

VLSI Test Technology and Reliability (ET4076)



Lecture 5

Combinational Circuit Test Generation

(Chapter 7)

Said Hamdioui

Computer Engineering Lab
Delft University of Technology
2009-2010

Learning aims of today's lecture

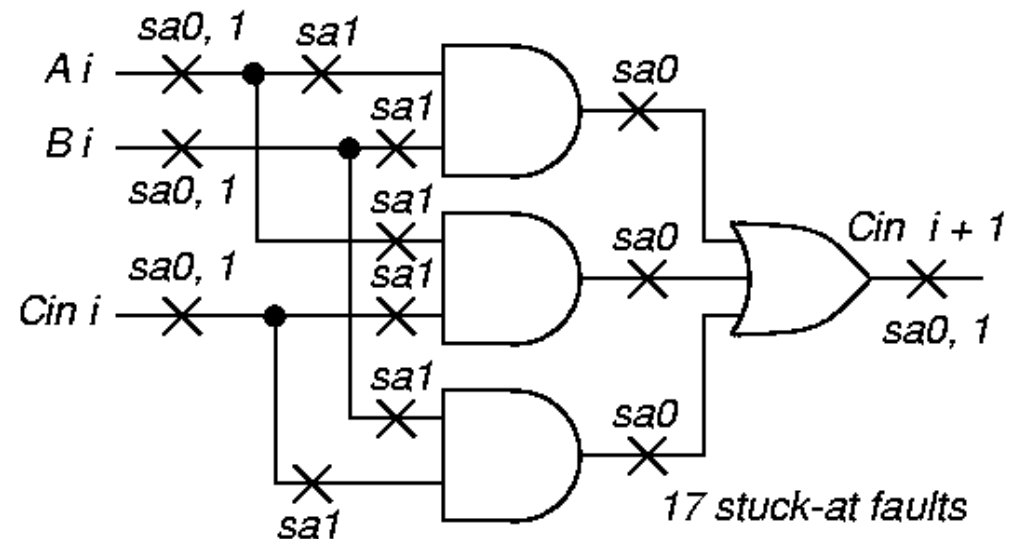
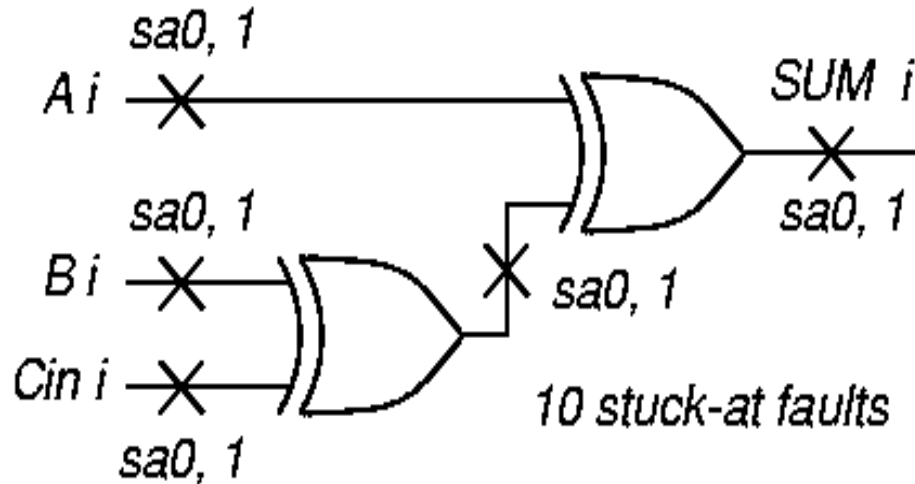
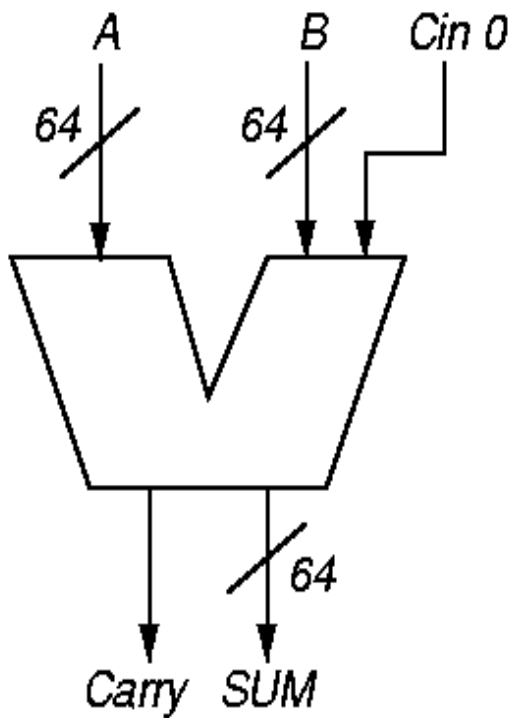
- Be able to
 - Define the different types of ATPG (concept, advantages, disadvantages)
 - Apply ATPG algorithms on combinational circuits
 - Explain the difference between the different path sanitization ATPG algorithms (D, PODEM, FAN)

Contents

- Algorithms and representations
 - Structural vs. functional test
 - Definition of ATPG
 - Search spaces
 - Algorithm completeness
 - ATPG algebra
 - Types of Algorithms
- Redundancy Identification
- Testing as a global problem
- ATPG algorithms
 - D-Calculus & D-Algorithm
 - PODEM Algorithm
 - FAN algorithm
- Advanced algorithms
- Test compaction

Functional vs. Structural ATPG

**Adder with 129 inputs
and 65 outputs**



Functional vs. Structural (Continued)

- **Functional ATPG** – generate complete set of tests for circuit input-output combinations
 - 129 inputs, 65 outputs:
 - $2^{129} = 680,564,733,841,876,926,926,749,214,863,536,422,912$ patterns
 - Using 1 GHz ATE, would take 2.15×10^{22} years
=> Not practical

- **Structural test:**
 - No redundant adder hardware, **64** bit slices
 - Each with **27** faults (using fault equivalence)
 - At most **$64 \times 27 = 1728$** faults (at the most 1728 tests)
 - Takes **0.000001728** s on 1 GHz ATE

- Designer gives small set of functional tests ($\sim 70\%$ FC)
 - Augment with structural tests to boost coverage to 98+ %

Definitions

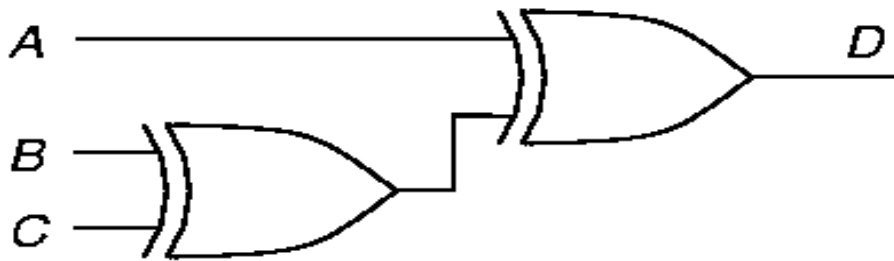
- **ATPG: Automatic Test pattern generator**
 - Operations on digital hardware:
 - Inject fault into circuit modeled in computer
 - Use various ways to activate and propagate fault effect through hardware to circuit output
 - Output flips from expected to faulty signal

- ***Electron-beam (E-beam) test***
 - Observes internal signals – “picture” of nodes charged to 0 and 1 in different colors
 - No need to propagate the fault effect to the PO
 - Too expensive

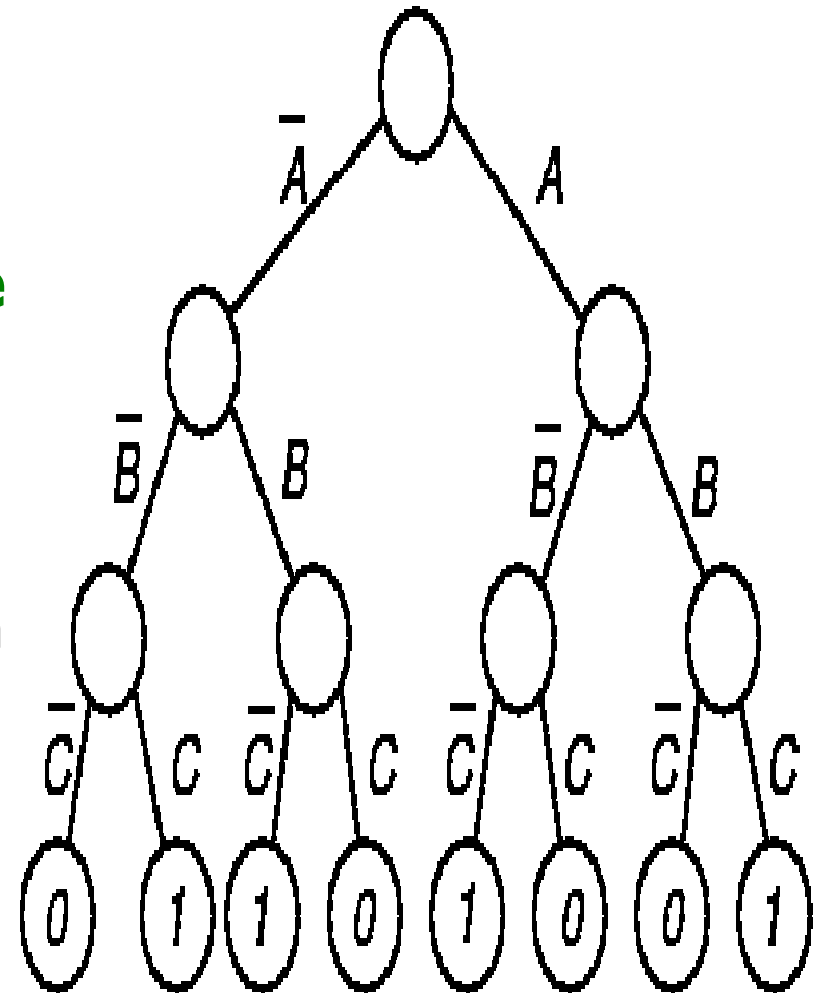
- ***Scan design***
 - Add test hardware to all FFs to make them a shift Register
 - Can shift state in, scan state out (test mode)
 - Widely used: makes sequential test combinational
 - Costs: 5 to 20% chip area, circuit delay, extra pin, longer test sequence

Search space abstractions

Binary Search tree

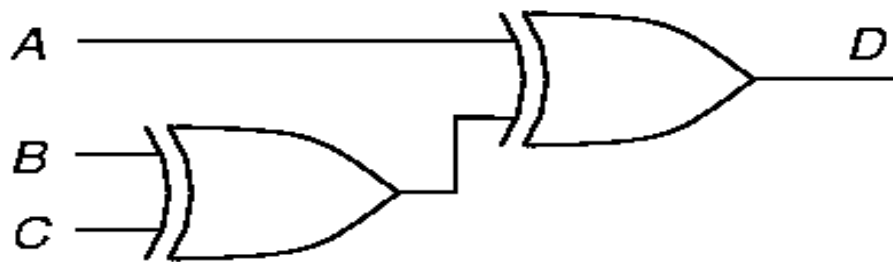


- The corresponding **binary tree**
 - Gives all possible input patterns
 - Leaf: labeled with good output
 - ATPG implicitly search the tree to find test pattern
 - Worst case: complete examination
 - Exponential rise
 - # of leaves: $2^{\#PI}$



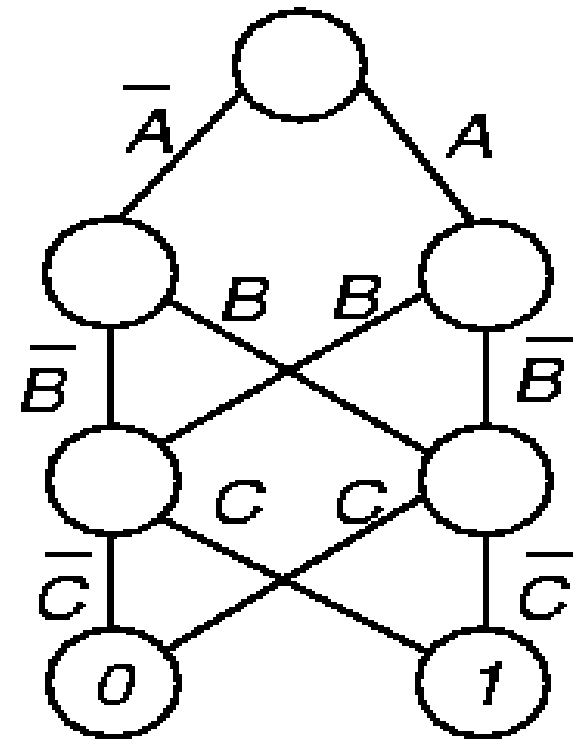
Search space abstractions

Binary decision tree



□ Binary decision tree

- Follow path from root/source to sink node – product of literals along path gives Boolean value at sink
- Rightmost path: $A B^* C^* = 1$
- Leftmost path: $A^* B^* C^* = 0$
- Problem: Size varies greatly with variable order



Algorithm Completeness

- An algorithm is **complete** if it ultimately can search entire binary decision tree, if needed, to generate a test
 - Completeness is important, otherwise ATPG may not achieve the required fault coverage
- **Untestable fault** – no test for it even after entire tree searched
 - Thus circuit behaves **correct** even in the presence of the fault
- Combinational circuits only – untestable faults are **redundant**, showing the presence of unnecessary hardware

ATPG Algebra

- Roth's 5-Valued and Muth's 9-Valued
- ATPG algebra: higher order Boolean set of notation to presents both **good** and **failing** circuit **simultaneously**

Symbol	Meaning	Good machine	Failing machine	
D	1/0	1	0	Roth's Algebra [5 values]
D*	0/1	0	1	
0	0/0	0	0	
1	1/1	1	1	
X	X/X	X	X	
G0	0/X	0	X	Muth's additions [Extended unknowns]
G1	1/X	1	X	
F0	X/0	X	0	
F1	X/1	X	1	

ATPG Algebra

Roth's and Muth's Higher-Order Algebras

- Represent two machines, which are simulated simultaneously by a computer program:
 - Good circuit machine (1st value)
 - Bad circuit machine (2nd value)

- Better to represent both in the algebra:
 - Need only 1 pass of ATPG to solve both

- Needed for complete ATPG:
 - Combinational: Roth Algebra, Multi-path sensitization,
 - Sequential: Muth Algebra -- good and bad machines may have different initial values due to fault

Type of algorithms

- Exhaustive
 - Expensive
 - Not practical, unless circuit partitioned

- Random-Pattern Generation
 - Easy to implement
 - Does not realize higher fault coverage

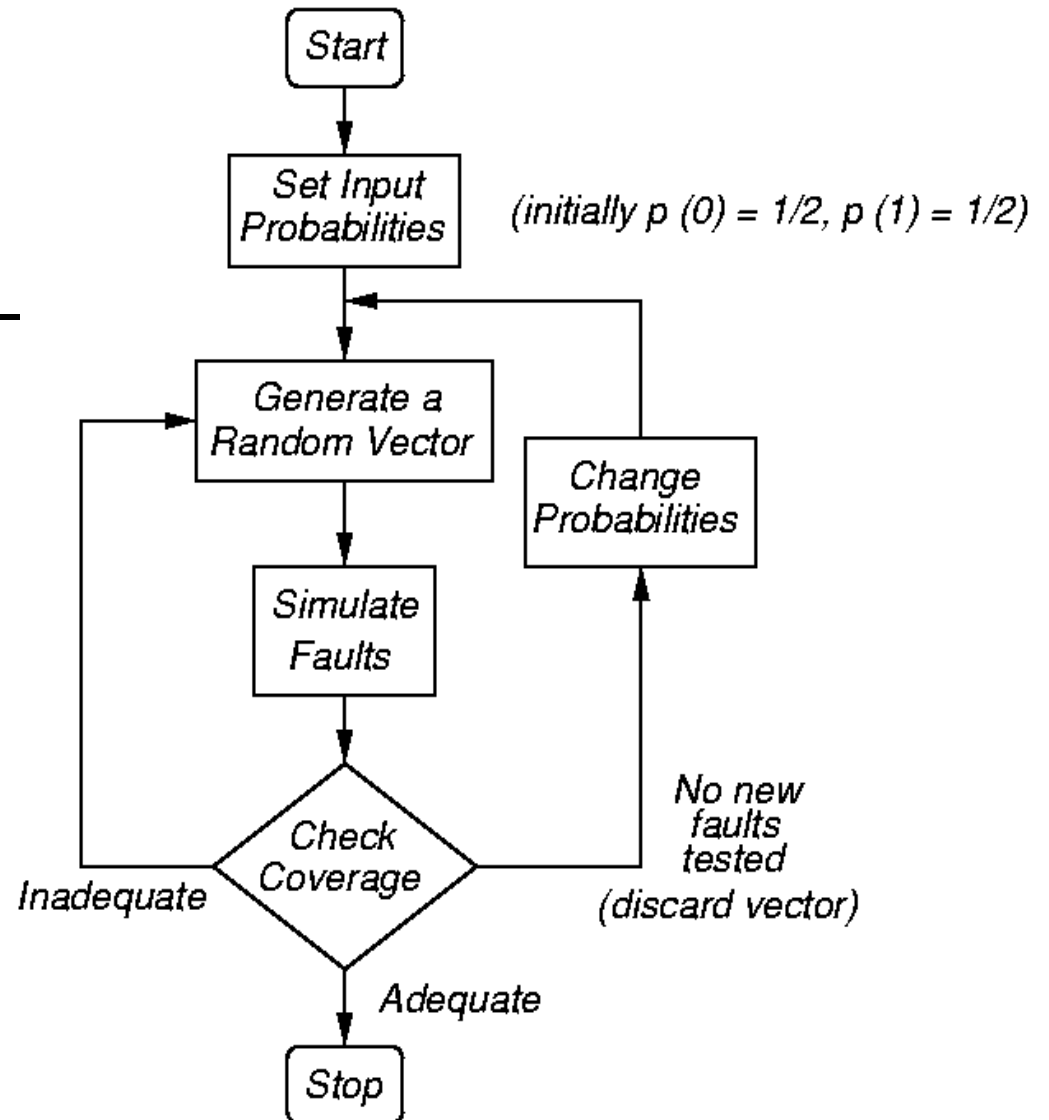
- Path Sensitization Method
 - Preferred ATPG method

Type of algorithms.... Exhaustive

- For n -input circuit, generate all 2^n input patterns
- Infeasible, unless circuit is **partitioned** into cones of logic, with ≤ 15 inputs
 - Perform exhaustive ATPG for each cone
 - Misses faults that require specific activation patterns for multiple cones to be tested

Type of algorithms....Random-Pattern Generation

- Flow chart for method
- Use to get tests for 60-80% of faults, then switch to D-algorithm or other ATPG for rest



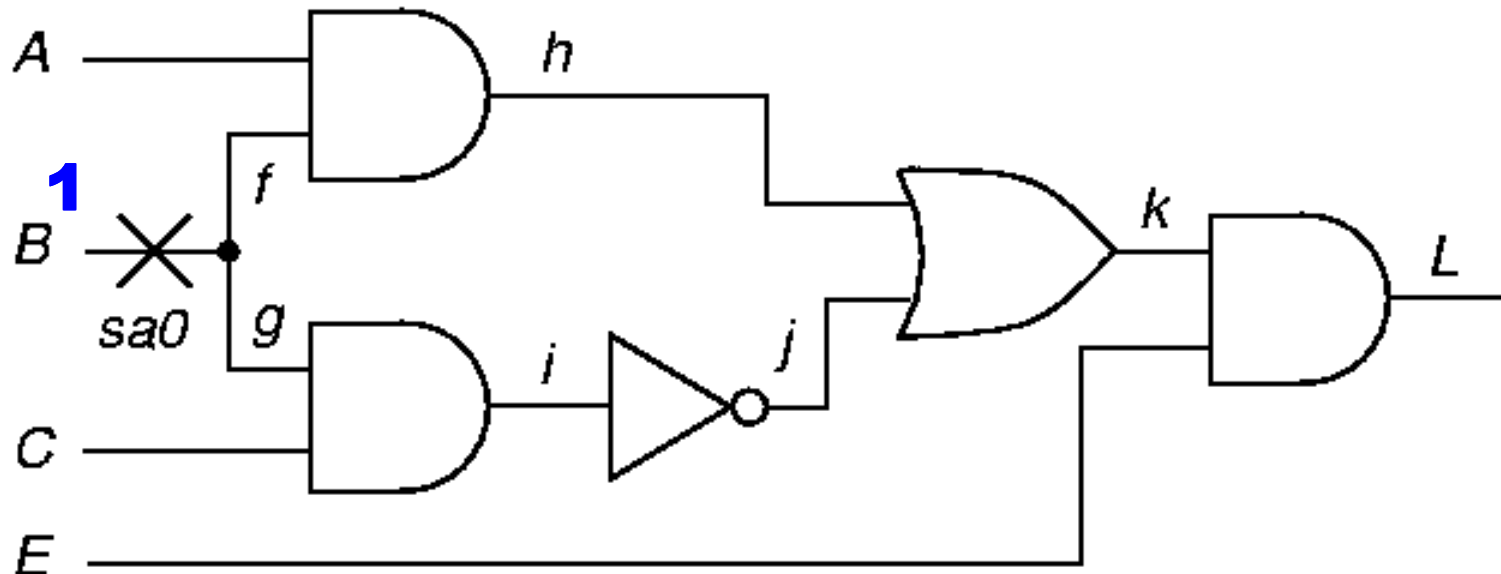
Type of algorithms..... Path Sensitization Method

- Is the preferred ATPG method
- Approach based on **three** steps
 1. **Fault sensitization**
 - Line SAF is activated by forcing it to an opposite value from the fault value
 - Also known as **Fault activation, fault excitation**
 2. **Fault propagation**
 - Fault effect is to be propagated through one or more paths to PO
 - Also known as **path sensitization**
 3. **Line justification**
 - Justify the sensitized fault by setting the PI of the circuits
- Steps 2 and 3 may find a **conflict**
 - ATPG has to **backup** or **backtrack**
 - i.e., discard previously made signal assignment and make an **alternative assignment**

Type of algorithms..... Path Sensitization Method

□ Example:

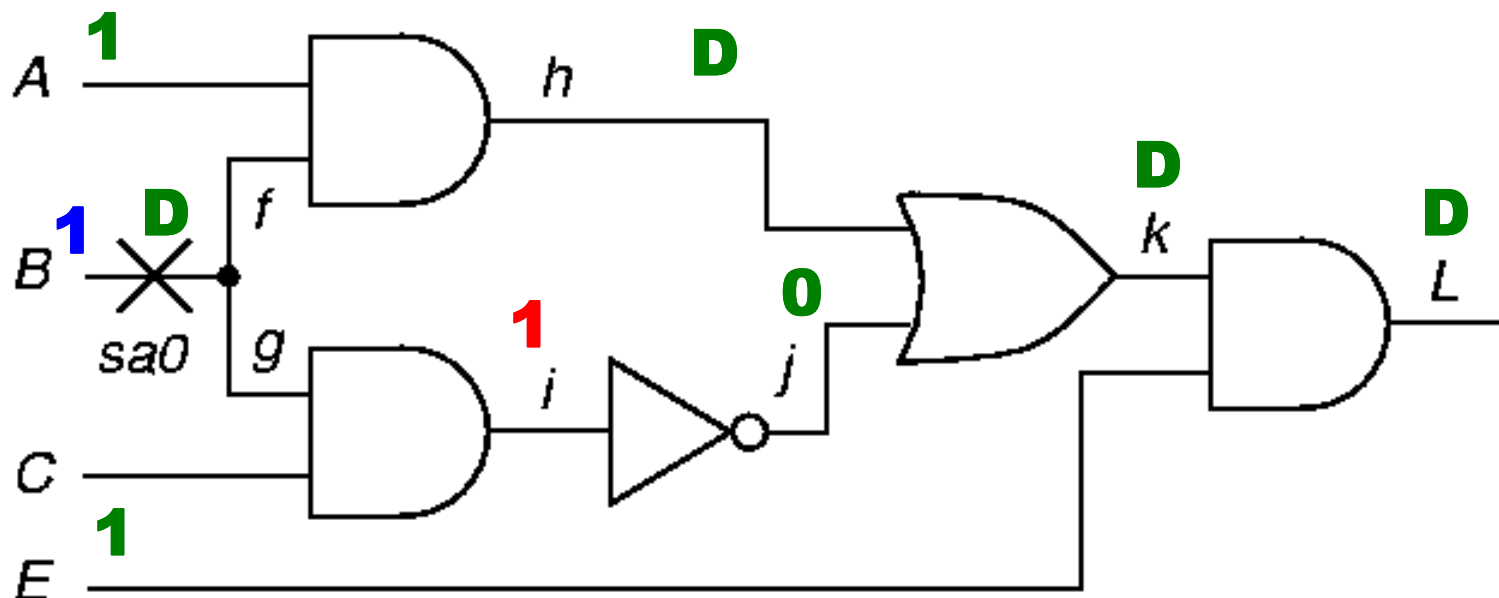
- **Fault sensitization:** set B to 1



Type of algorithms..... Path Sensitization Method

□ Example:

