** between 2 and 4 represented in the matrix?



	1	2	3	4	5	6	7	8
1	0	1	1	0	0	0	0	0
2	1	0	1	1	1	0	0	0
3	1	1	0	0	1	0	1	1
4	0	1	0	0	1	0	0	0
5	0	1	1	1	0	1	0	0
6	0	0	0	0	1	0	0	0
7	0	0	1	0	0	0	0	1
8	0	0	1	0	0	0	1	0

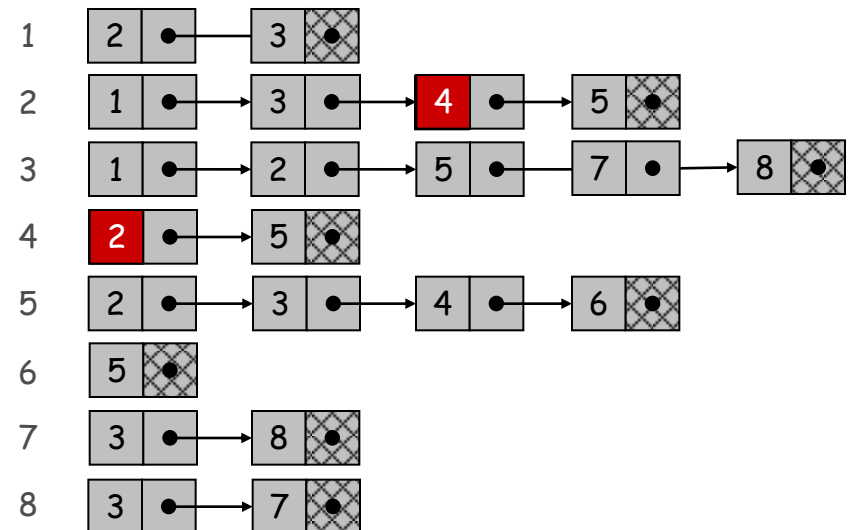
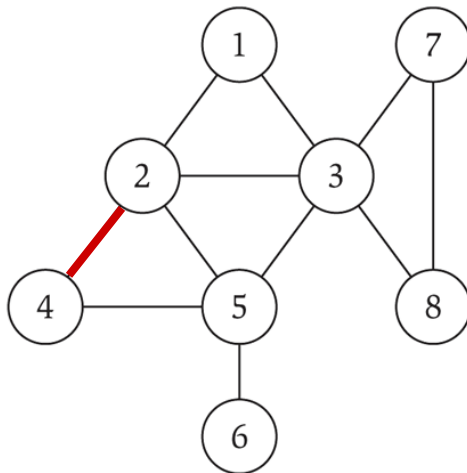
Graph Representation: Adjacency List

Adjacency list. Node indexed array of lists.

- Two representations of each edge.
- Space proportional to $m + n$.
- Checking if (u, v) is an edge takes $O(\text{deg}(u))$ time.
- Identifying all edges takes $\Theta(m + n)$ time.

degree = number of neighbors of u

Q. How is **this edge** between 2 and 4 represented in the adjacency list?



Paths and Connectivity

Def. A **path** in an undirected graph $G = (V, E)$ is a sequence P of nodes $v_1, v_2, \dots, v_{k-1}, v_k$ with the property that each consecutive pair v_i, v_{i+1} is joined by an edge in E .

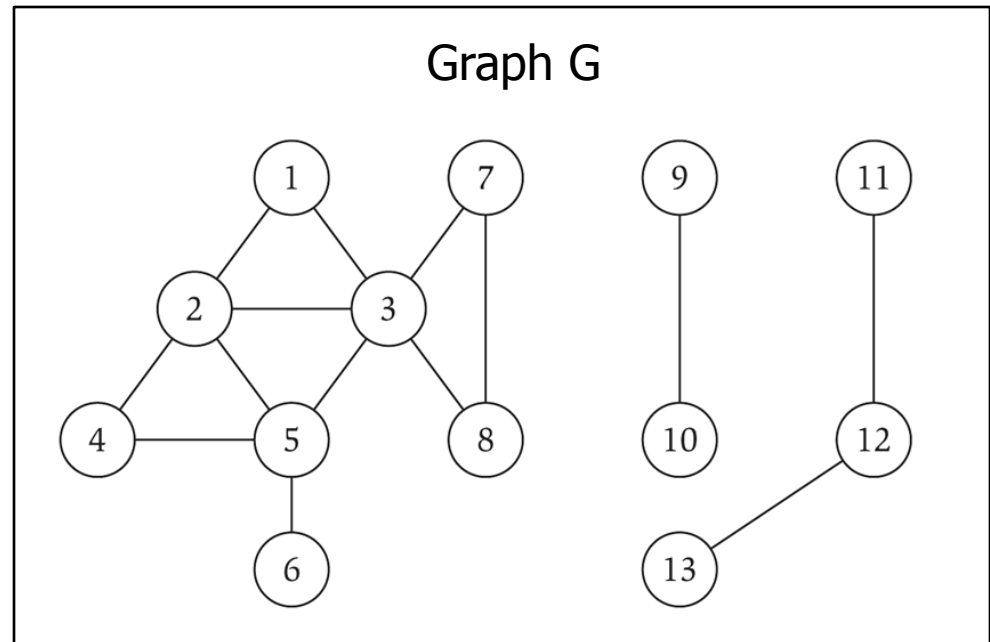
Def. A path is **simple** if all nodes are distinct.