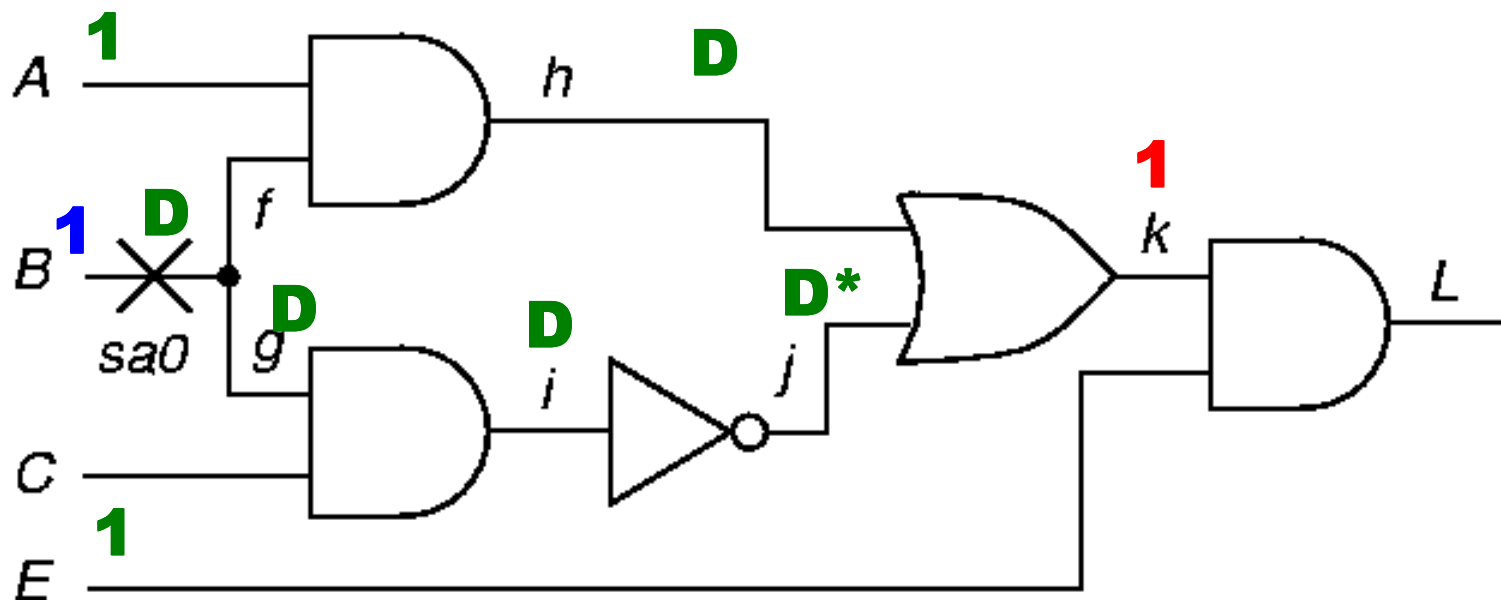
- **Fault sensitization:** set B to 1
 - **Fault propagation:** Try path $f - h - k - L$
 - **Line justification:** There is no way to justify the 1 on i
- ⇒ **Conflict**
- ⇒ Backtrack



Type of algorithms..... Path Sensitization Method

□ Example:

- **Fault sensitization:** set B to 1
- **Fault propagation:**
 - Try simultaneous paths $f - h - k - L$ and $g - i - j - k - L$
 - Blocked at k because *D-frontier* (D or D^*) **disappears**

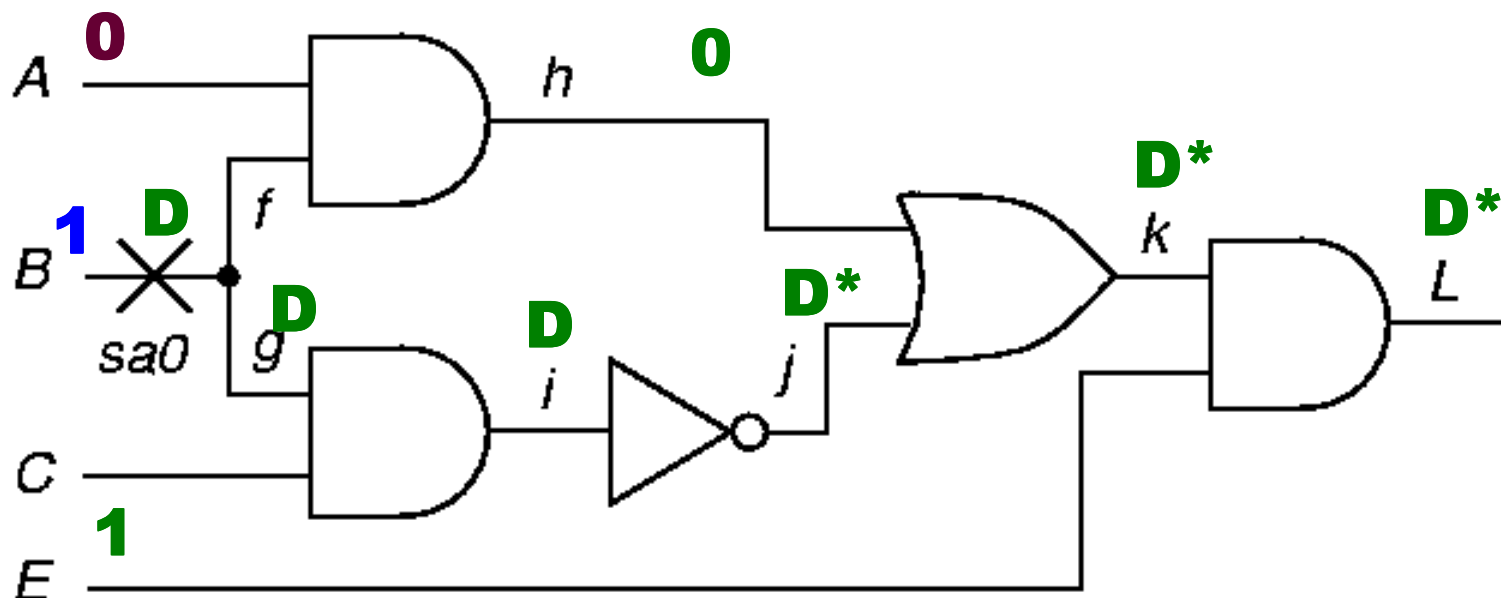


Type of algorithms..... Path Sensitization Method

□ Example:

- **Fault sensitization:** set B to 1
- **Fault propagation:** Try path $g - i - j - k - L$
- **Line justification:** set A to 0

=> **Test found and fault detected**



ATPG Computational Complexity

- Ibarra and Sahni analysis [1975] – ***NP-Complete***
 - No polynomial expression found for compute time, presumed to be exponential

- Worst case:
 - *no_pi* inputs, 2^{no_pi} input combinations
 - *no_ff* flip-flops, 4^{no_ff} initial flip-flop states
 - (good machine 0 or 1 x bad machine 0 or 1)
 - work to forward simulate or reverse simulate all logic gates *n* rises proportional with *n*
 - Complexity: $O(n \times 2^{no_pi} \times 4^{no_ff})$

ATPG Computational Complexity (Cnt)

□ History ATPG

- Improve heuristic algorithms & procedures to:
 - Find all necessary signal assignments as early as possible
 - Search as little as possible of the decision space
 - Worst case complexity: $O(n \times 2^{no_pi} \times 4^{no_ff})$

□ Fault simulation

- For combinational circuits: $O(n^2)$
- For sequential circuits: estimated between $O(n^2)$ and $O(n^3)$
 - Based on empirical measurements

⇒ whenever possible use fault simulation to avoid ATPG computations

E.g.,

- Use RPG and fault simulation to get tests
- Use ATPG for hard-to-test faults

ATPG Computational Complexity* (Cnt)

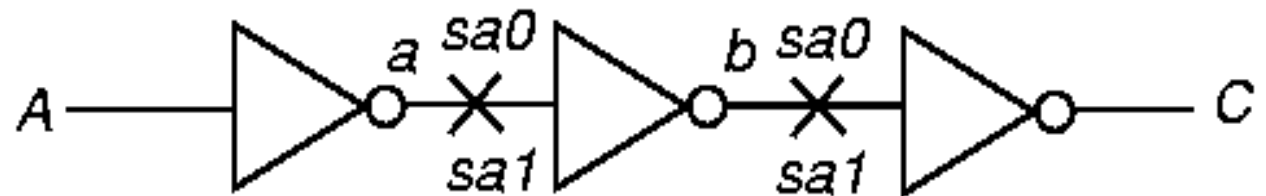
ATPG for analog:

Analog Fault Modeling Impractical for Logic ATPG

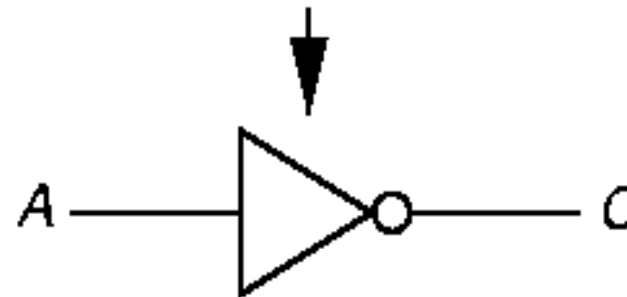
- Problems with modeling actual defects in analog circuits
 - Huge # of different possible analog faults in digital circuit
 - Exponential complexity of ATPG algorithm – a 20 flip-flop circuit can take days of computing
 - Cannot afford to go to a lower-level model
 - Most test-pattern generators for digital circuits cannot even model at the transistor switch level
 - More complex
 - (see textbook for 5 examples of switch-level ATPG)

Redundancy identification... Untestable faults(1)

- Combinational ATPG can find redundant (unnecessary) hardware
- Untestable faults in combinational circuits indicates redundant hardware



- Fault Test
- a sa1, b sa0* $A = 1$
- a sa0, b sa1* $A = 0$



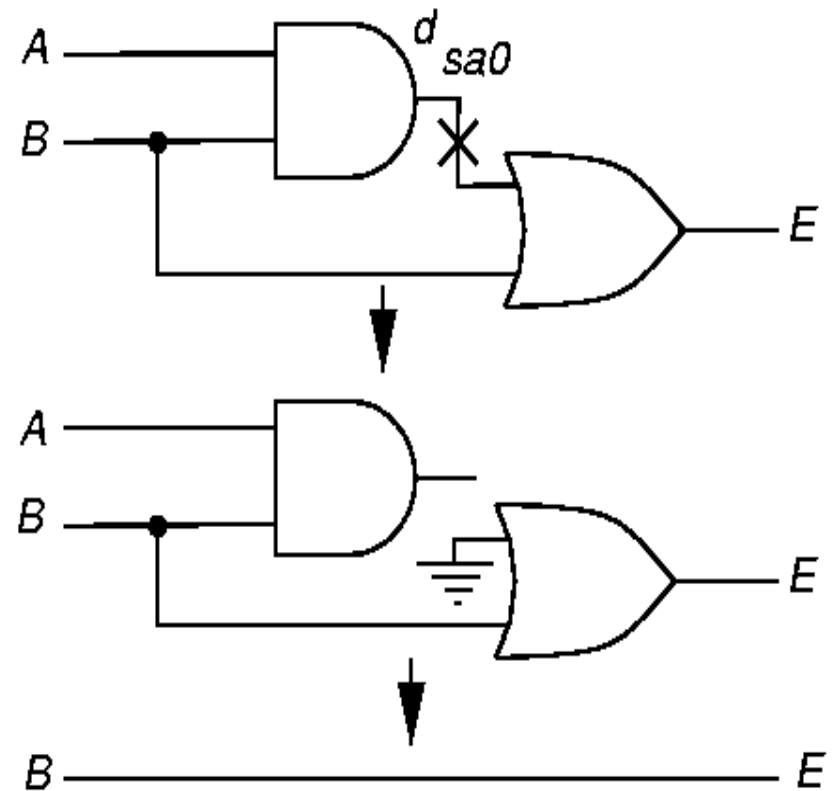
- All faults are testable
- Therefore, these faults are not redundant

Redundancy identification... Untestable faults(2)

Test fault: s-a-0 on line d

- Set A & B to 1
- However, E will be 1
- Fault cannot be propagated to the output
 - => Untestable fault
 - => Redundant Hardware

- Remove the redundant hardware
 - d always 0: ground it permanently
 - Remove OR gate and replace with a wire
 - And discard AND gate and input A



Redundancy identification Untestable faults(3)

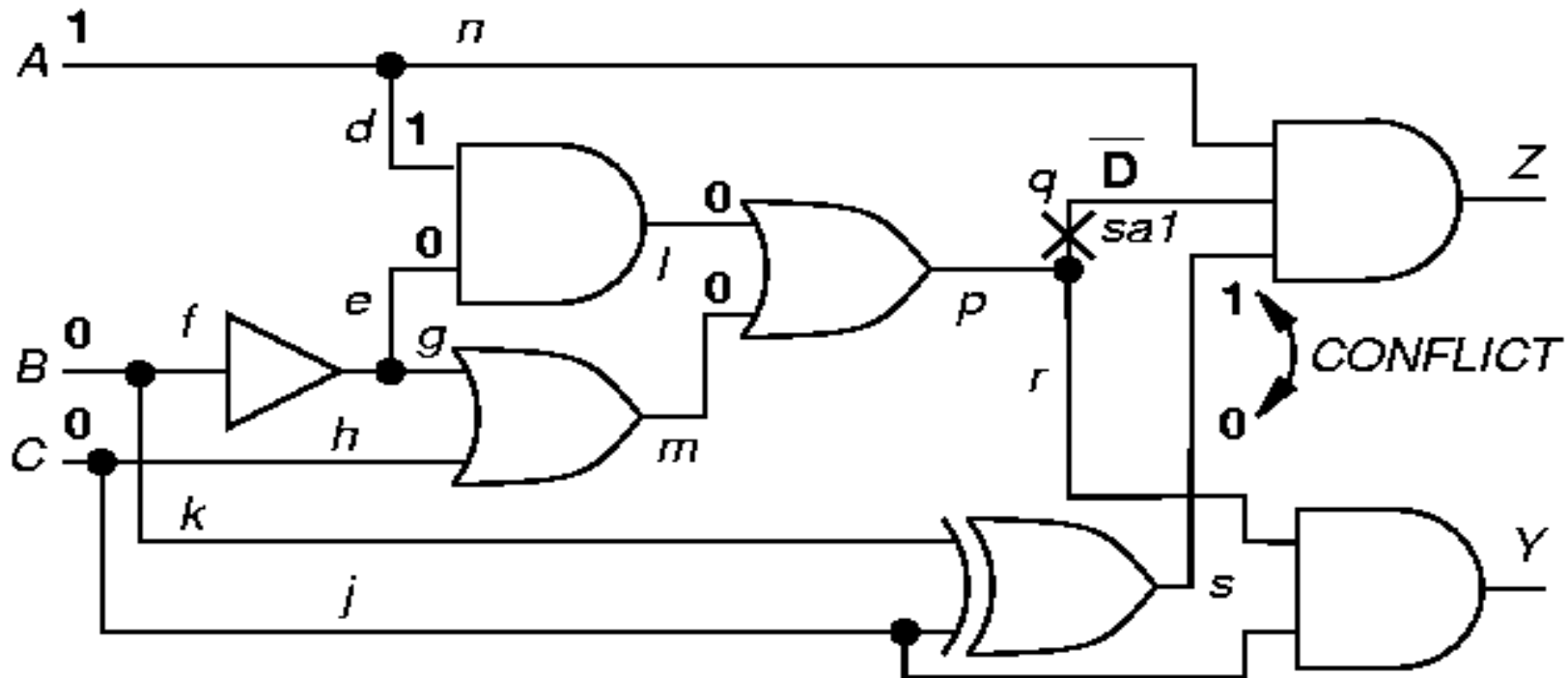
□ Test fault q s-a-1

- **Fault sensitization**: set q to 0
- **Fault propagation**: set A to 1 and **s to 1**
- **Line justification**: set B to 0 and C to 0

⇒ **Conflict at s**

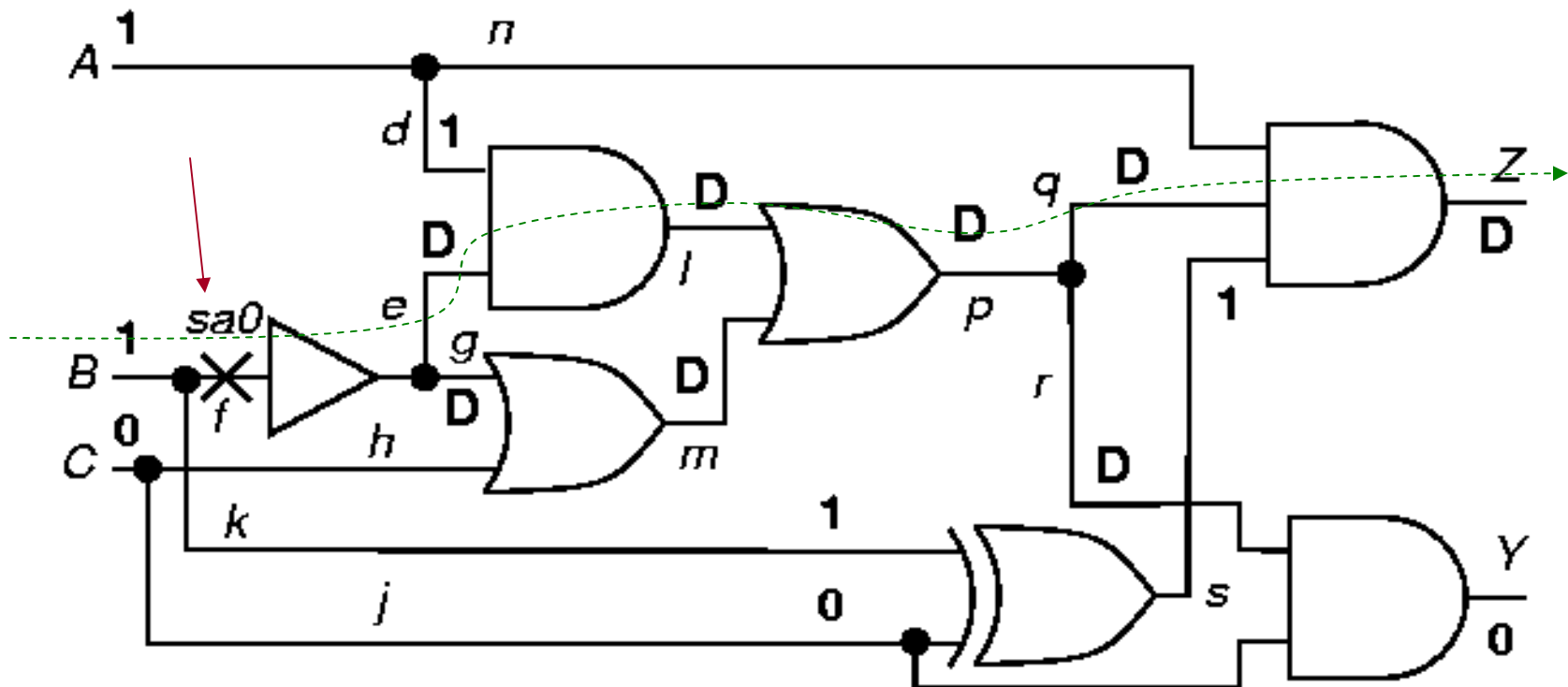
No other possible alternative assignments

⇒ **Redundant fault**



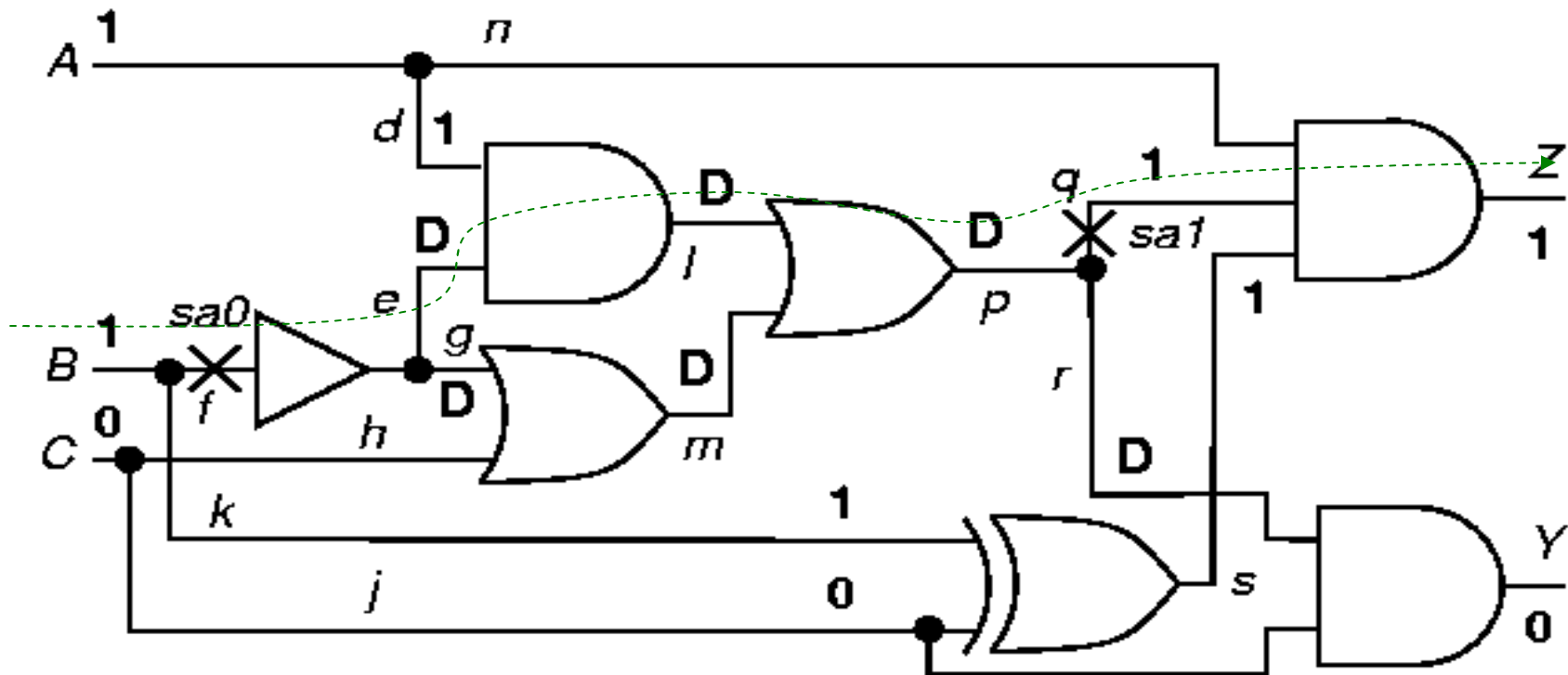
Redundancy identification Fault masking (1)

- **f** s-a-0 tested when fault **q** s-a-1 **not** there
 - Fault sensitization: set **f** to 1
 - Fault propagation (via **felqz**): set **d** to 1, **m** to 0, **s** to 1, **n** to 1
 - Line justification: **set A to 1, B to 1 and C to 0**
- ⇒ Fault detected at **Z**



Redundancy identification Fault masking (2)

- **f** s-a-0 tested when fault **q** s-a-1 also **present**
 - Fault sensitization: set f to 1
 - Fault propagation (via felqz): set d to 1, m to 0, ...
 - **Blocked** by q s-a-1 => Fault can **NOT** be detected
- ⇒ **q s-a-1 (redundant) masks f s-a-0**



Redundancy identification Hazards

Eliminates **hazards** in circuit output

□ $OUT0 = A.B + A*.C (+B.C)$

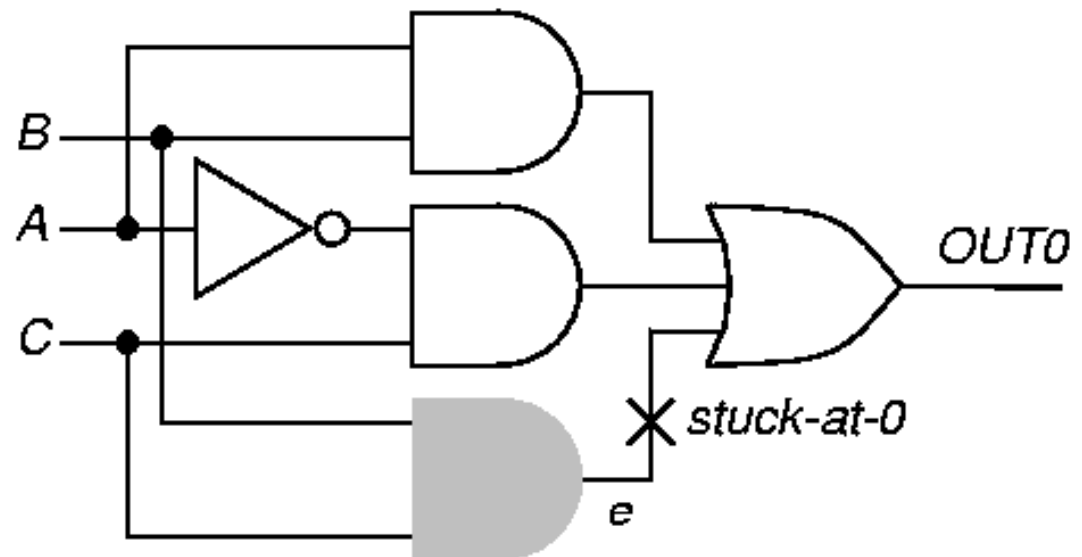
- B.C is added to prevent hazards (when $A 0 \rightarrow 1$)

□ **Redundant** fault s-a-0 on line e

- No effect on the function
- But **masks** any testable fault using line e to propagate to OUT
- Additional area and power due to redundancy

		BC			
		00	01	11	10
A	0	0	1	1	0
	1	0	0	1	1

(a)



(b)

Redundancy identificationImpact

□ Redundancy impacts:

- Performance
- Power consumption
- Area overhead
- Reliability
 - A redundant fault can **mask** the presence of other testable faults
 - Unwanted in systems critical to human safety or wealth!!