Def. An undirected graph is **connected** if for every pair of nodes u and v , there is a path between u and v .

Q. Is 1-3-8-7-3-5 a path?

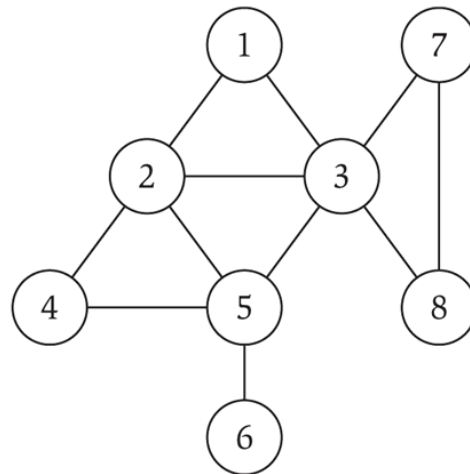
Q. Is it a simple path?

Q. Is this graph connected?



Cycles

Def. A **cycle** is a path $v_1, v_2, \dots, v_{k-1}, v_k$ in which $v_1 = v_k$, $k > 2$, and the first $k-1$ nodes are all distinct.



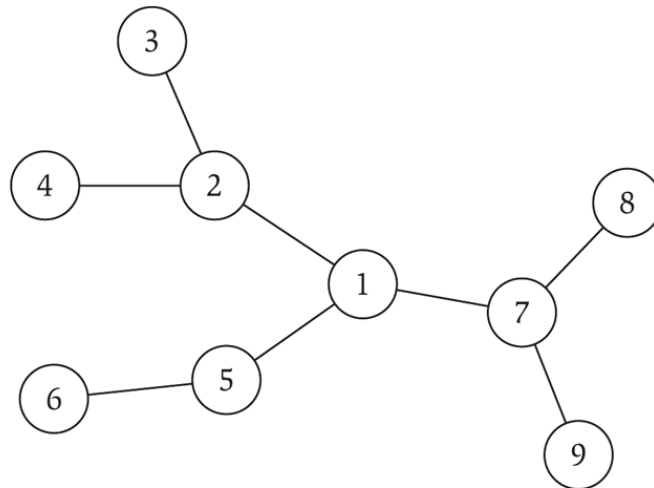
cycle $C = 1-2-4-5-3-1$

Trees

Def. An undirected graph is a **tree** if it is connected and does not contain a cycle.

Theorem. Let G be an undirected graph on n nodes. Any two of the following statements imply the third.

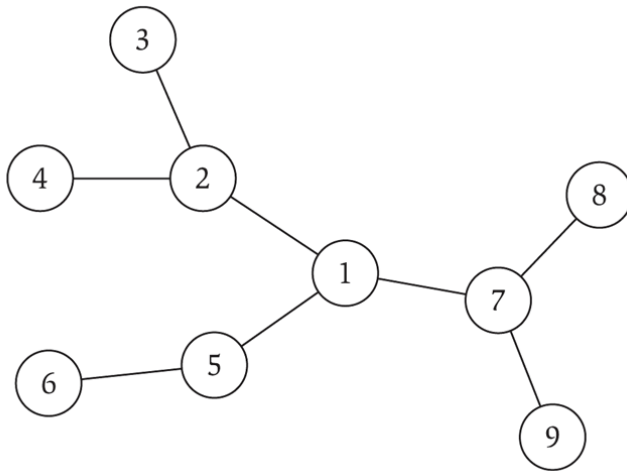
- G is connected.
- G does not contain a cycle.
- G has $n-1$ edges.



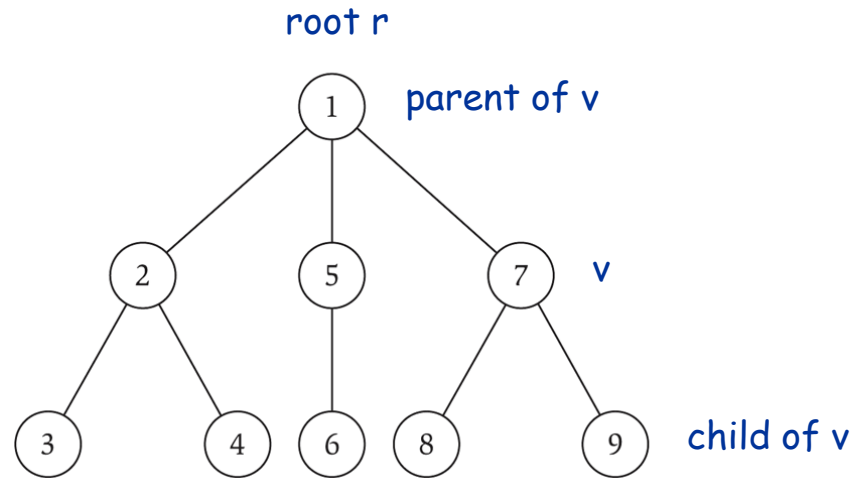
Rooted Trees

Rooted tree. Given a tree T , choose a root node r and orient each edge away from r .