□ How to deal with it?

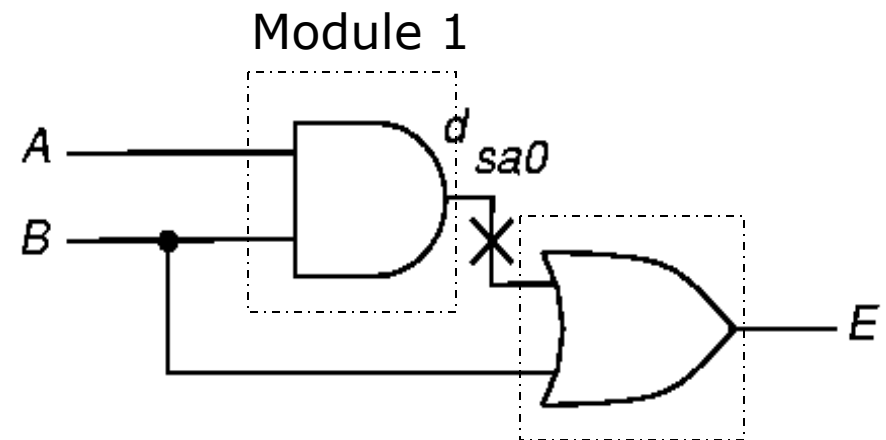
- Impossible to do that manually (e.g., Millions of gates)
- Commercial tools
 - Synthesis moderately-sized designs to irredundant hardware
 - Big designs have to be partitioned before synthesis
 - May introduce redundant hardware
 - ATPG is one of the best ways to find this redundancy

Testing as a global problem

- Testing is a **global problem**
- Combination of **fully testable** modules in a logic circuit are
 - **Not** necessary fully testable
 - May be not be testable with the same patterns that would test the modules individually

Example

- Module 1
 - Requires $AB = 11$ for d s-a-0
- The test cannot be used for the entire circuit



ATPG algorithms

- Definitions
- D-Algorithm (Roth) -- 1966
 - D-cubes
 - Bridging faults
 - Logic gate function change faults
- PODEM (Goel) -- 1981
 - X-Path-Check
 - Backtracing
- Summary

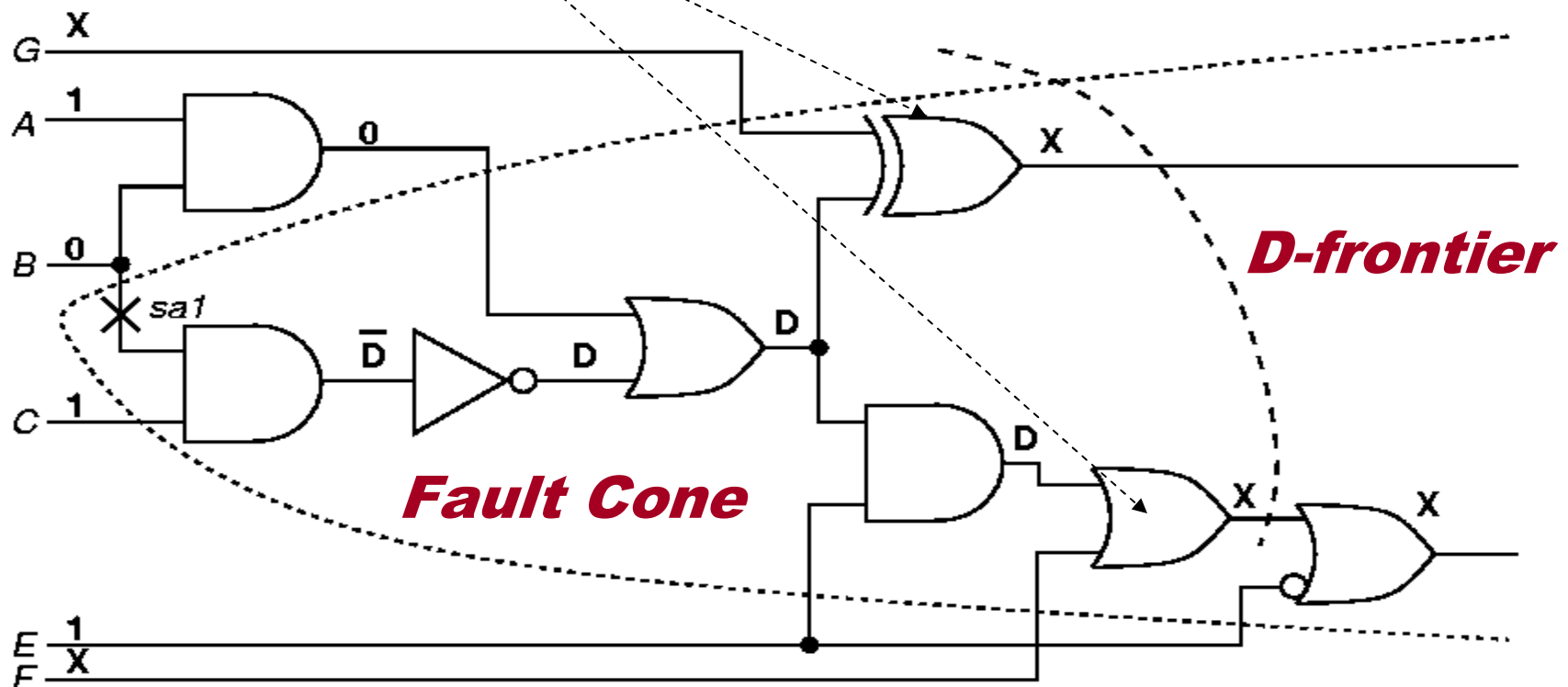
ATPG algorithms... History & Speedups

Algorithm	Est. speedup over D-ALG	Year	
D-ALG	1	1966	
PODEM	7	1981	
FAN	23	1983	
TOPS	292	1987	
SOCRATES	1574	ATPG System	1988
Waicukauski et al.	2189	ATPG System	1990
EST	8765	ATPG System	1991
TRAN	3005	ATPG System	1993
Recursive learning	485		1995
Tafertshofer et al.	25057		1997

ATPG algorithms.....Definitions (1)

Fault Cone and D-frontier

- **Fault Cone**: Set of hardware affected by fault
- **D-frontier**: Set of all gates with D or D* at inputs and X at the output



ATPG algorithms.....Definitions(2)

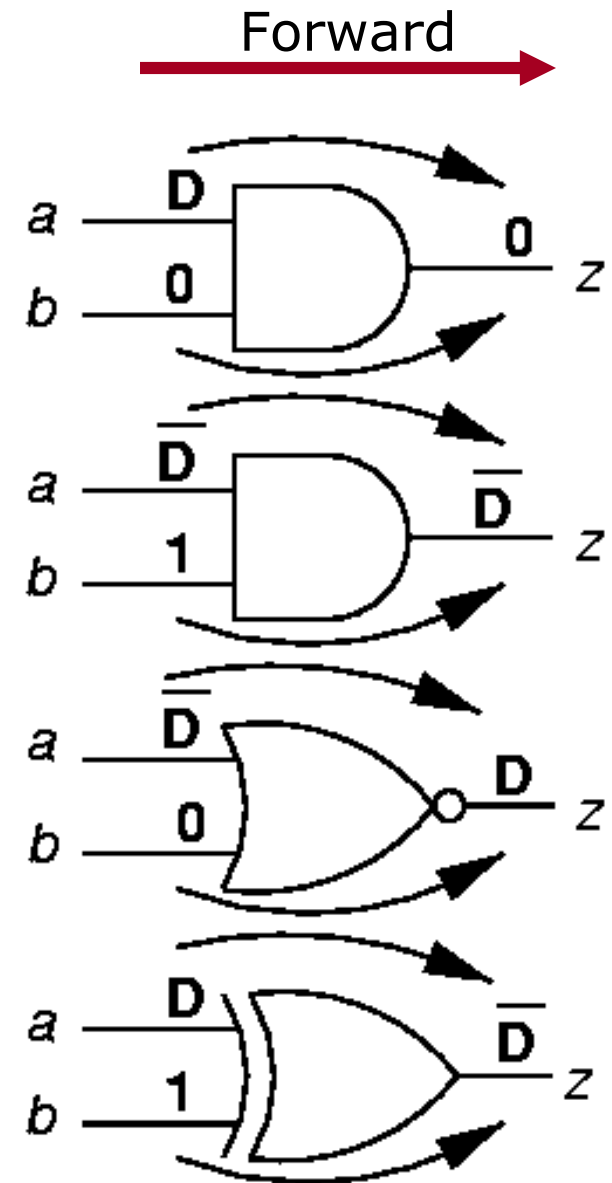
Forward Implication

Results in logic gate inputs that are significantly labeled so that output is **uniquely** determined

Example:

AND gate forward implication table:

a \ b	0	1	X	D	\bar{D}
0	0	0	0	0	0
1	0	1	X	D	\bar{D}
X	0	X	X	X	X
D	0	D	X	D	0
\bar{D}	0	\bar{D}	X	0	\bar{D}

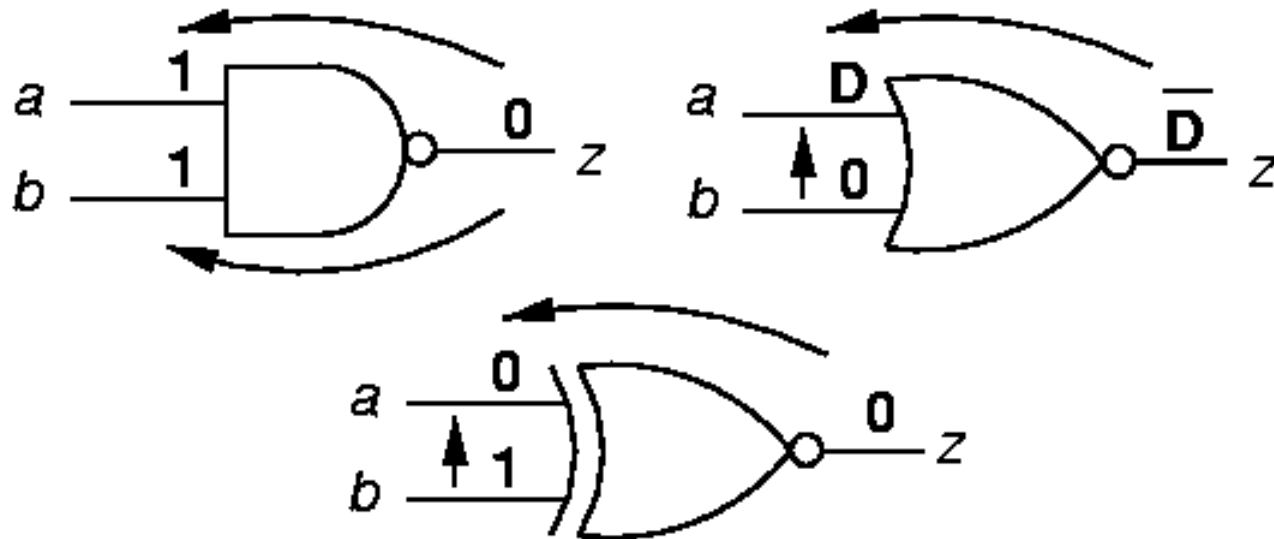


ATPG algorithms.....Definitions(3)

Backward Implication

- **Unique** determination of all gate inputs when the gate output and *some* of the inputs are given

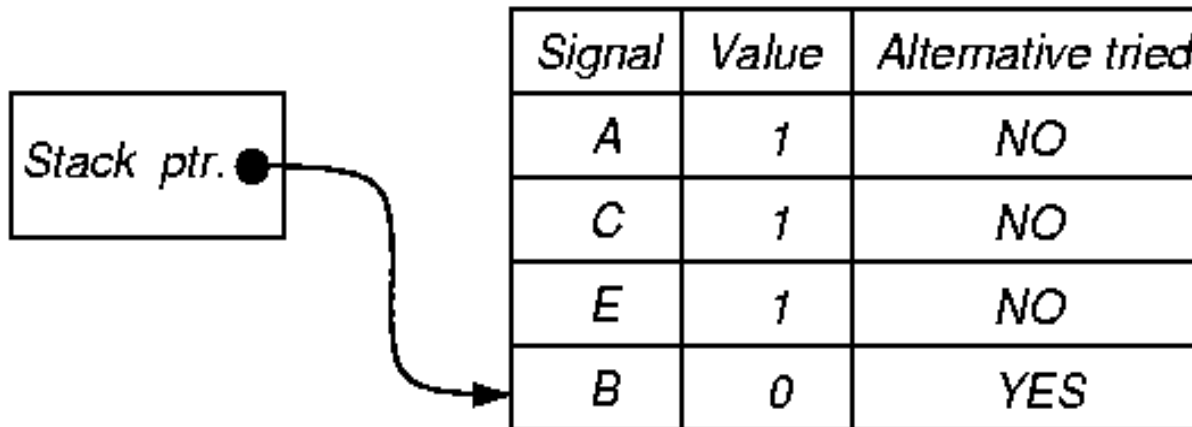
- Example



ATPG algorithms.....Definitions(4)

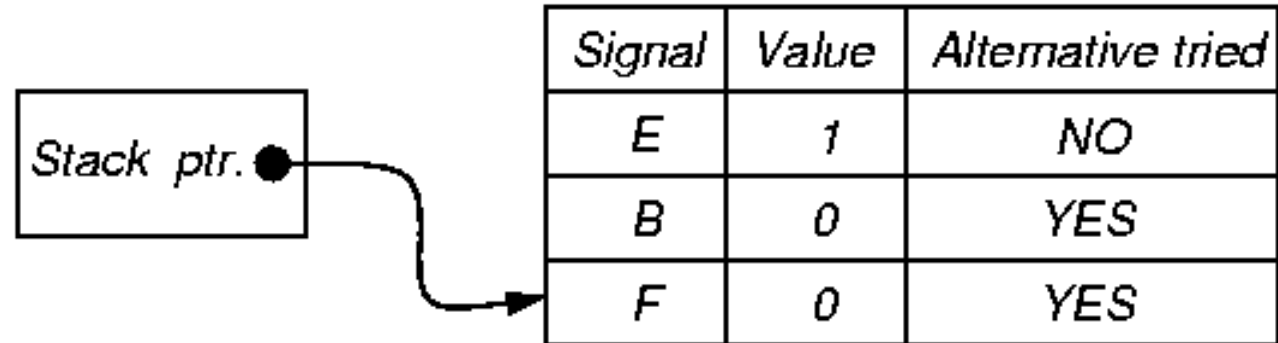
Implication Stack

- It is a Push-down stack; it records:
 - Each signal set in circuit by ATPG
 - Whether alternate signal value already tried
 - Portion of binary search tree already searched

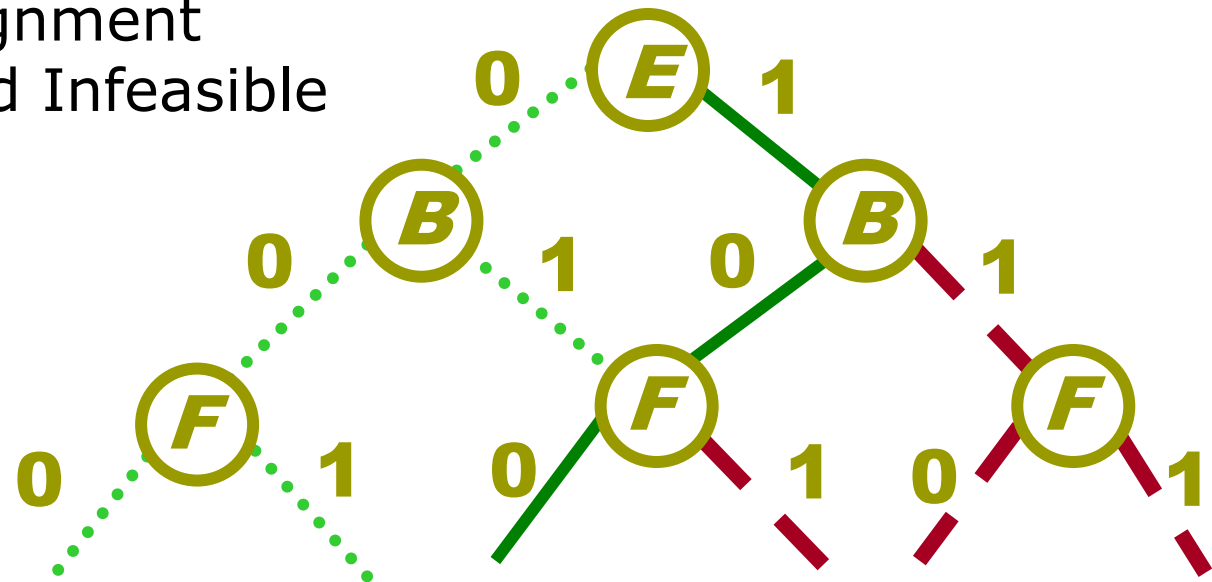


ATPG algorithms.....Definitions(5)

Implication Stack after backtrack



- Unexplored
- Present Assignment
- Searched and Infeasible



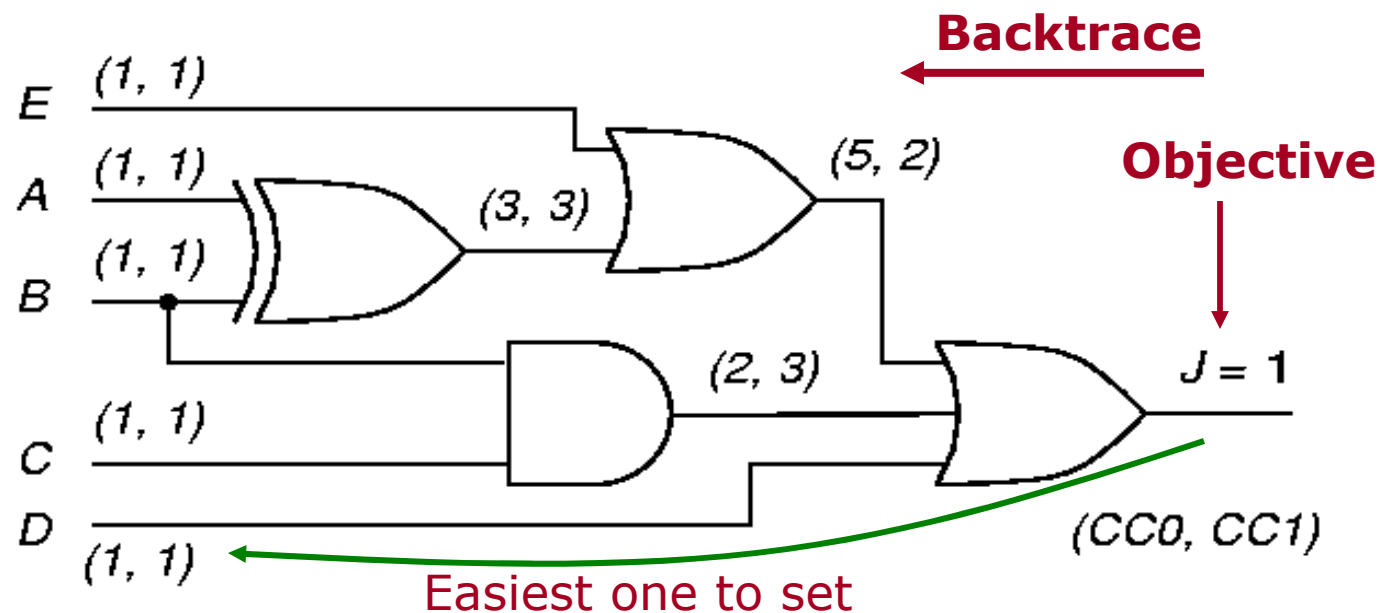
ATPG algorithms.....Definitions(6)

□ **ATPG Objective**

- Desired signal value to be achieved
- Guides it away from infeasible/hard solutions

□ **Backtrace**

- Determines which primary input and value to set to achieve objective
- Use testability measures



Branch-and-Bound Search

- Efficiently searches binary search tree
 - **Branching** – At each tree level, selects which input variable to set to what value
 - **Bounding** – Avoids exploring large tree portions by artificially restricting search decision choices
 - Complete exploration is impractical
 - Uses **heuristics**

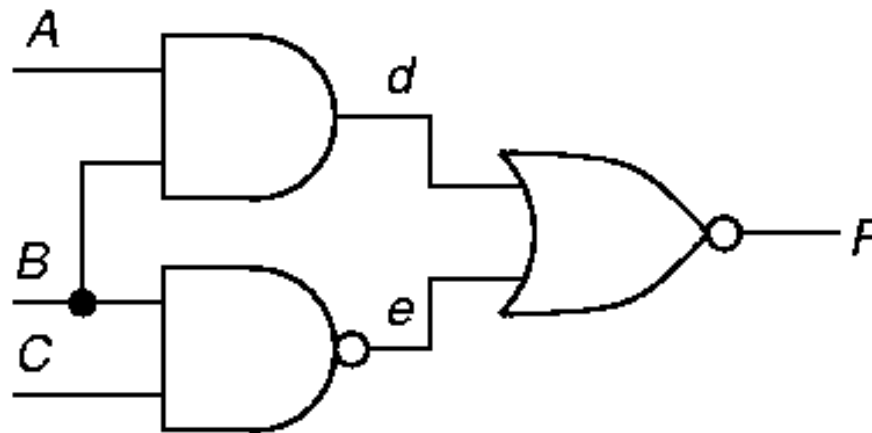
D-Algorithm..... Roth IBM (1966)

- Fundamental concepts invented:
 - First complete ATPG algorithm
 - *D-Cube*
 - *D-Calculus*
 - *Implications* – forward and backward
 - *Implication stack*
 - *Backtrack*
 - Test Search Space

D-Algorithm..... Definitions(1)

□ Singular Cover:

Minimal set of logic signal assignments to show *essential prime implicants* of *Karnaugh map*



Gate	Inputs	Output	Gate	Inputs	Output		
AND	A	B	d	NOR	d	e	F
1	0	X	0	1	1	X	0
2	X	0	0	2	X	1	0
3	1	1	1	3	0	0	1

D-Algorithm..... Definitions(2)

D-Cube: Collapsed truth table entry to characterize logic block

- Use Roth's 5-valued algebra
- Can change all **D's** to **D*'s** and **D*'s** to **D's** (do both)

□ Example: AND gate **propagation** :

	A	B	d
Combine rows 3 (good) & 1(faulty)	D	1	D
Reverse inputs	1	D	D
And two cubes	D	D	D
Interchange D and D*	D*	D*	D*
	1	D*	D*
	D*	1	D*

D-Algorithm..... Definitions(3)

D-Cube Operation of D-Intersection

- **D-intersection:** $0 \cap 0 = 0 \cap X = X \cap 0 = 0$
 $1 \cap 1 = 1 \cap X = X \cap 1 = 1$
 $X \cap X = X$

- ψ – undefined (same as ϕ)
- If both μ and λ occur, then cubes incompatible
- If only μ occur, then $D \cap D = D$ and $D^* \cap D^* = D^*$
- If only λ occur, then inverse D and D^* in second cube

- **D-containment:**

Cube a contains
 Cube b if b is a
 subset of a

\cap	0	1	X	D	D*
0	0	ϕ	0	ψ	ψ
1	ϕ	1	1	ψ	ψ
X	0	1	X	D	D*
D	ψ	ψ	D	μ	λ
D*	ψ	ψ	D*	λ	μ

D-Algorithm..... Definitions(4)

Primitive D-Cube of Failure

- Models circuit faults:
 - *Stuck-at-0*
 - *Stuck-at-1*
 - *Bridging fault* (short circuit)
 - Arbitrary change in logic function
- AND Output sa0: "1 1 D"
- AND Output sa1: "0 X D*" and "X 0 D*"
- Wire sa0: "D"; Wire sa1: "D*";

Test cube:

Refers to the set of PI, intermediate, and PO circuit signals that are set to get a test for the fault

Propagation D-cube: models conditions under which fault effect propagates through gate

D-Algorithm..... Definitions(5)

Implication Procedure

1. Model fault with appropriate *primitive D-cube of failure* (PDF)
 2. Select *propagation D-cubes* to propagate fault effect to a circuit output (*D-drive* procedure)
 3. Select *singular cover* cubes to justify internal circuit signals (*Consistency* procedure)
 - If Cube intersection fail, back up to the last decision point and select an alternative cube
- Put signal assignments in *test cube*
 - Regrettably, cubes are selected very **arbitrarily** by D-ALG

D-AlgorithmTop Level

1. Number all circuit lines in increasing level order from PIs to POs;
2. Select a primitive D-cube of the fault to be the *test cube*;
 - Put logic outputs with inputs labeled as D (D*) onto the ***D-frontier***;
//[provoke the fault and D-cube]
3. *D-drive* (); // [Propagate D or D* to PO]
4. *Consistency* (); // [Tracking backwards to PI]
5. return ();

D-Algorithm*D-drive*

```
while (untried fault effects on D-frontier)
  select next untried D-frontier gate for propagation;
  while (untried fault effect fanouts exist)
    select next untried fault effect fanout;
    generate next untried propagation D-cube;
    D-intersect selected cube with test cube;
    if (intersection fails or is undefined) continue;
    if (all propagation D-cubes tried & failed) break;
    if (intersection succeeded)
      add propagation D-cube to test cube -- recreate D-frontier;
      Find all forward & backward implications of assignment;
      save D-frontier, algorithm state, test cube, fanouts, fault;
      break;
    else if (intersection fails & D and D* in test cube) Backtrack ();
    else if (intersection fails) break;
  if (all fault effects unpropagatable) Backtrack ();
```

D-Algorithm.....*Consistency*

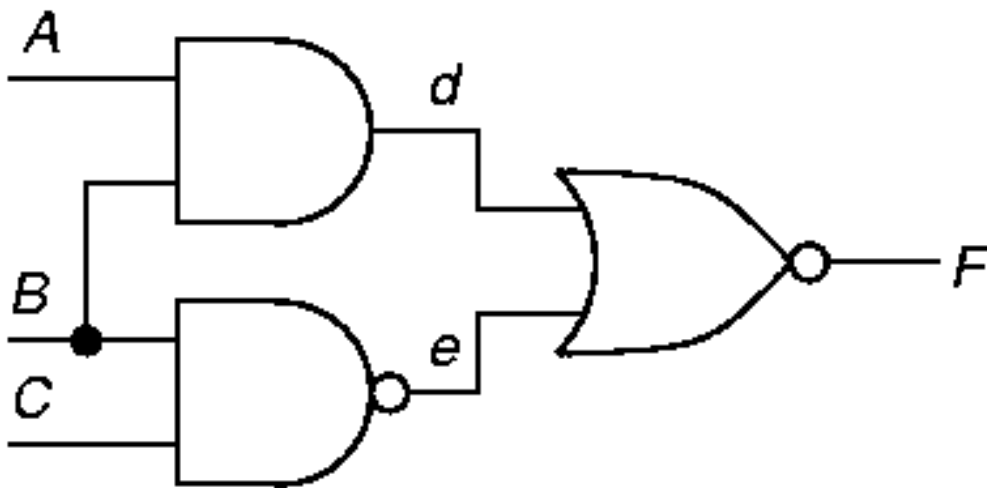
g = coordinates of test cube with 1's & 0's;
if (g is only PIs) **fault testable & stop**;
for (each unjustified signal in g)
 Select highest # unjustified signal z in g , not a PI;
 if (inputs to gate z are both **D** and **D***) break;
 while (untried singular covers of gate z)
 select next untried singular cover;
 if (no more singular covers)
 If (no more stack choices) **fault untestable & stop**;
 else if (untried alternatives in *Consistency*)
 pop implication stack -- try alternate assignment;
 else
 Backtrack ();
 D-drive ();
 If (singular cover D-intersects with z) delete z from g , add inputs
 to singular cover to g , *find all forward and backward implications
 of new assignment*, and break;
 If (intersection fails) mark singular cover as failed;

D-Algorithm..... *Backtrack*

```
if (PO exists with fault effect) Consistency ();  
else pop prior implication stack setting to try  
    alternate assignment;  
if (no untried choices in implication stack)  
    fault untestable & stop;  
else return;
```

D-Algorithm..... *Example 1(1)*

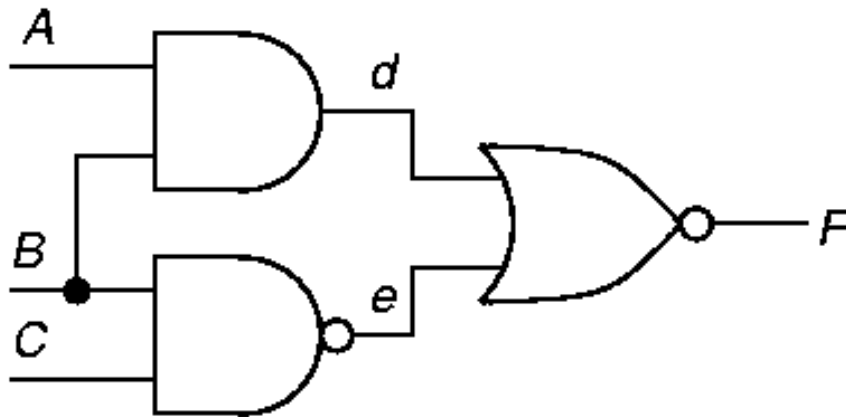
- Circuit Example 7.1 and Truth Table



Inputs			Output
<i>a</i>	<i>b</i>	<i>c</i>	<i>F</i>
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	0

D-Algorithm..... Example 1(2)

- **Singular cover** – Used for **justifying** lines

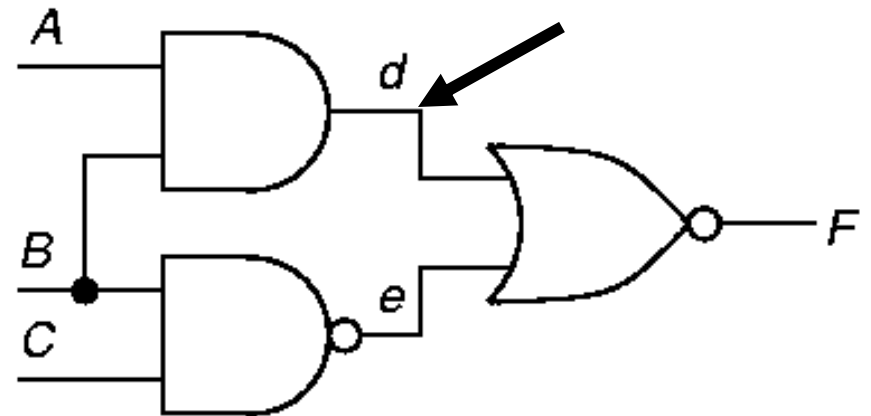


- **Propagation D-cubes** – Conditions under which difference between good/failing machines **propagates**

A	B	C	d	e	F
1	1		1		
0			0		
	0	1		0	
	1	0		1	
			1	1	0
			0	0	0
					1
D	1		D		
1	D		D		
D	D		D		
	D	1		D*	
	1	D		D*	
	D	D		D*	
			D	0	D*
			0	D	D*
			D	D	D*

D-Algorithm..... Example 1 (3)

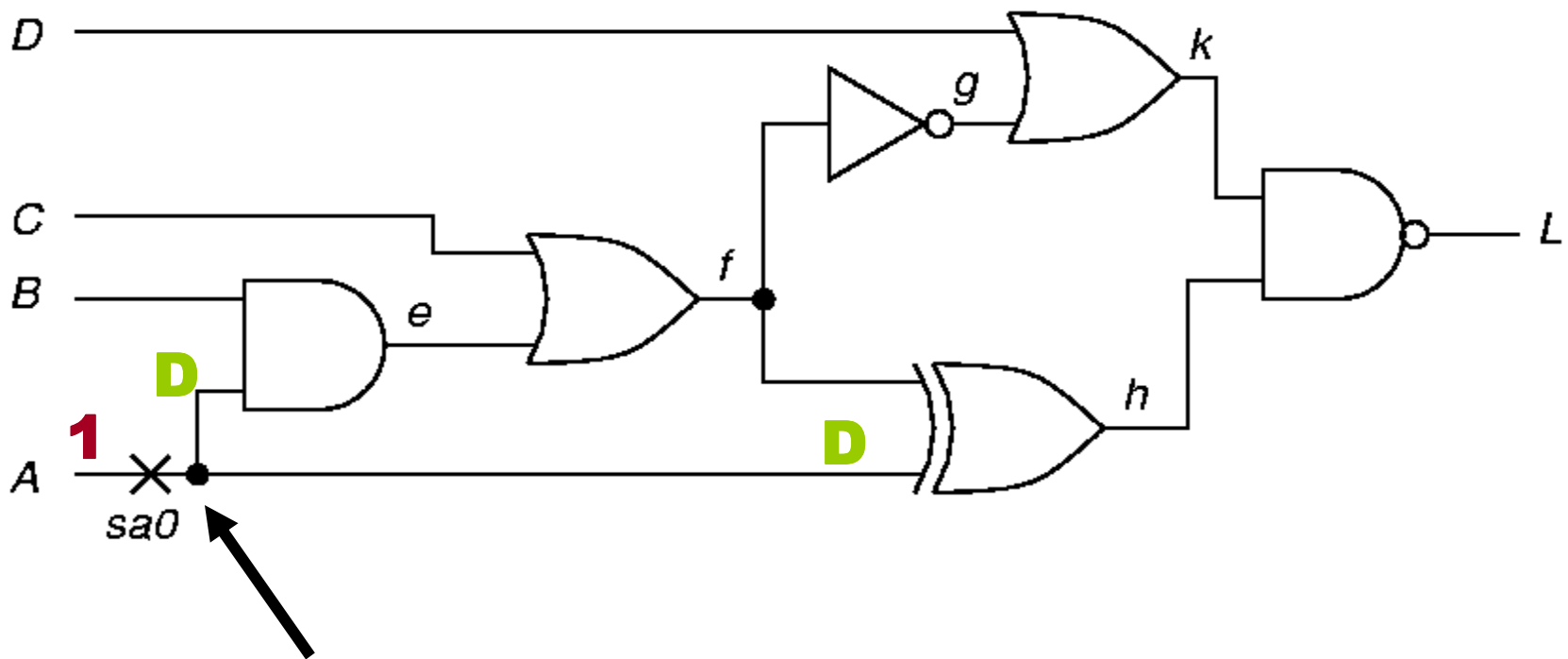
- Steps for Fault d sa0



Step	A	B	C	d	e	F	Cube type
1	1	1		D			PDF of AND gate
2				D	0	D*	Prop. D-cube for NOR
3		1	1		0		Singular Cover of NAND

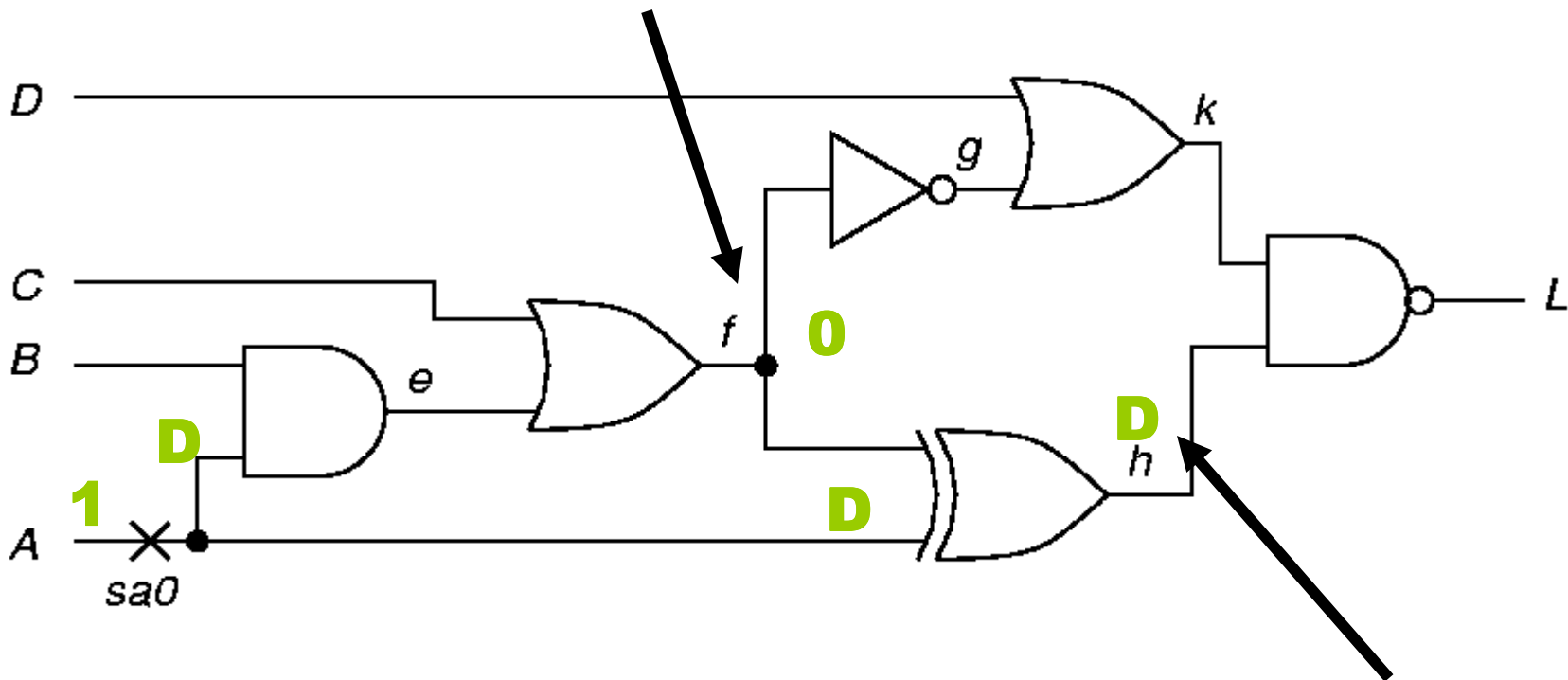
D-Algorithm..... Example2(1)

- Example 7.2 Fault A sa0
- **Step 1:** Select a primitive D-cube of the fault - Set **A = 1**



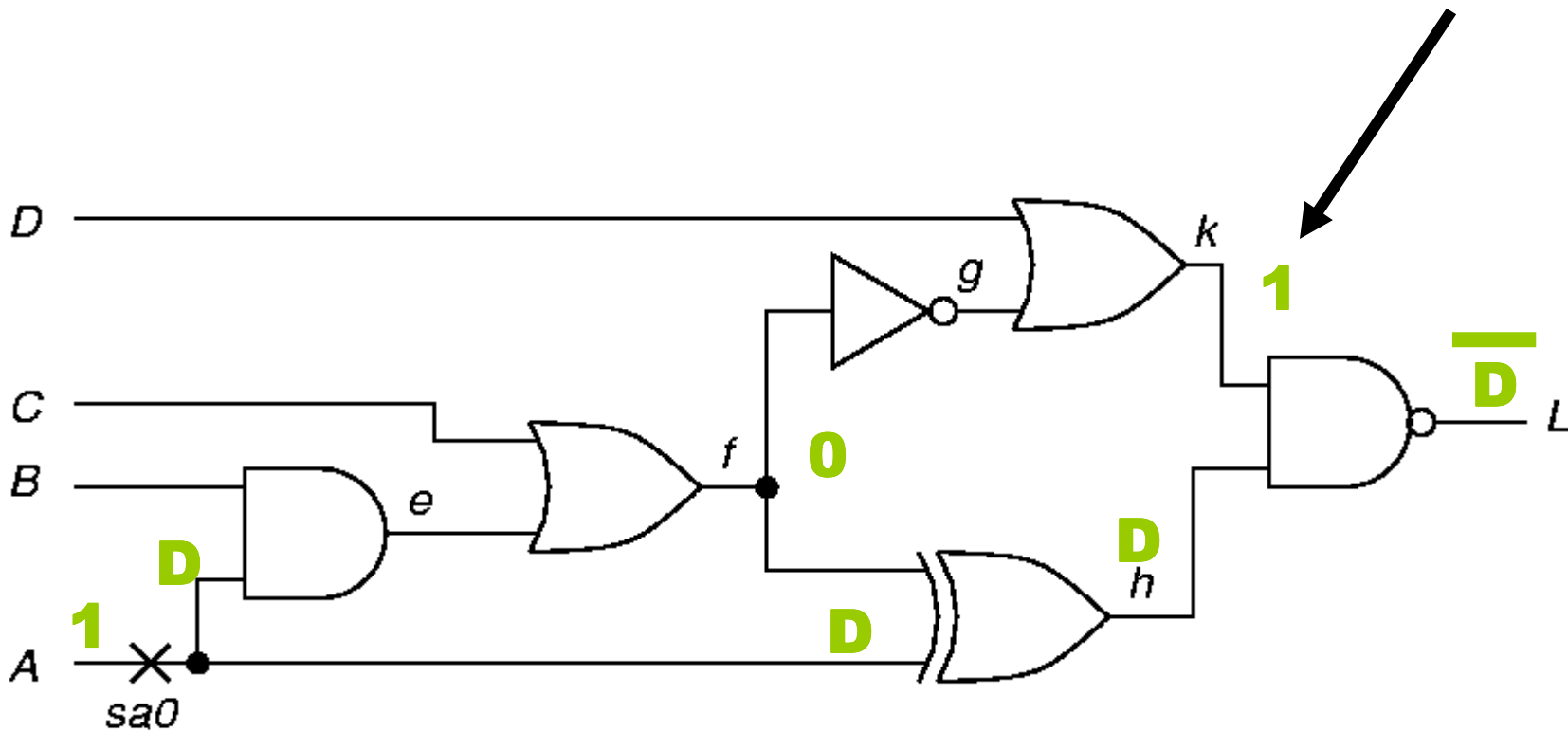
D-Algorithm..... Example2(2)

□ Step 2: *D-Drive* – Set $f = 0$



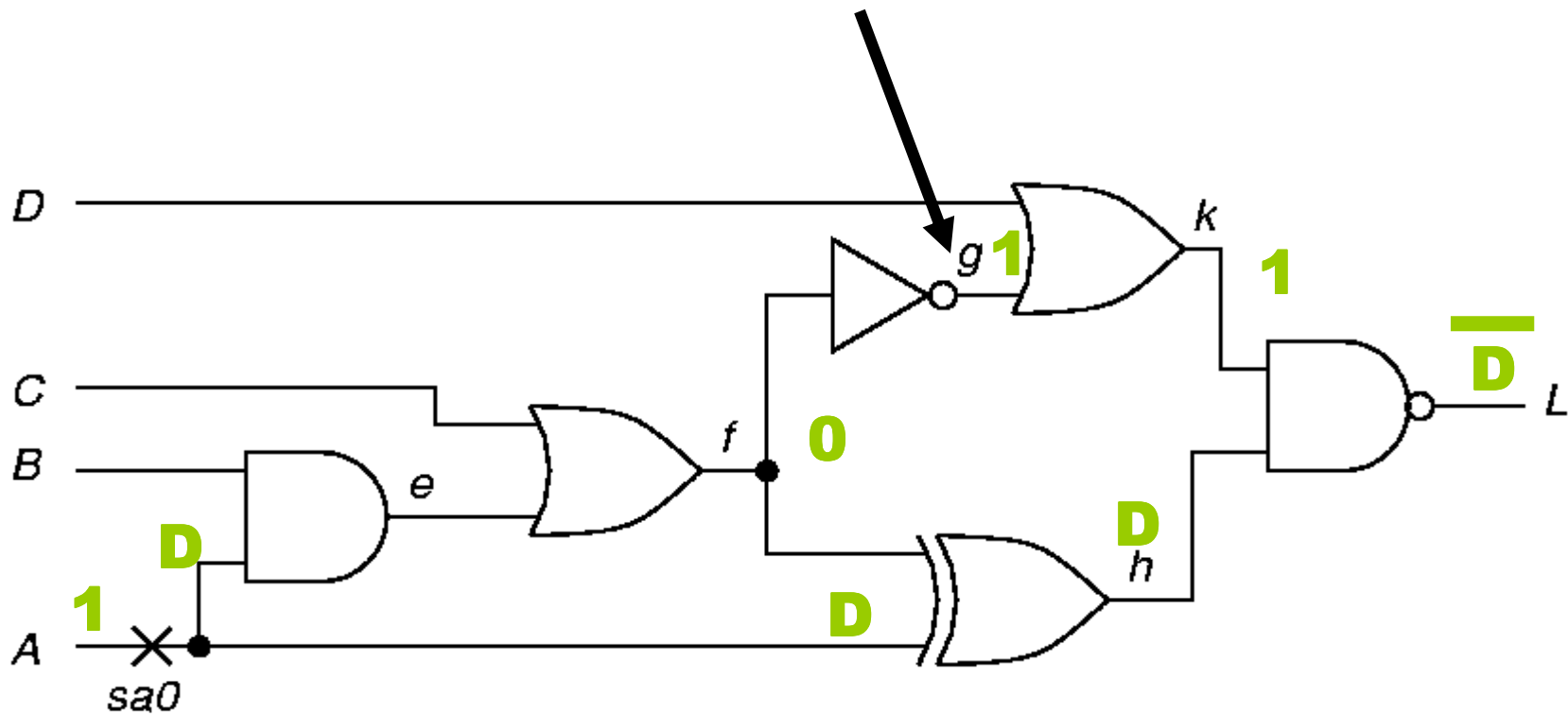
D-Algorithm..... Example2(3)