Importance. Models hierarchical structure.



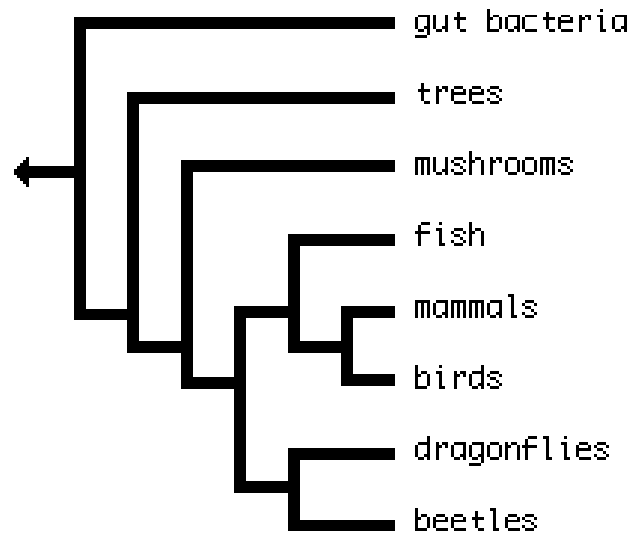
a tree



the same tree, rooted at 1

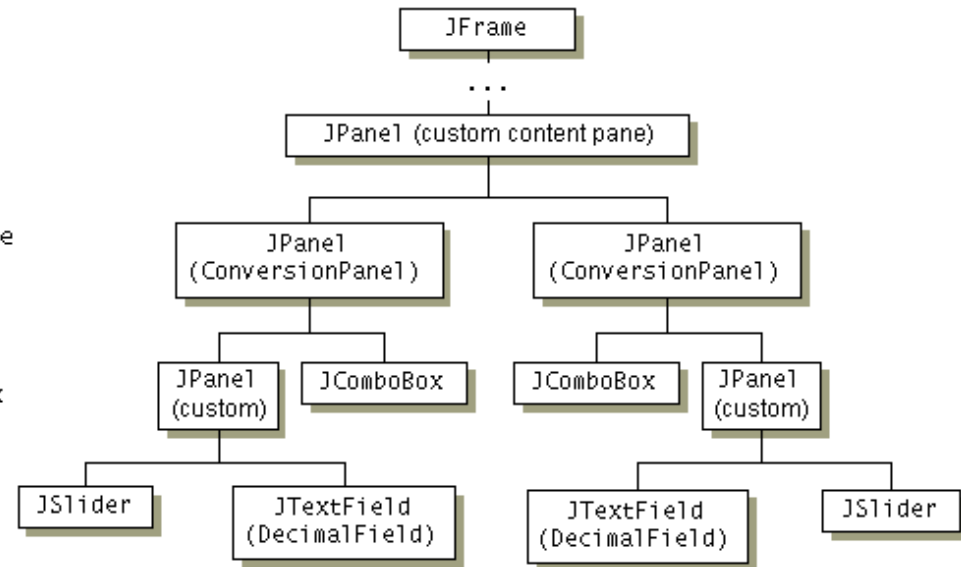
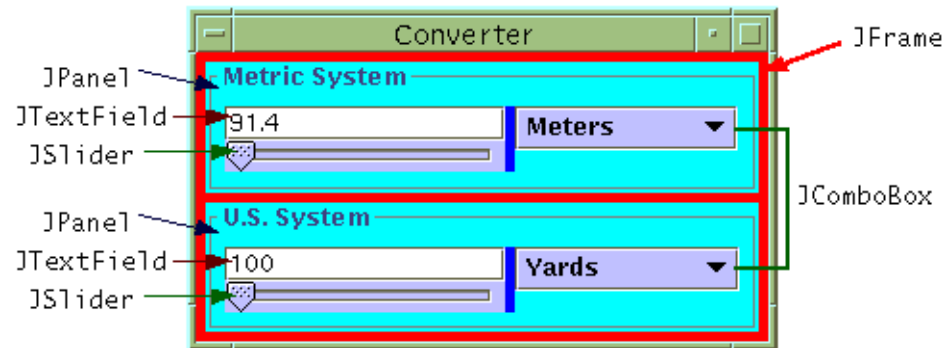
Phylogeny Trees

Phylogeny trees. Describe evolutionary history of species.



GUI Containment Hierarchy

GUI containment hierarchy. Describe organization of GUI widgets.



Reference: <http://java.sun.com/docs/books/tutorial/uiswing/overview/anatomy.html>