□ Step 3: *D-Drive* – Set $k = 1$



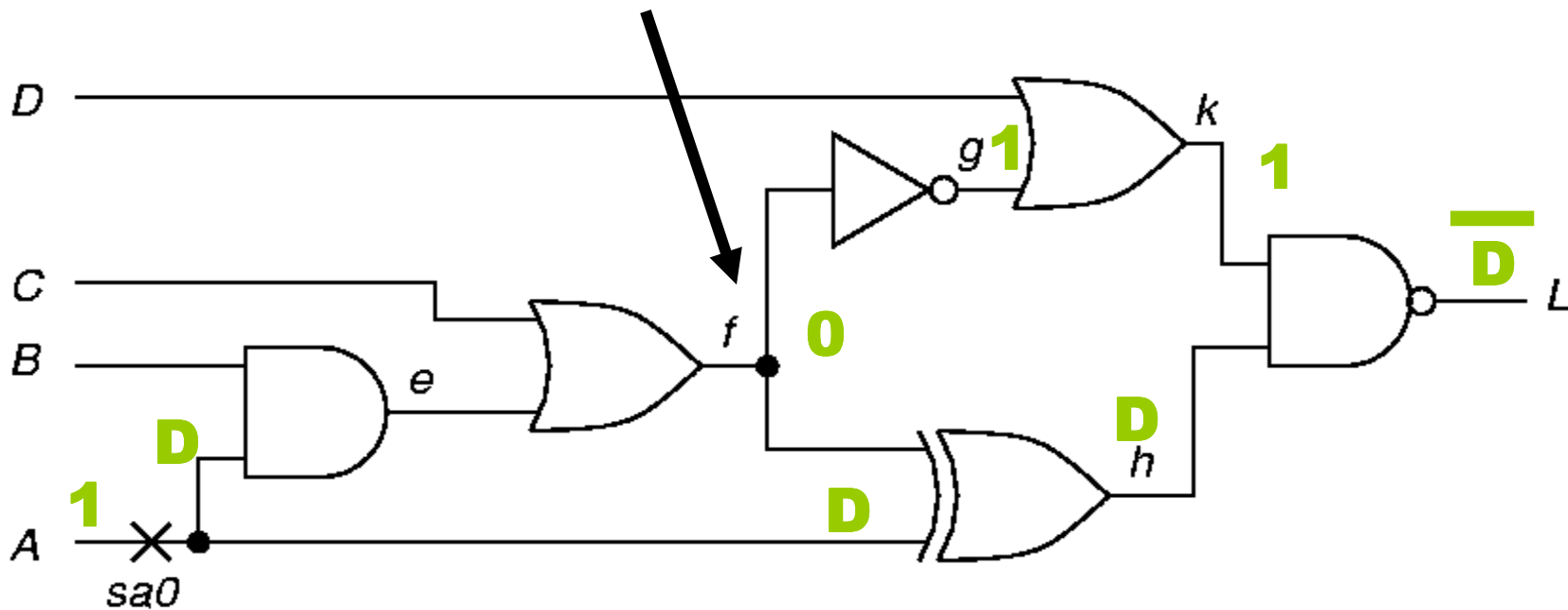
D-Algorithm..... Example2(4)

□ Step 4: *Consistency* – Set $g = 1$



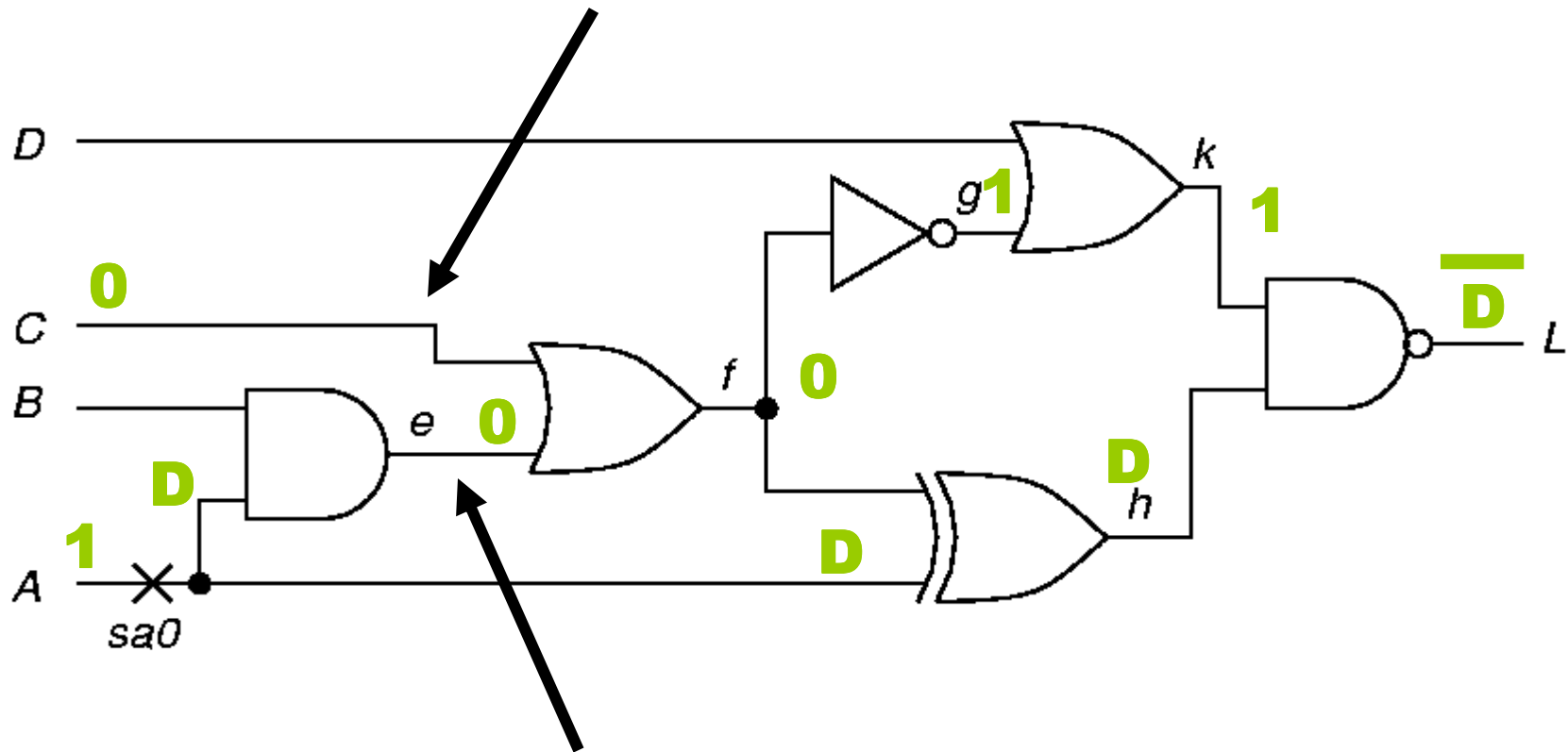
D-Algorithm..... Example2(5)

□ **Step 5: *Consistency* – $f = 0$ Already set**



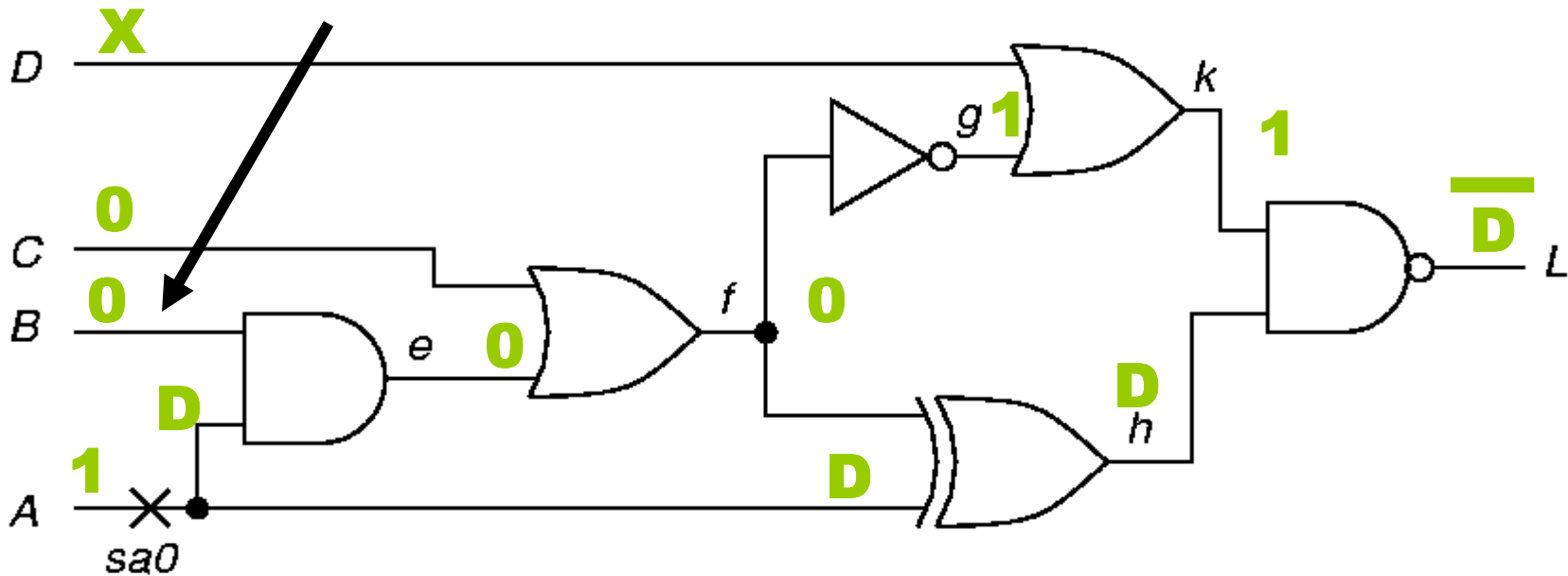
D-Algorithm..... Example2(6)

- **Step 6: Consistency** – Set $c=0$, Set $e=0$



D-Algorithm..... Example2(7)

- **Step 7: Consistency** – Set **B = 0**
- **No more line to justify: pattern found**



- **Test cube:** $A, B, C, D, e, f, g, h, k, L$
- $TC^7 = D00X001D1D^*$

PODEM.....Motivation

- IBM introduced semiconductor DRAM memory into its mainframes – late 1970's

- Memory had error correction and translation circuits – improved reliability
 - D-ALG unable to test these circuits
 - Search too undirected
 - Large XOR-gate trees
 - Must set all external inputs to define output
 - Needed a better ATPG tool

PODEM.....Goel IBM (1981)

□ **New concepts introduced:**

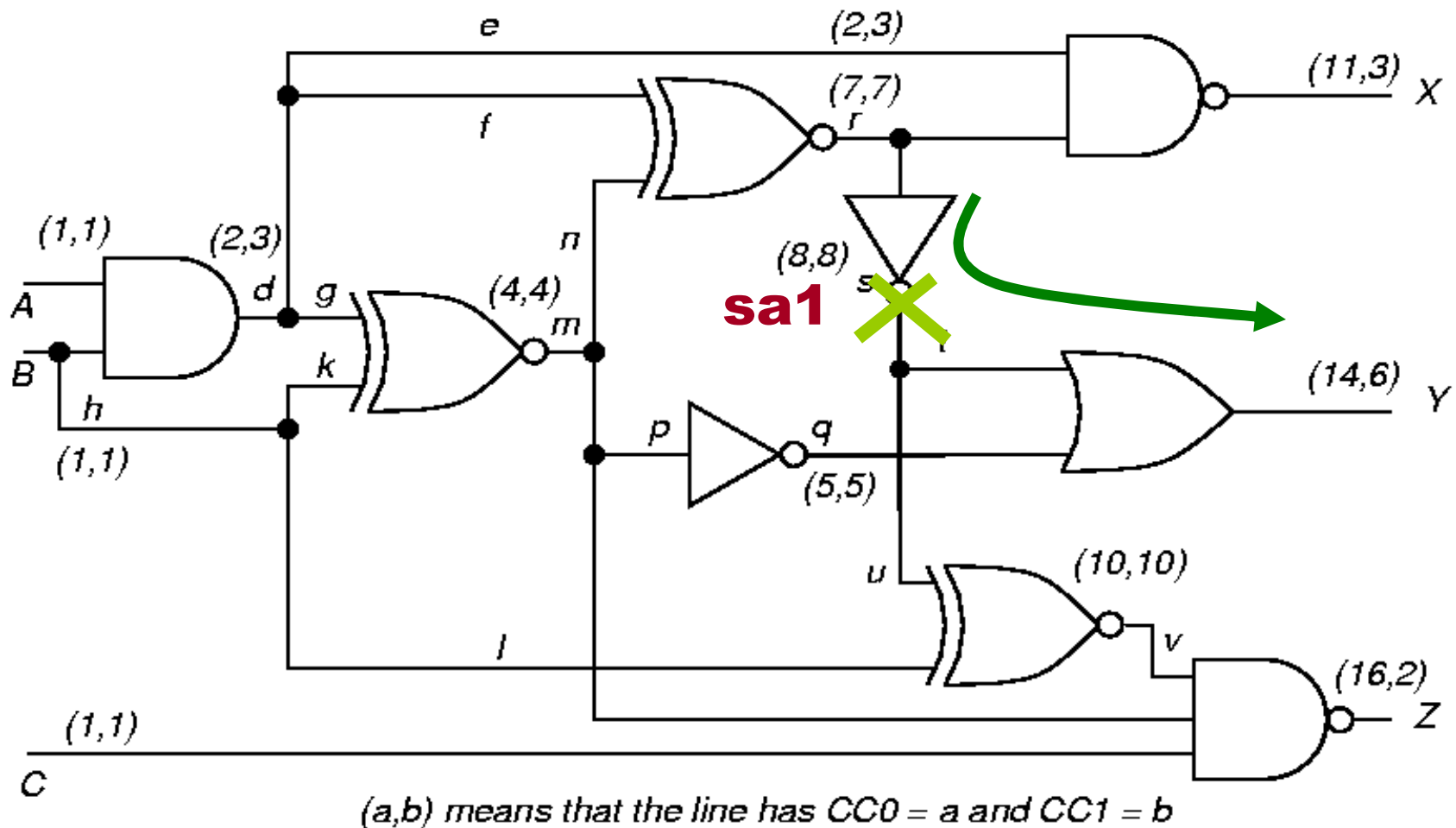
- Expand binary decision tree only around primary inputs
 - Reduction from 2^n to $2^{\#PI}$ (n: # of logic gates)
- Use *X-PATH-CHECK* to test whether *D-frontier* still there
- *Objectives*: bring ATPG closer to propagating D (D^*) to PO
- *Backtracing*
 - Use # of levels or CC measures to assign PI objectives

PODEM..... High-Level Flow

1. Assign binary value to unassigned PI
2. Determine implications of all PIs
3. Test Generated? If so, **done**.
4. Test possible with additional assigned PIs? If maybe, go to **Step 1**
5. Is there untried combination of values on assigned PIs? If not, **exit: untestable fault**
6. Set untried combination of values on assigned PIs using objectives and backtrace. Then, go to **Step 2**

PODEM..... Example(1)

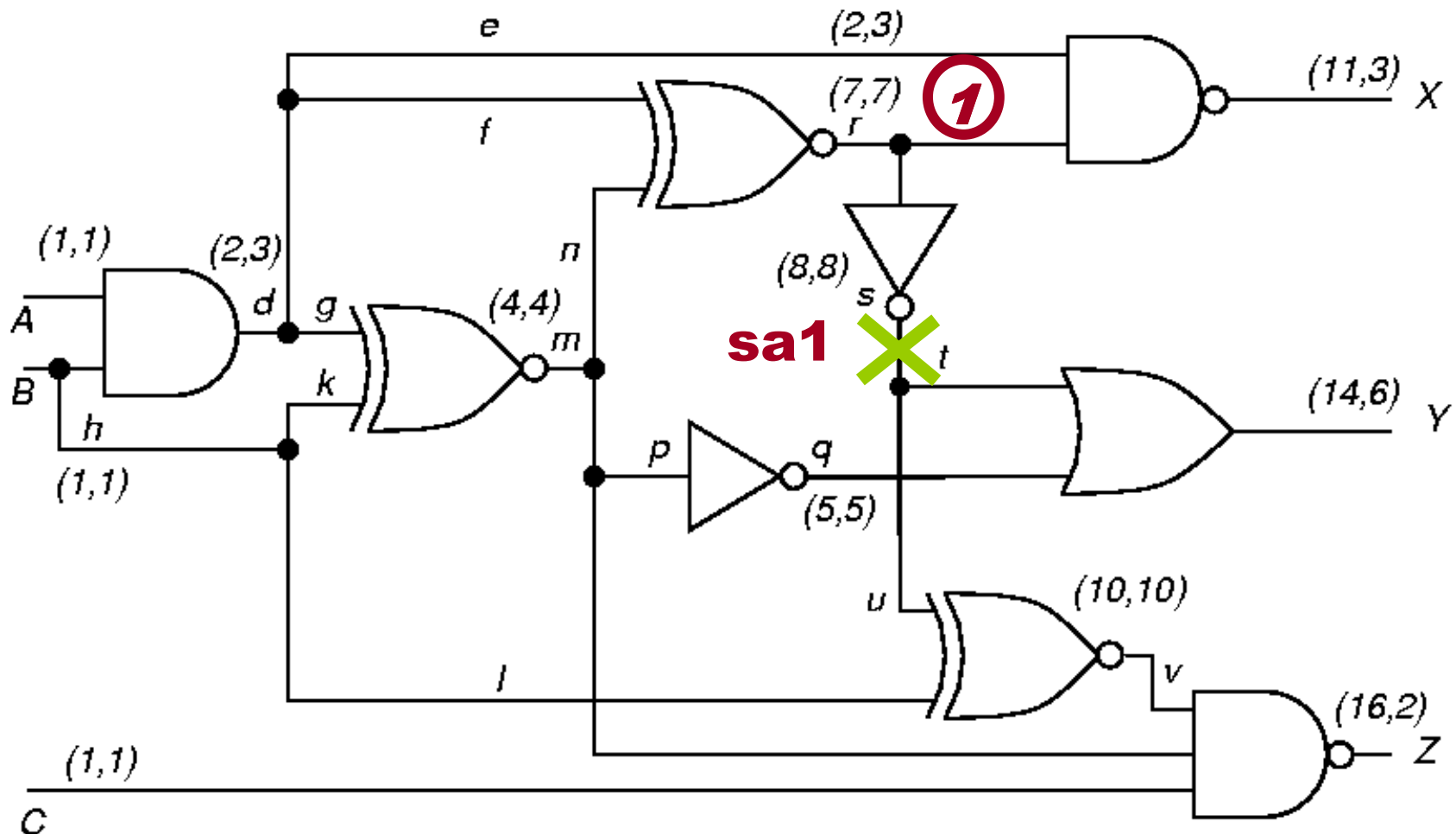
- Example 7.3
- Step 1: Select path **s - Y** for fault propagation; not s-u-v-Z (level distance from PO is 1 versus 2)



PODEM..... Example(2)

Step 2:

- Initial objective: Set r to **1** to sensitize fault

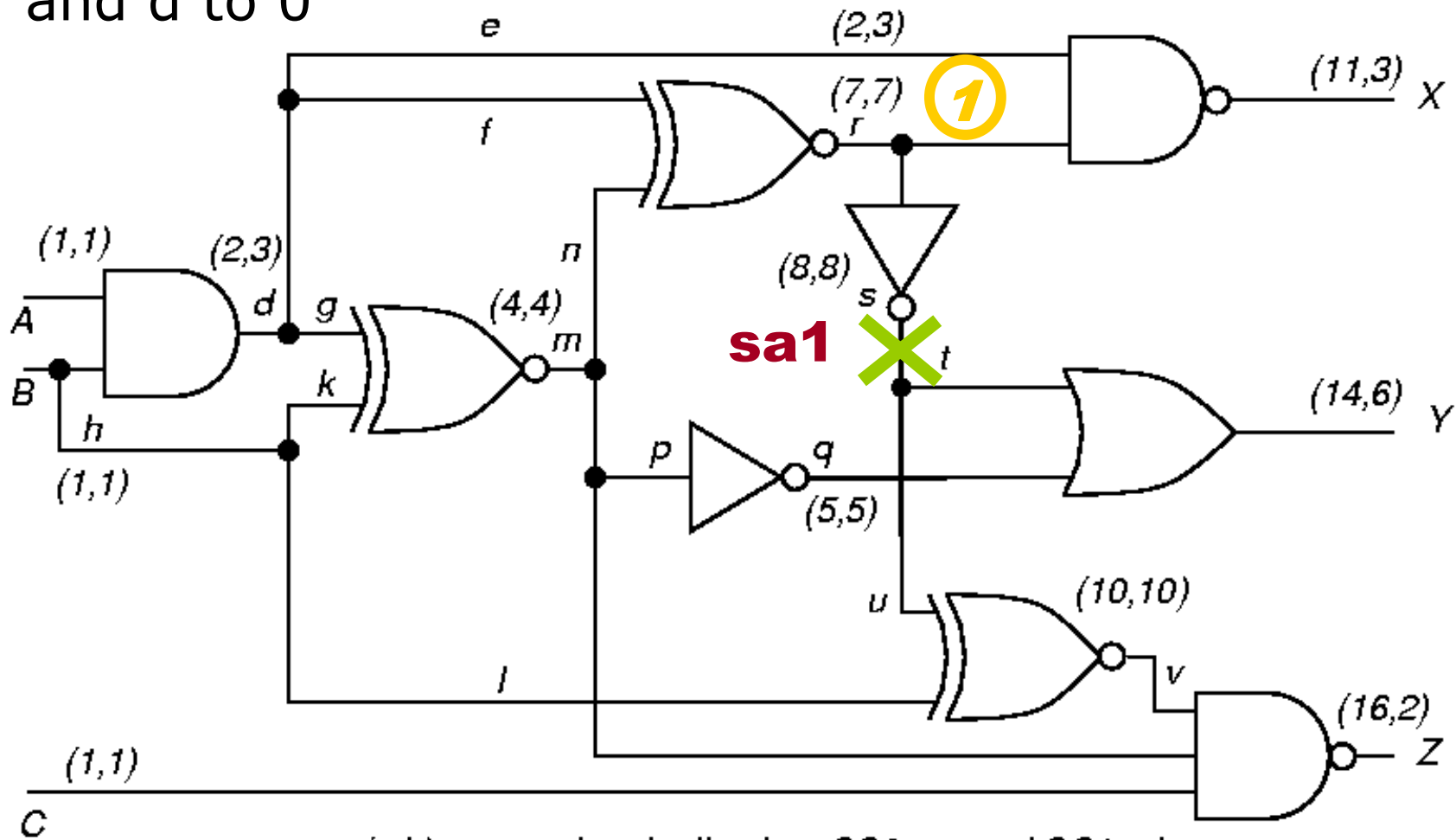


(a,b) means that the line has $CC0 = a$ and $CC1 = b$

PODEM..... Example(3)

Step 3: Backtrace from r

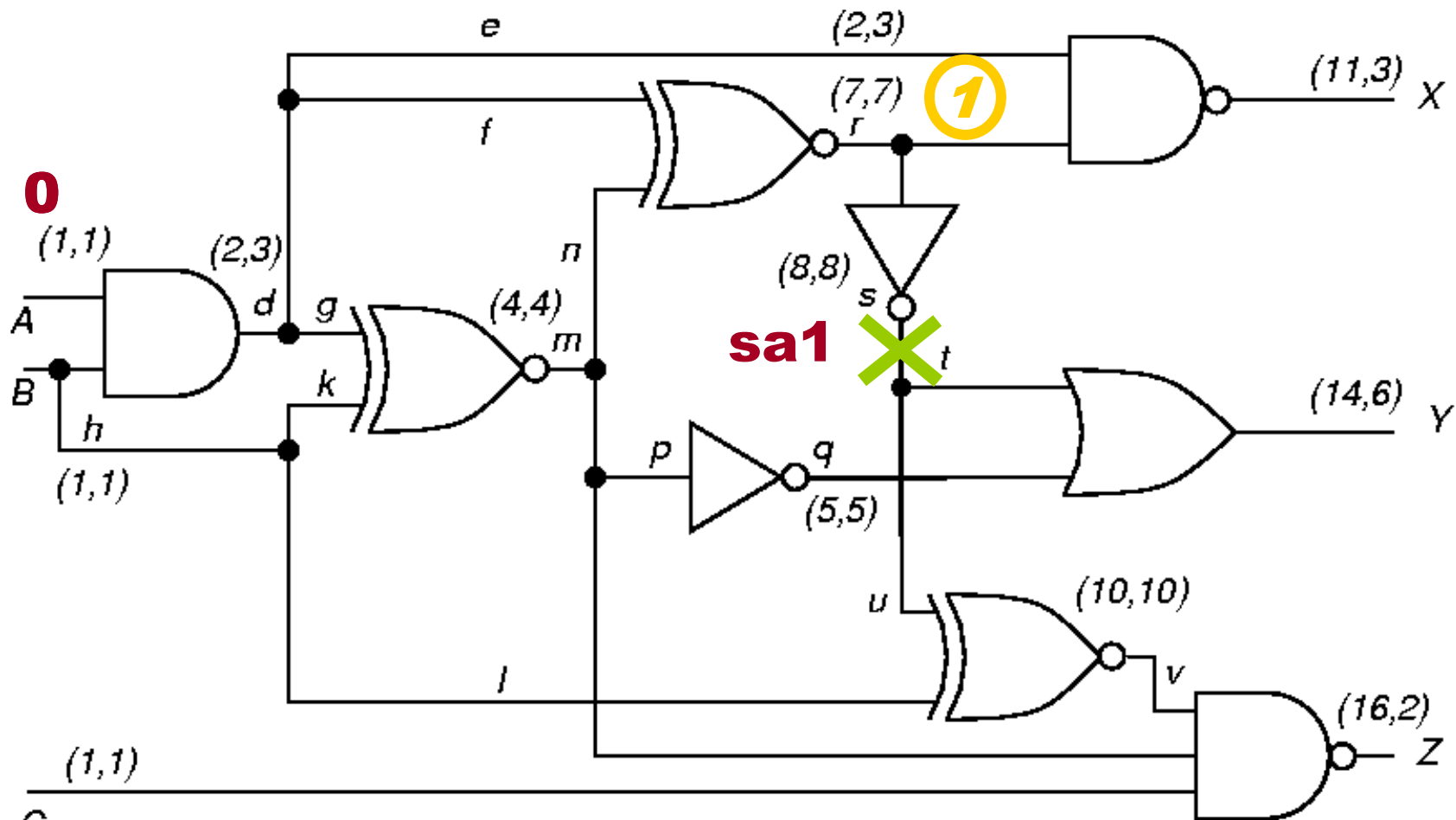
=> Intermediate objectives: Set n to 0, m to 0, g to 0 and d to 0



(a,b) means that the line has $CC0 = a$ and $CC1 = b$

PODEM..... Example(4)

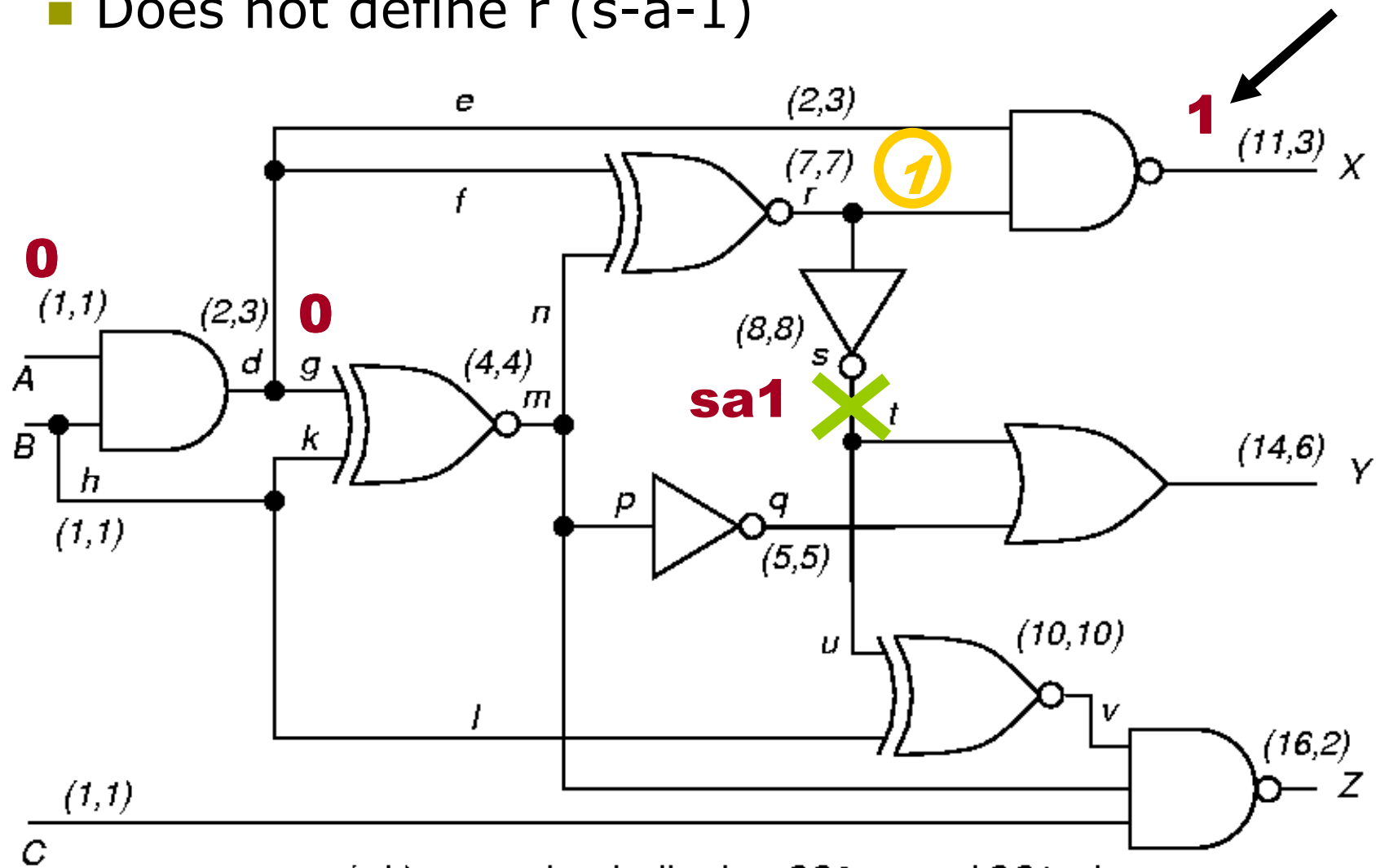
- Step 4: Set $A = 0$ in implication stack



(a,b) means that the line has $CC0 = a$ and $CC1 = b$

PODEM..... Example(5)

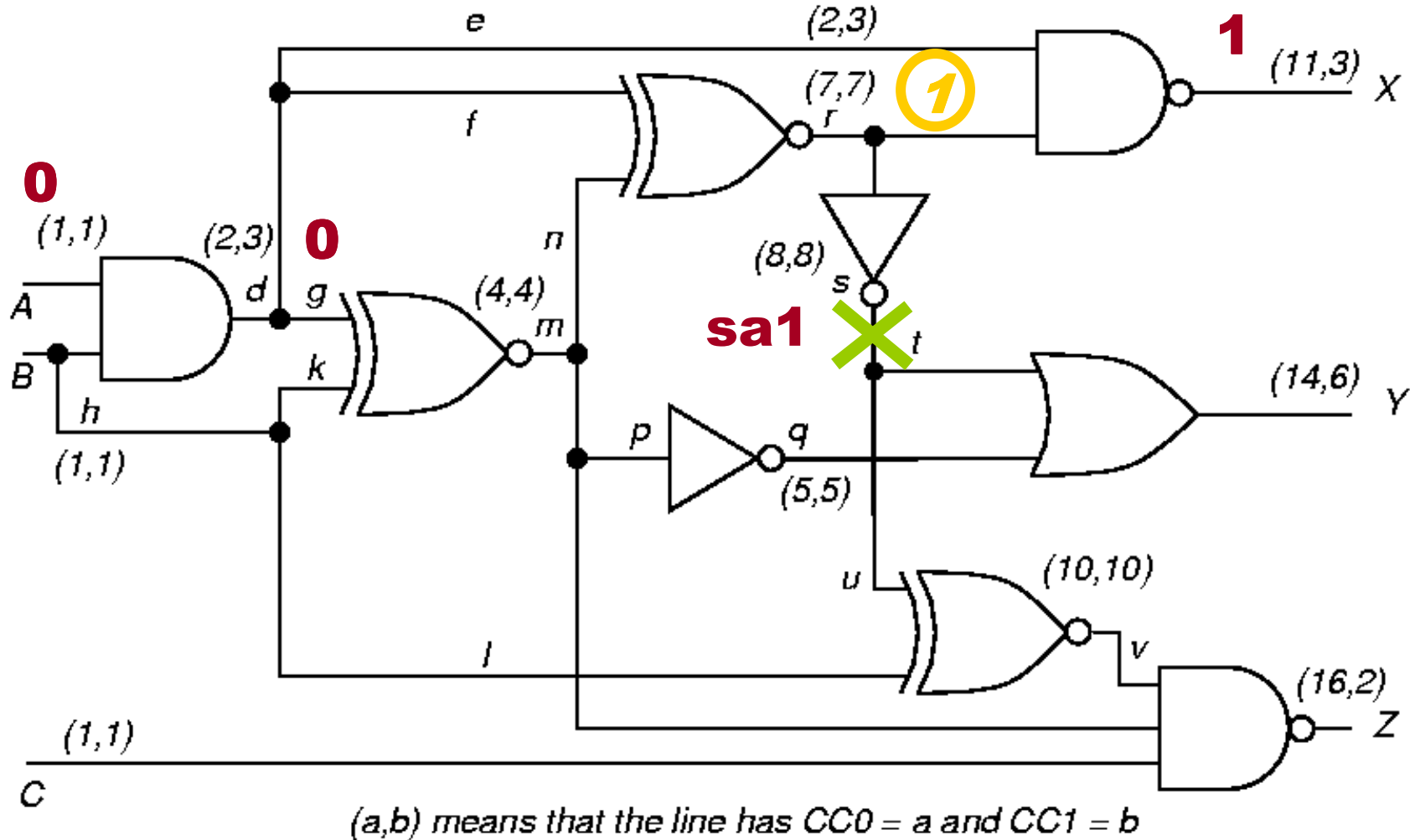
- **Step 5: Forward implications: $d = 0, X = 1$**
 - Does not define r ($s-a-1$)



(a,b) means that the line has $CC0 = a$ and $CC1 = b$

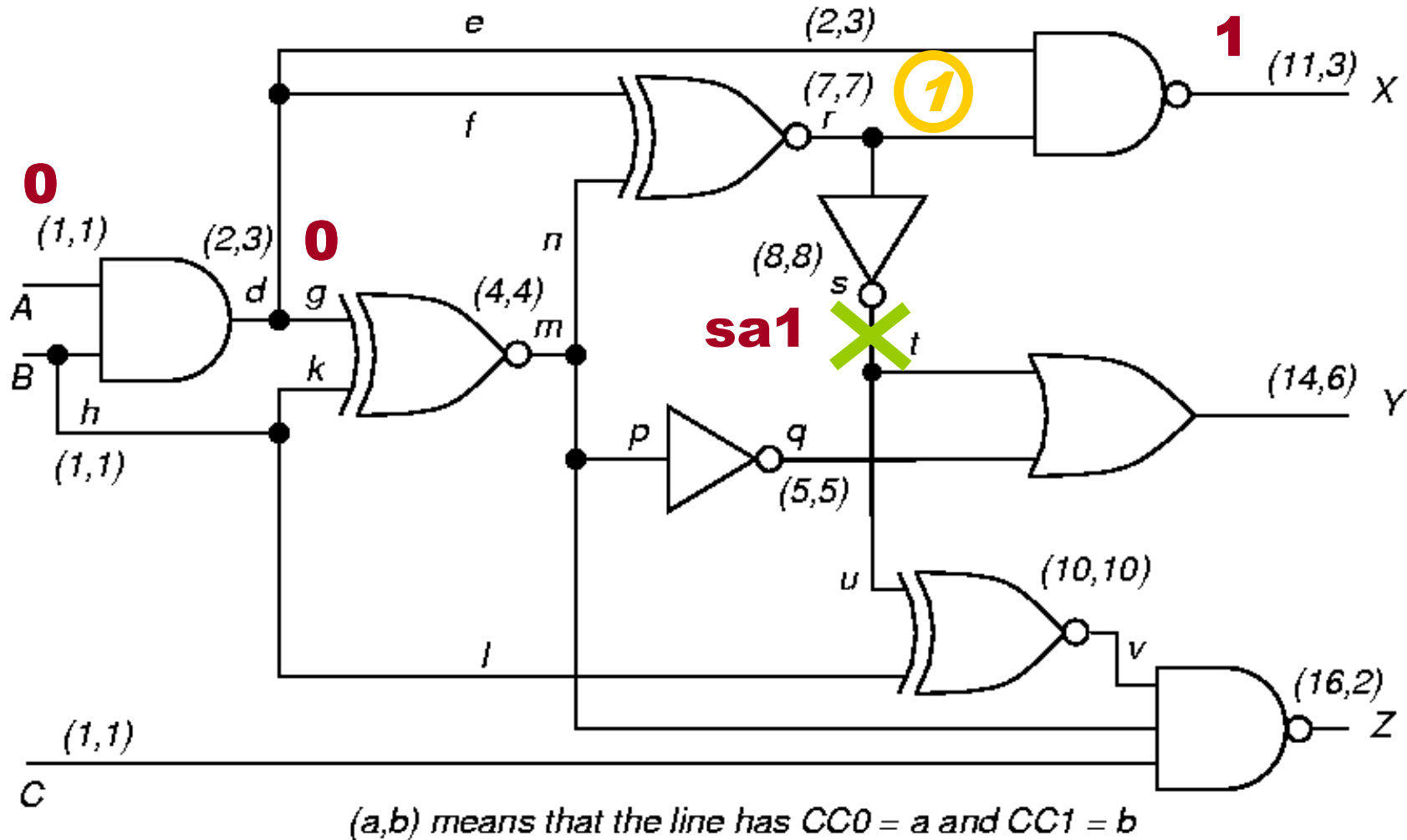
PODEM..... Example(6)

Step 6: Initial objective: set r to 1



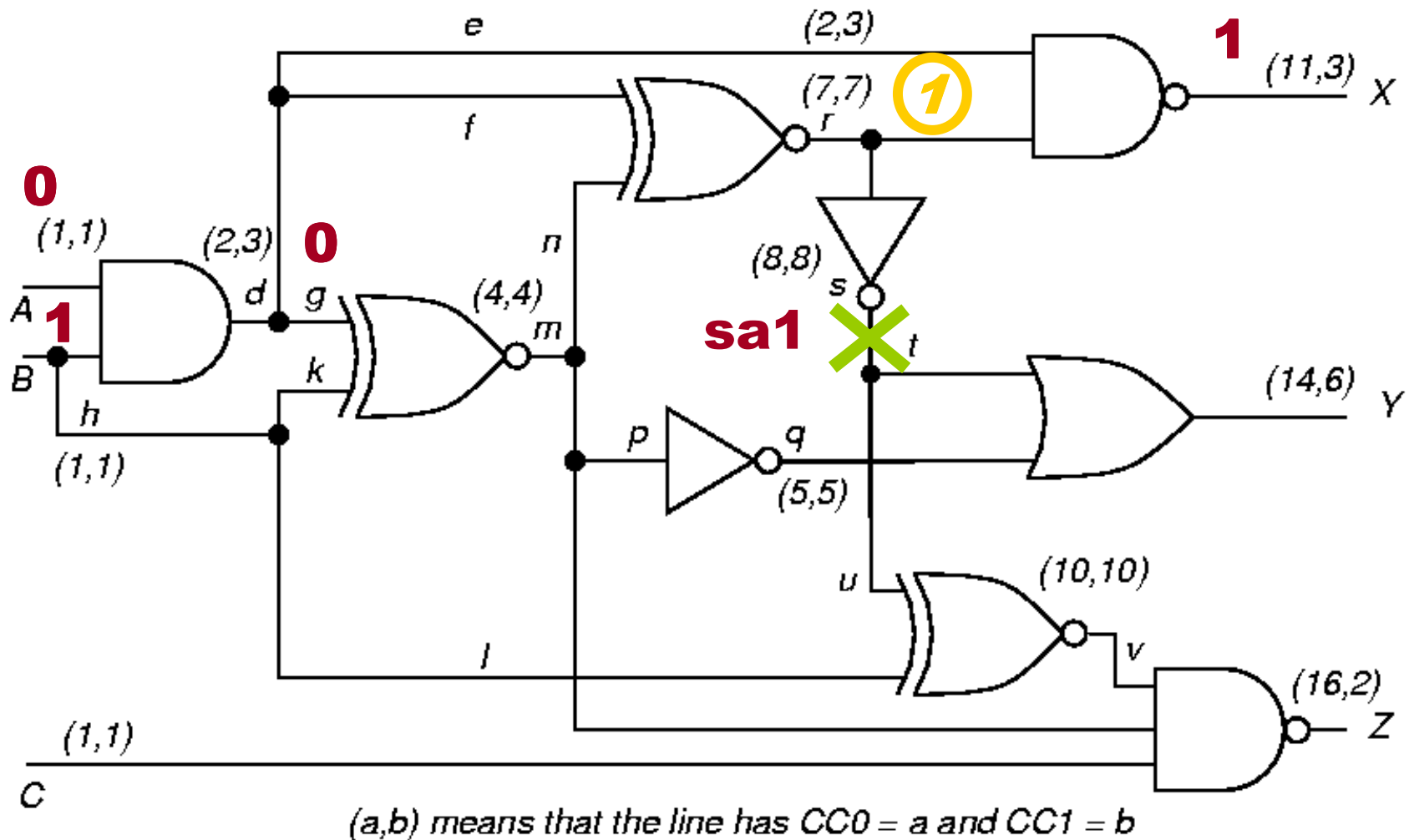
PODEM..... Example(7)

□ Step 7: Backtrace from *r* again



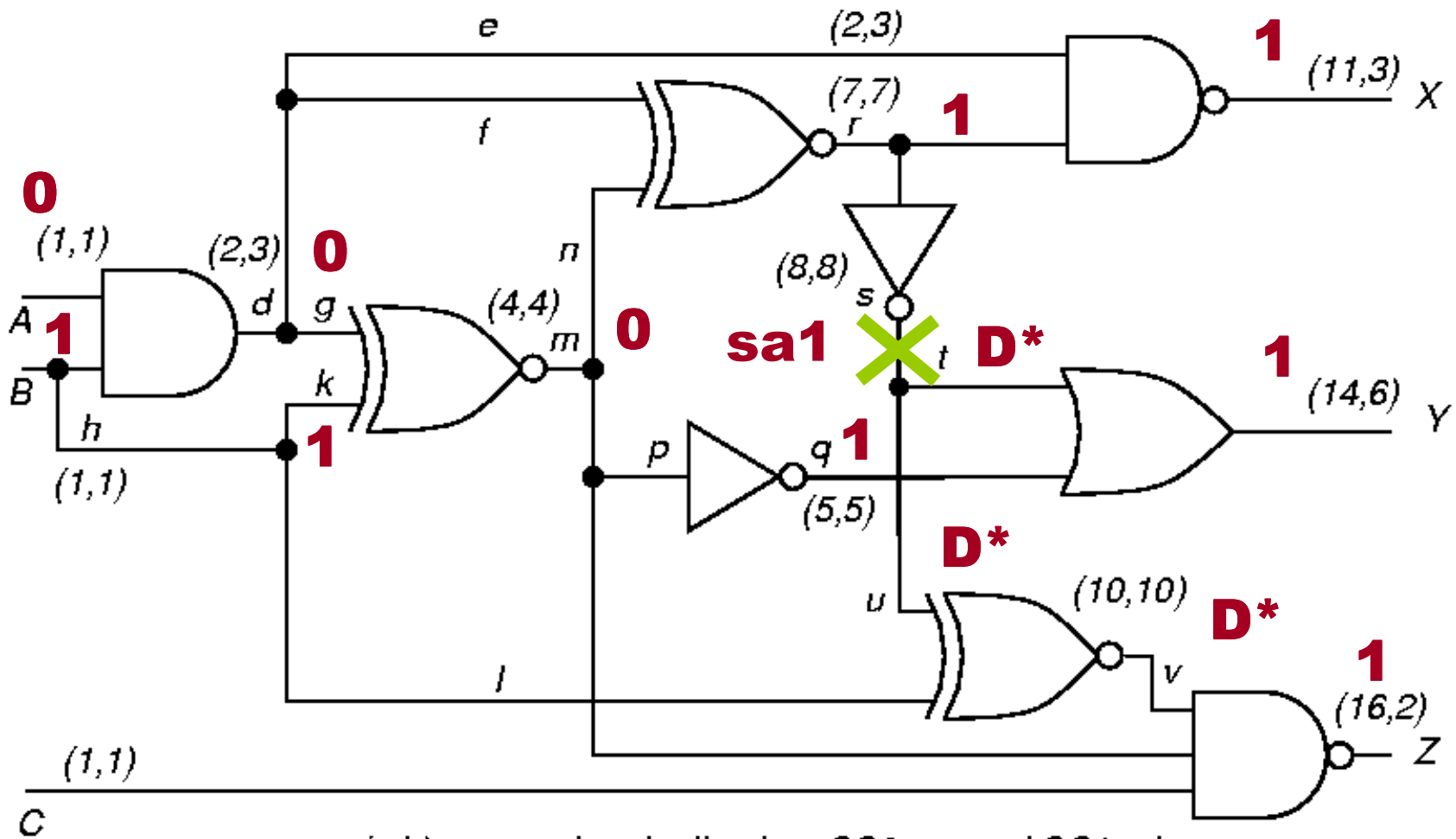
PODEM..... Example(8)

- Step 8: Set **B** to 1.
- Implications in stack: **A = 0, B = 1**



PODEM..... Example(9)

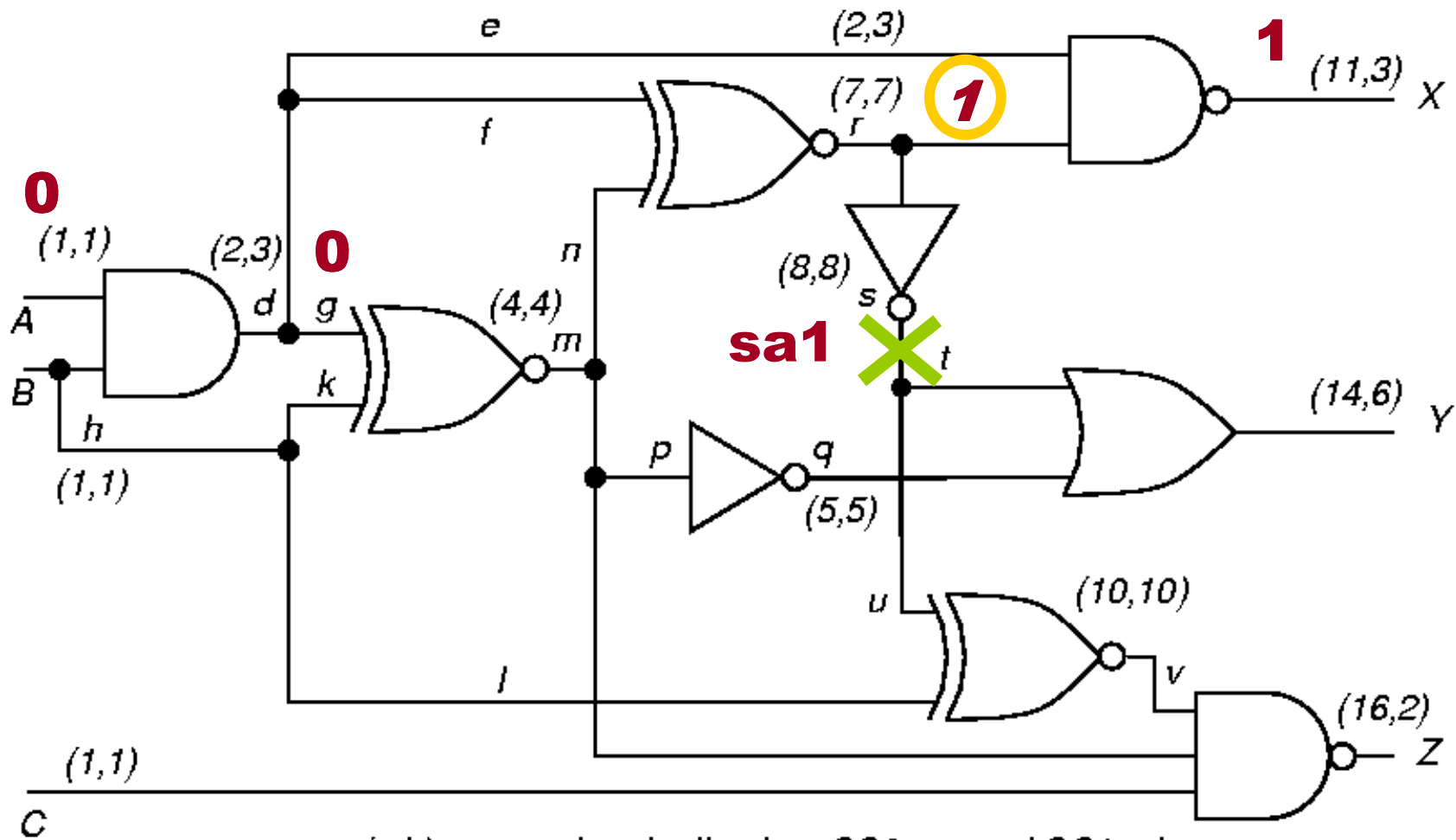
- Step 9: Forward implications: $k=1, m=0, r=1, q=1, Y=1, s=D^*, u=D^*, v^*=D, Z=1$



(a,b) means that the line has CC0 = a and CC1 = b

PODEM..... Example(10)

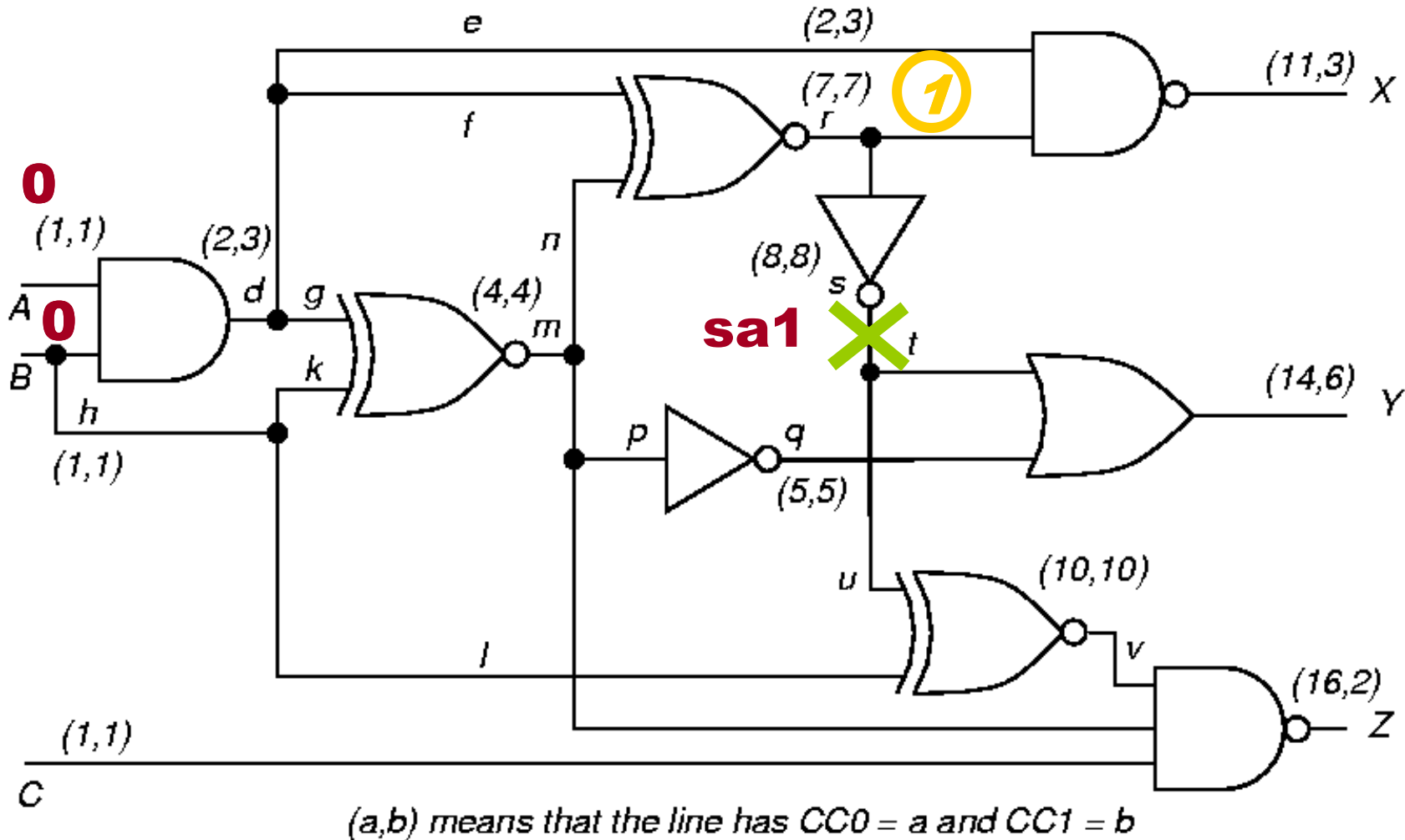
- Step 10: *X-PATH-CHECK* shows paths *s-Y* and *s-u-v-Z* blocked (*D-frontier* disappeared)



(a,b) means that the line has $CC0 = a$ and $CC1 = b$

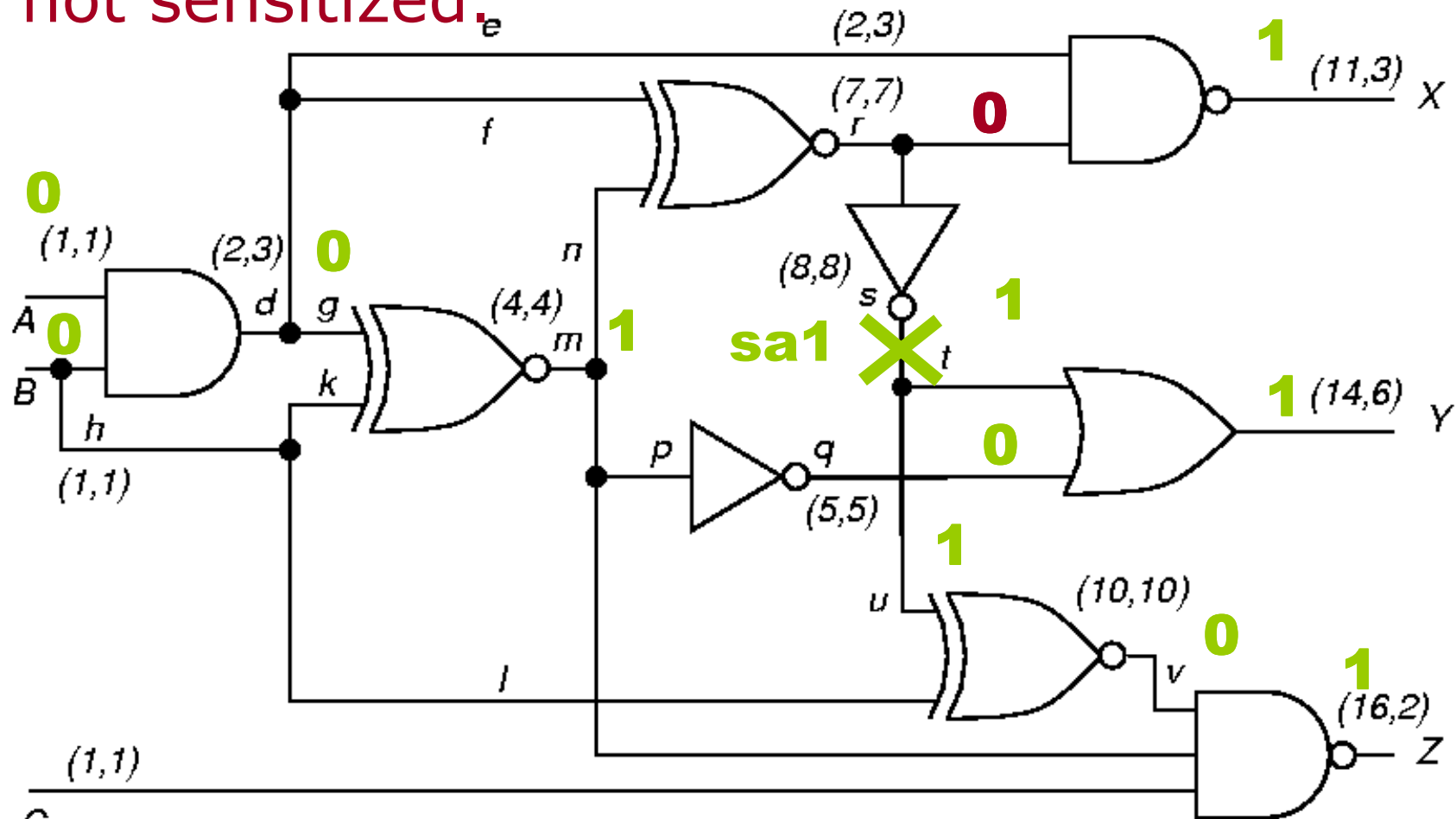
PODEM..... Example(11)

□ **Step11:** Set $B = 0$ (alternate assignment)



PODEM..... Example(12)

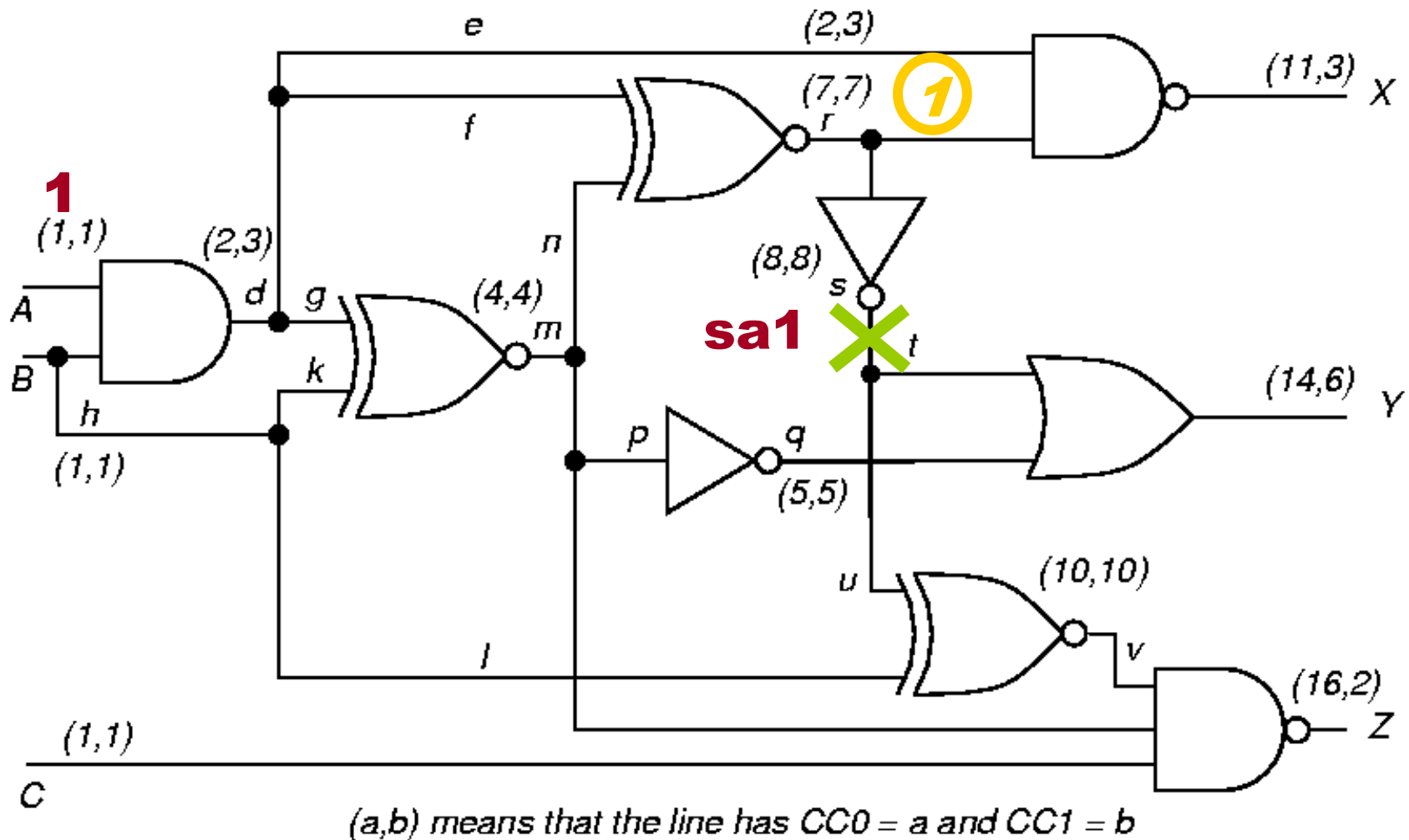
- Step 12: Forward implications: $d=0$, $X=1$, $m=1$, $r=0$, $s=1$, $q=0$, $Y=1$, $v=0$, $Z=1$. Fault not sensitized.



(a,b) means that the line has $CC0 = a$ and $CC1 = b$

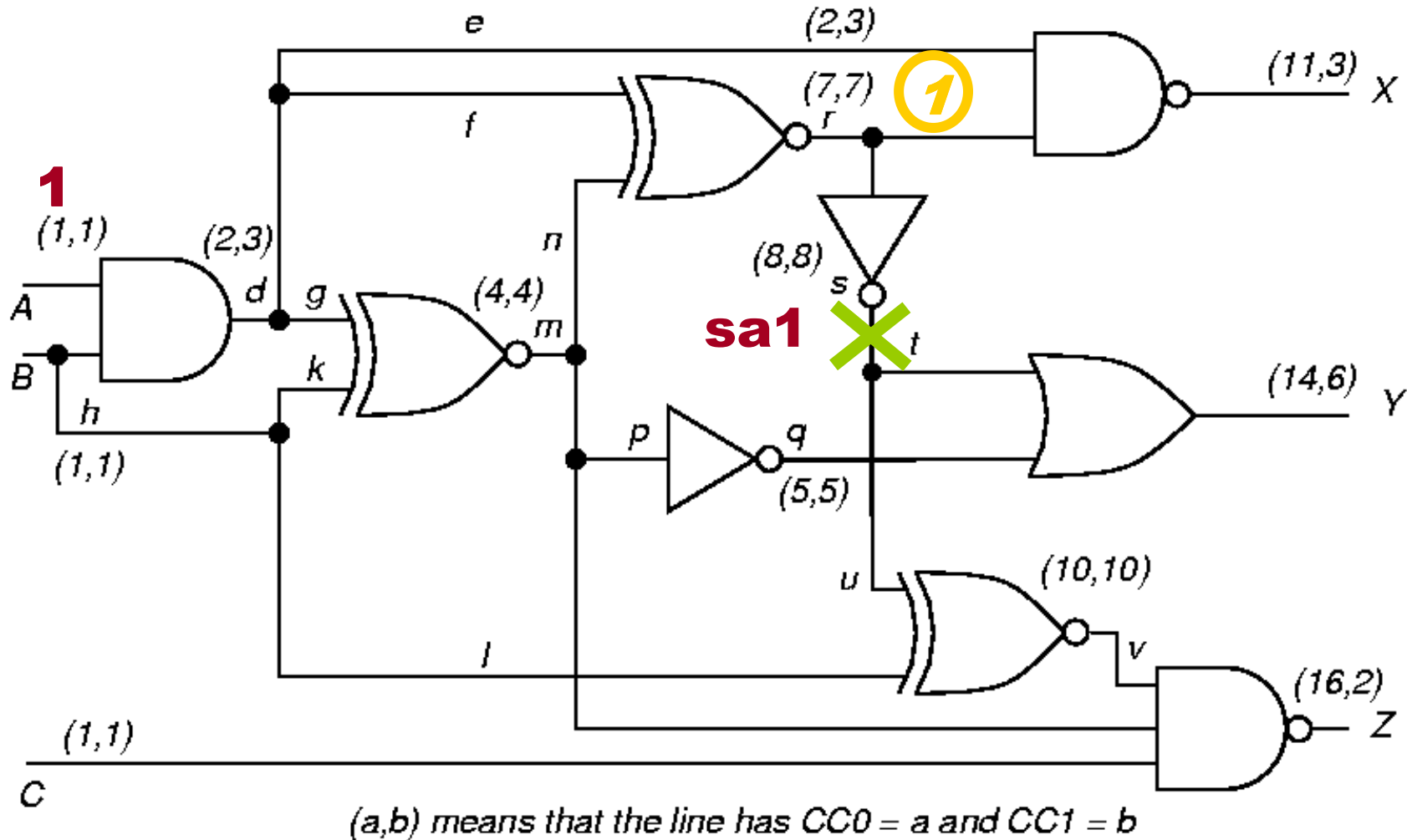
PODEM..... Example(13)

- Step 13: Set **A = 1** (alternate assignment)
- No implications possible -> Backtrace r



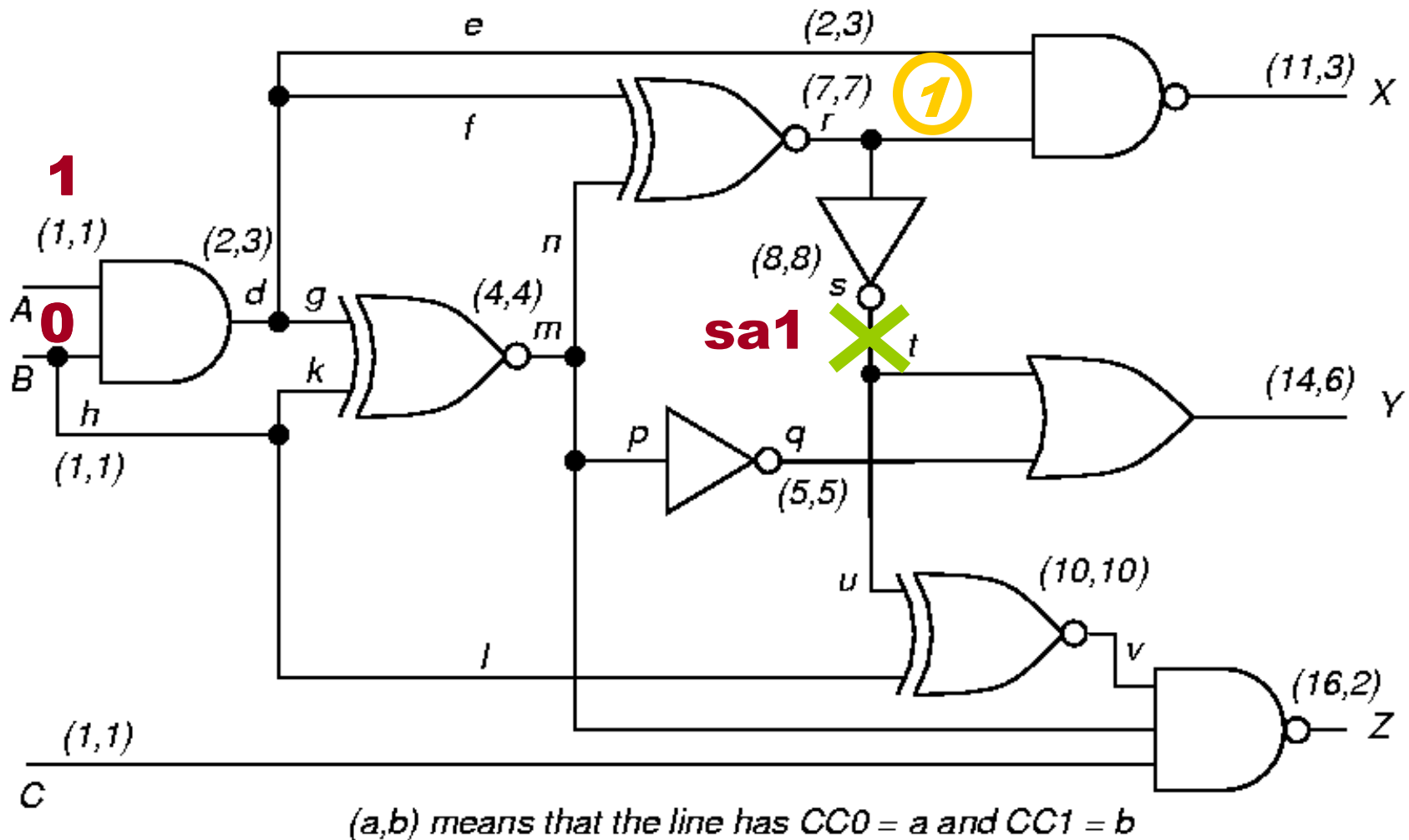
PODEM..... Example(14)

Step 14: Backtrace from *r* again



PODEM..... Example(15)

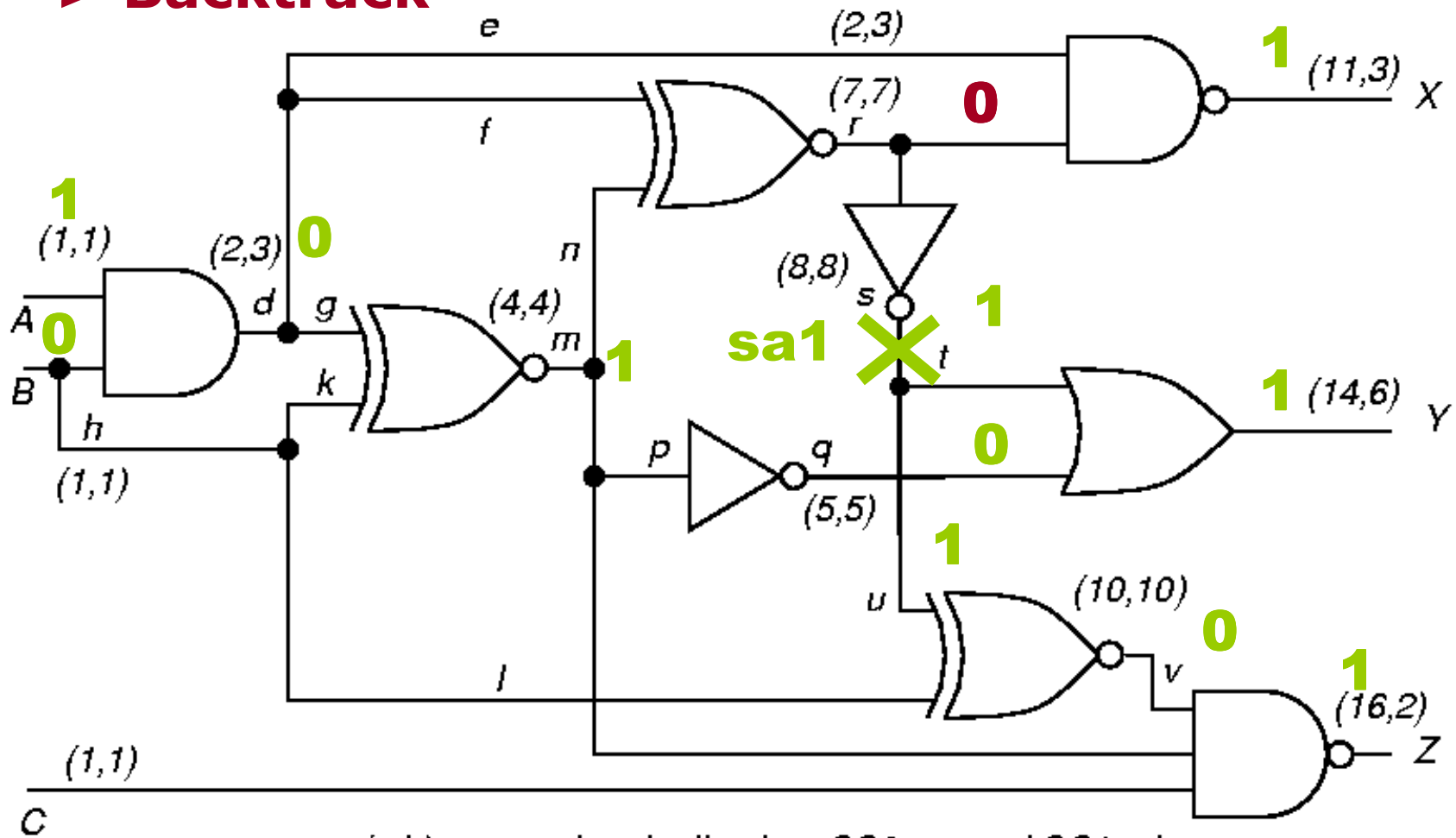
- Step 15: Set $B = 0$. Implications in stack: $A = 1, B = 0$



PODEM..... Example(16)

□ **Step 16:** Forward implications: $d=0, X=1, m=1, r=0$. **Conflict: fault not sensitized.**

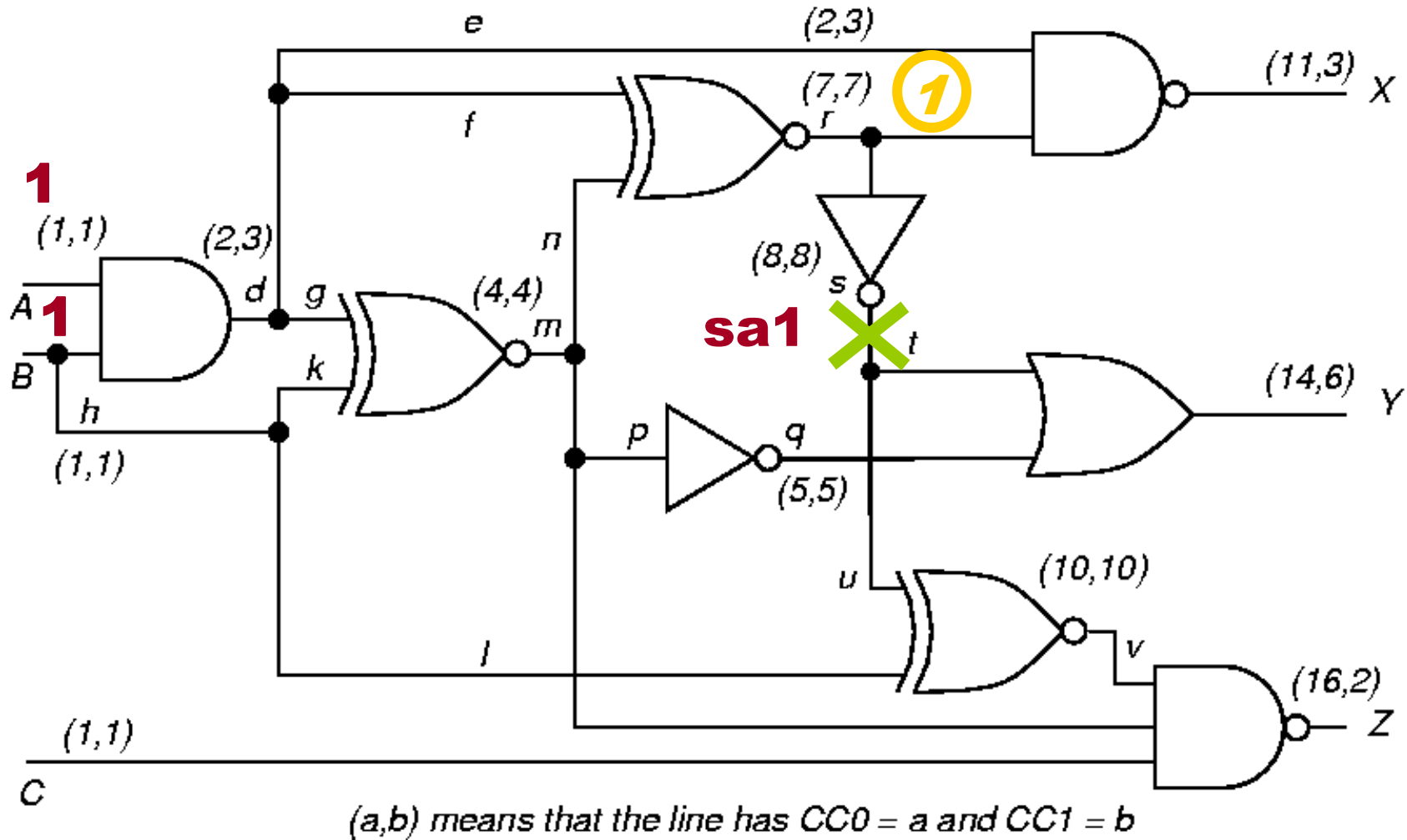
-> **Backtrack**



(a,b) means that the line has $CC0 = a$ and $CC1 = b$

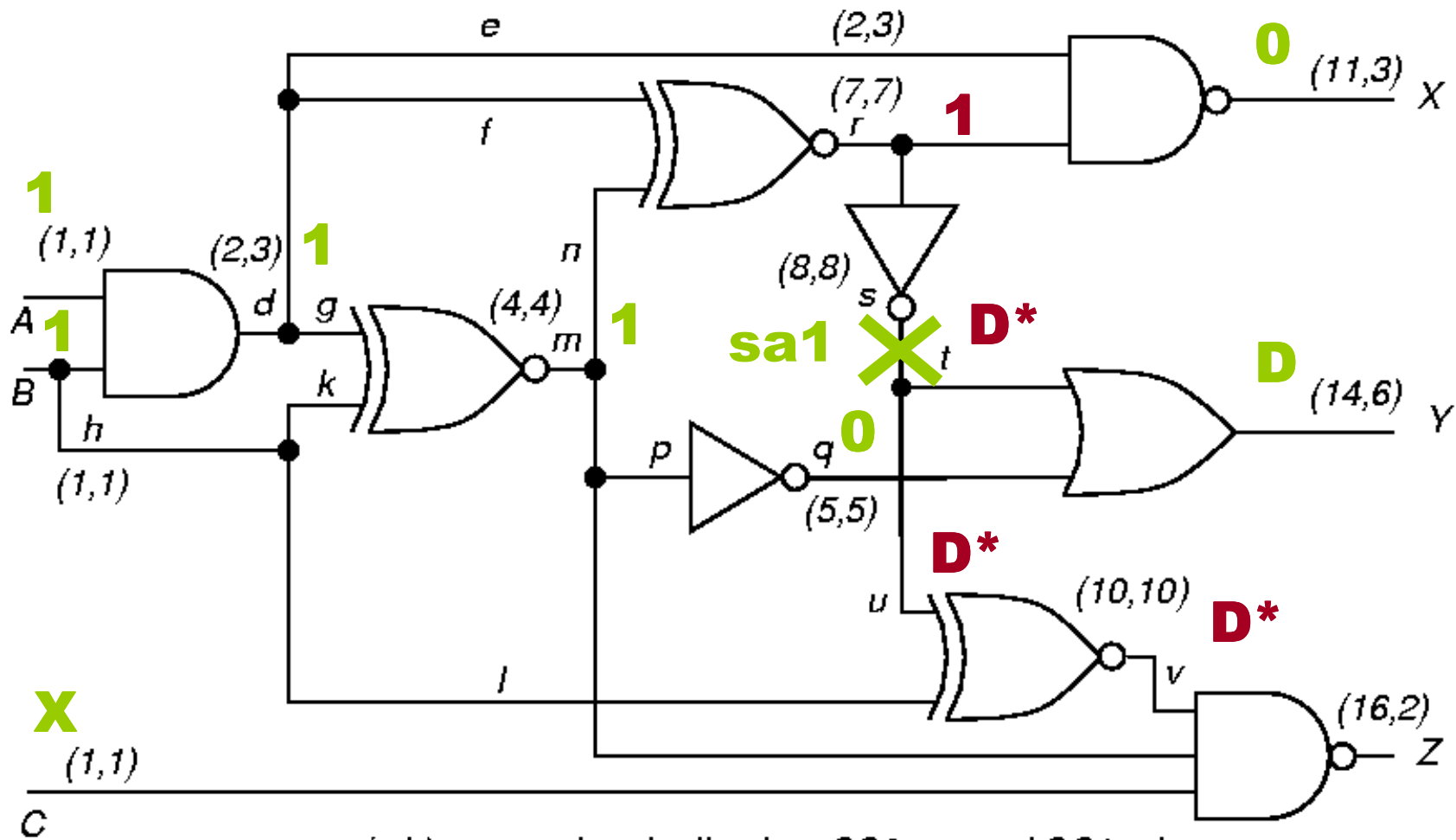
PODEM..... Example(17)

- Step 17: Set **B = 1** (alternate assignment)



PODEM..... Example(18)

- Step 18: Forward implications: $d=1$, $m=1$, $r=1$, $q=0$, $s=D^*$, $v=D^*$, $X=0$, $Y=D^*$



(a,b) means that the line has $CC0 = a$ and $CC1 = b$

PODEMAlgorithm

```
while (no fault effect at POs)
  if (xpathcheck (D-frontier))
    (l, vl) = Objective (fault, vfault);
    (pi, vpi) = Backtrace (l, vl);
    Imply (pi, vpi);
    if (PODEM (fault, vfault) == SUCCESS) return (SUCCESS);
    (pi, vpi) = Backtrack ();
    Imply (pi, vpi);
    if (PODEM (fault, vfault) == SUCCESS) return (SUCCESS);
    Imply (pi, "X");
    return (FAILURE);
  else if (implication stack exhausted)
    return (FAILURE);
  else Backtrack ();
return (SUCCESS);
```

Summary classical ATPG algorithms

- D-ALG – First complete ATPG algorithm
 - *D-Cube*
 - *D-Calculus*
 - *Implications* – forward and backward
 - *Implication stack*
 - *Backup*

- PODEM
 - Expand decision tree only around PIs
 - Use *X-PATH-CHECK* to see if *D-frontier* exists
 - *Objectives* -- bring ATPG closer to getting D (D^*) to PO
 - *Backtracing*

Advanced combinational ATPG

- FAN – *Multiple Backtrace* (1983)
- TOPS – *Dominators* (1987)
- SOCRATES – *Learning* (1988)
- Legal Assignments (1990)
- EST – *Search space learning* (1991)
- BDD Test generation (1991)
- *Implication Graphs* and *Transitive Closure* (1988 - 97)
- *Recursive Learning* (1995)

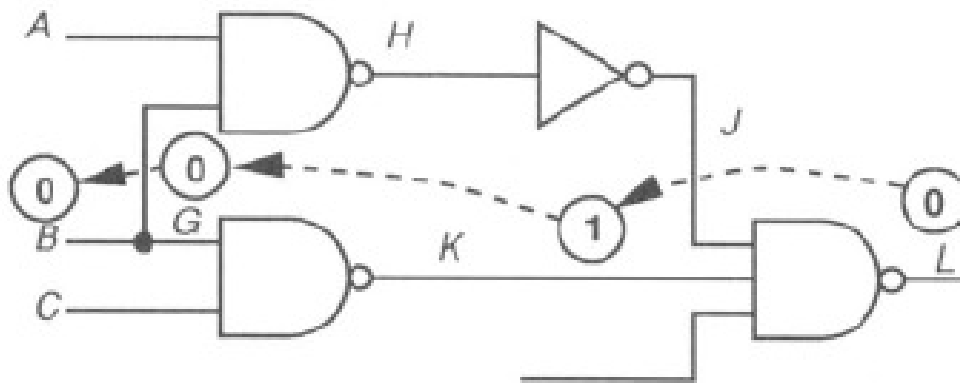
Advanced comb. ATPG.....FAN

□ FAN(1983)

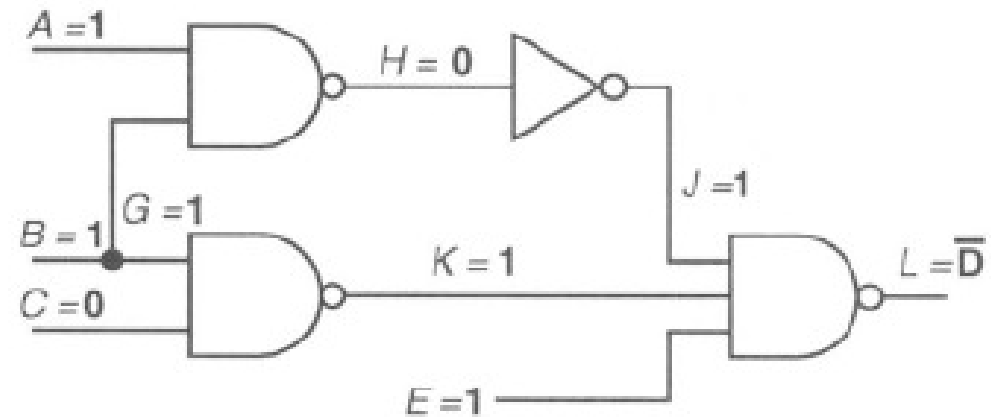
- Further limit the ATPG space and accelerate Backtracing
- Use **immediate implications** of uniquely-determined signals
- **Unique sensitization**
- Assign values to **headlines** (instead of PI)
- **Multiple Backtrace**

Advanced comb. ATPG.....FAN

- Use **immediate implications** of uniquely-determined signals



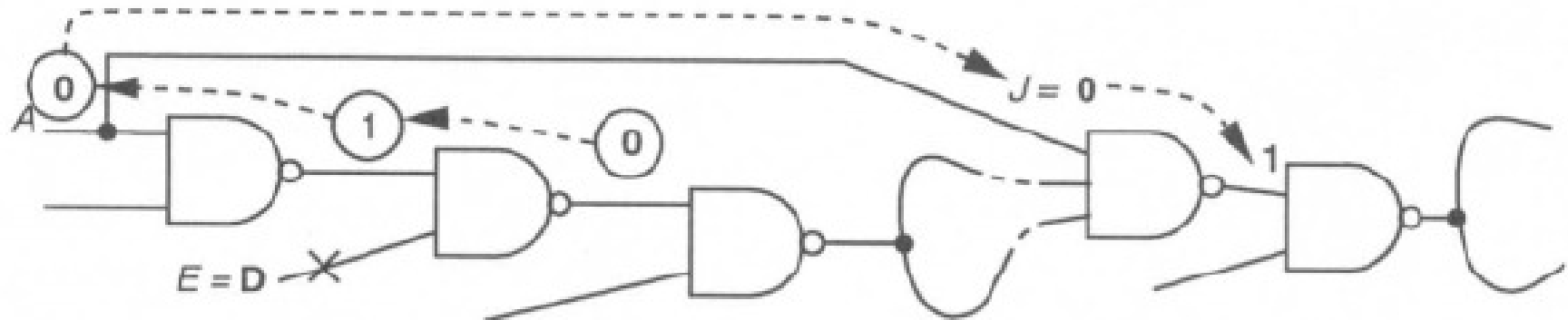
(a) Behavior of PODEM



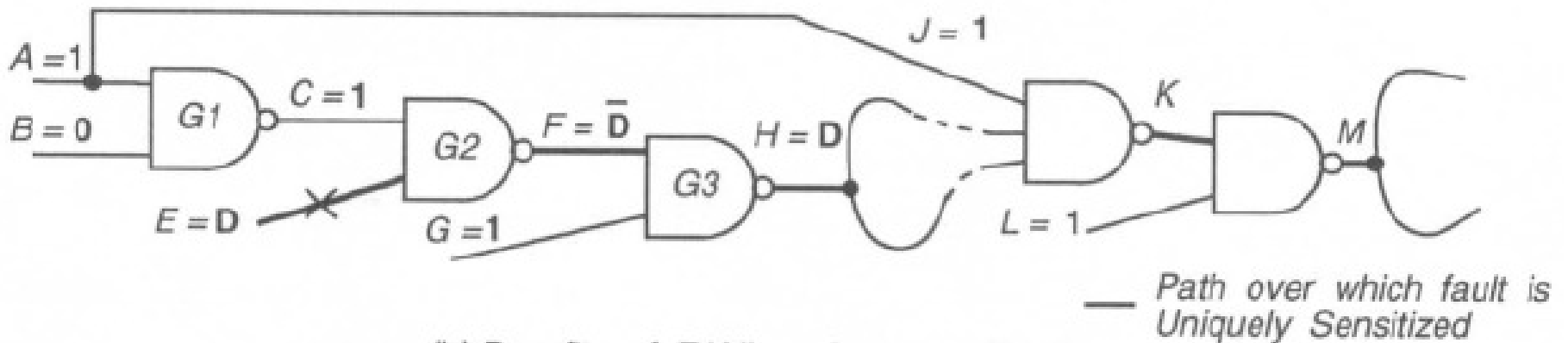
(b) Behavior of FAN.

Advanced comb. ATPG.....FAN

Unique sensitization



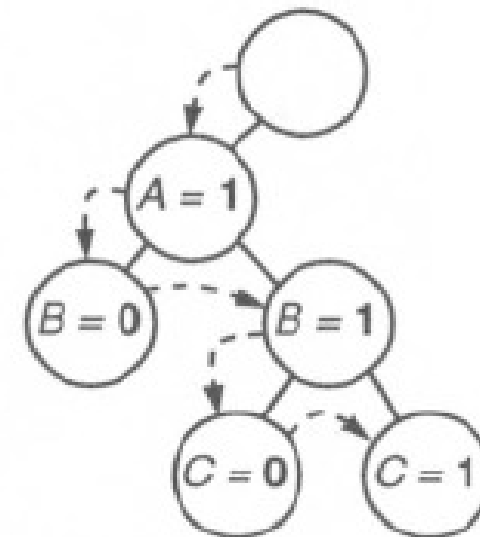
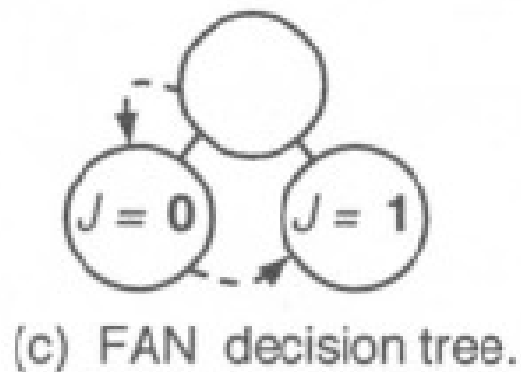
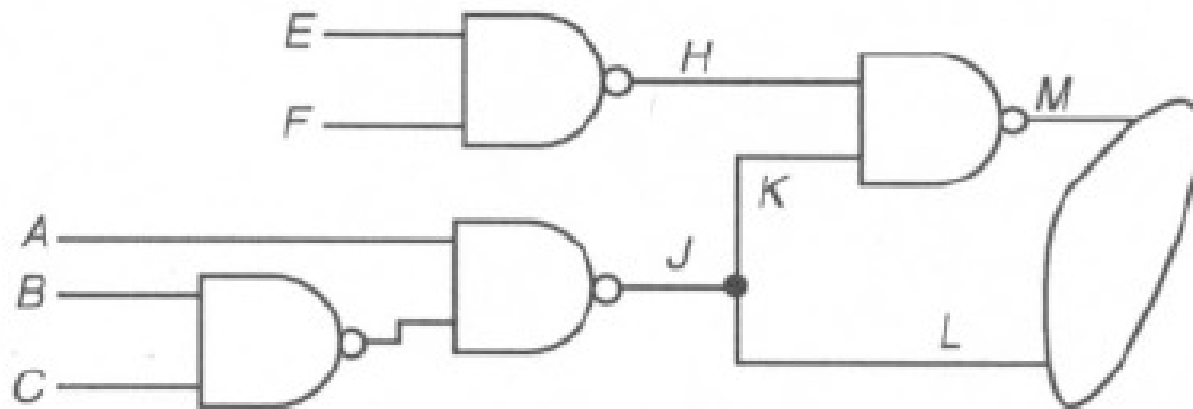
(a) Problems with PODEM.



(b) Benefits of FAN's unique sensitization.

Advanced comb. ATPG.....FAN

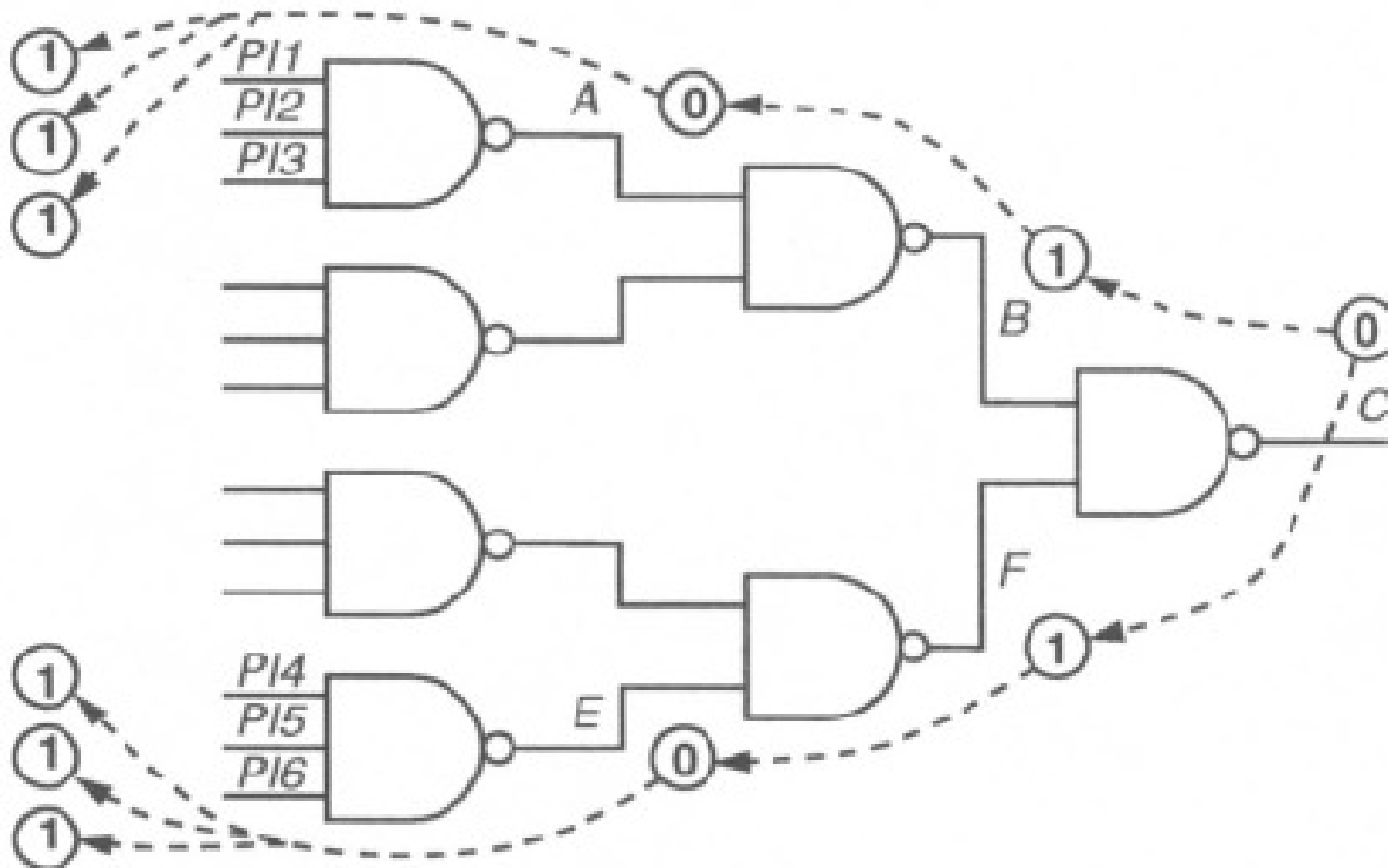
- Assign values to **headlines** (instead of PI)



(b) PODEM decision tree.

Advanced comb. ATPG.....FAN

Multiple Backtrace



Summary

- ❑ Structural vs. functional test
- ❑ Three types of ATPG algorithm
- ❑ Path sensitization algorithms are the preferred ATPG
 - D-Algorithm, PODEM, FAN
 - They use ATPG algebra
- ❑ Redundancy reduces the fault coverage
 - Masking effect
 - Reliability problems
- ❑ Testing is a global problem