VLSI Test Technology and Reliability (ET4076)

Lecture 5

Combinational Circuit Test Generation
(Chapter 7)

Said Hamdioui

Computer Engineering Lab
Delft University of Technology
2009-2010
Learning aims of today’s lecture

- Be able to
  - Define the different types of ATPG (concept, advantages, disadvantages)
  - Apply ATPG algorithms on combinational circuits
  - Explain the difference between the different path sanitization ATPG algorithms (D, PODEM, FAN)
Contents

- Algorithms and representations
  - Structural vs. functional test
  - Definition of ATPG
  - Search spaces
  - Algorithm completeness
  - ATPG algebra
  - Types of Algorithms

- Redundancy Identification

- Testing as a global problem

- ATPG algorithms
  - D-Calculus & D-Algorithm
  - PODEM Algorithm
  - FAN algorithm

- Advanced algorithms

- Test compaction
Functional vs. Structural ATPG

Adder with 129 inputs and 65 outputs

10 stuck-at faults

17 stuck-at faults
Functional vs. Structural (Continued)

- **Functional ATPG** – generate complete set of tests for circuit input-output combinations
  - 129 inputs, 65 outputs:
  - \(2^{129} = 680,564,733,841,876,926,926,749,214,863,536,422,912\)
  - Using 1 GHz ATE, would take \(2.15 \times 10^{22}\) years
  => Not practical

- **Structural test:**
  - No redundant adder hardware, 64 bit slices
  - Each with 27 faults (using fault equivalence)
  - At most \(64 \times 27 = 1728\) faults (at the most 1728 tests)
  - Takes \(0.000001728\) s on 1 GHz ATE

- Designer gives small set of functional tests (\(~ 70\%\) FC)
  - Augment with structural tests to boost coverage to 98\%
Definitions

- **ATPG: Automatic Test pattern generator**
  - Operations on digital hardware:
  - Inject fault into circuit modeled in computer
  - Use various ways to activate and propagate fault effect through hardware to circuit output
  - Output flips from expected to faulty signal

- **Electron-beam (E-beam) test**
  - Observes internal signals – “picture” of nodes charged to 0 and 1 in different colors
  - No need to propagate the fault effect to the PO
  - Too expensive

- **Scan design**
  - Add test hardware to all FFs to make them a shift Register
  - Can shift state in, scan state out (test mode)
  - Widely used: makes sequential test combinational
  - Costs: 5 to 20% chip area, circuit delay, extra pin, longer test sequence
Search space abstractions

Binary Search tree

- The corresponding binary tree
  - Gives all possible input patterns
  - Leaf: labeled with good output
  - ATPG implicitly search the tree to find test pattern
  - Worst case: complete examination
    - Exponential rise
    - # of leaves: $2^{#PI}$
Search space abstractions

Binary decision tree

- **Binary decision tree**
  - Follow path from root/source to sink node – product of literals along path gives Boolean value at sink
  - Rightmost path: $A \cdot B^* \cdot C^* = 1$
  - Leftmost path: $A^* \cdot B \cdot C^* = 0$
  - Problem: Size varies greatly with variable order
Algorithm Completeness

- An algorithm is **complete** if it ultimately can search entire binary decision tree, if needed, to generate a test
  - Completeness is important, otherwise ATPG may not achieve the required fault coverage

- **Untestable fault** – no test for it even after entire tree searched
  - Thus circuit behaves **correct** even in the presence of the fault

- Combinational circuits only – untestable faults are **redundant**, showing the presence of unnecessary hardware
# ATPG Algebra

- Roth’s 5-Valued and Muth’s 9-Valued
- ATPG algebra: higher order Boolean set of notation to presents both **good** and **failing** circuit **simultaneously**

<table>
<thead>
<tr>
<th>Symbol</th>
<th>Meaning</th>
<th>Good machine</th>
<th>Failing machine</th>
</tr>
</thead>
<tbody>
<tr>
<td>D</td>
<td>1/0</td>
<td>1</td>
<td>0</td>
</tr>
<tr>
<td>D*</td>
<td>0/1</td>
<td>0</td>
<td>1</td>
</tr>
<tr>
<td>0</td>
<td>0/0</td>
<td>0</td>
<td>0</td>
</tr>
<tr>
<td>1</td>
<td>1/1</td>
<td>1</td>
<td>1</td>
</tr>
<tr>
<td>X</td>
<td>X/X</td>
<td>X</td>
<td>X</td>
</tr>
<tr>
<td>G0</td>
<td>0/X</td>
<td>0</td>
<td>X</td>
</tr>
<tr>
<td>G1</td>
<td>1/X</td>
<td>1</td>
<td>X</td>
</tr>
<tr>
<td>F0</td>
<td>X/0</td>
<td>X</td>
<td>0</td>
</tr>
<tr>
<td>F1</td>
<td>X/1</td>
<td>X</td>
<td>1</td>
</tr>
</tbody>
</table>

Roth’s Algebra [5 values]

Muth’s additions [Extended unknowns]
Roth’s and Muth’s Higher-Order Algebras

- Represent two machines, which are simulated simultaneously by a computer program:
  - Good circuit machine (1st value)
  - Bad circuit machine (2nd value)

- Better to represent both in the algebra:
  - Need only 1 pass of ATPG to solve both

- Needed for complete ATPG:
  - Combinational: Roth Algebra, Multi-path sensitization,
  - Sequential: Muth Algebra -- good and bad machines may have different initial values due to fault
Type of algorithms

- **Exhaustive**
  - Expensive
  - Not practical, unless circuit partitioned

- **Random-Pattern Generation**
  - Easy to implement
  - Does not realize higher fault coverage

- **Path Sensitization Method**
  - Preferred ATPG method
Type of algorithms…: Exhaustive

- For $n$-input circuit, generate all $2^n$ input patterns

- Infeasible, unless circuit is **partitioned** into cones of logic, with $\leq 15$ inputs
  - Perform exhaustive ATPG for each cone
  - Misses faults that require specific activation patterns for multiple cones to be tested
Type of algorithms….Random-Pattern Generation

- Flow chart for method
- Use to get tests for 60-80% of faults, then switch to D-algorithm or other ATPG for rest
Type of algorithms...... Path Sensitization Method

- Is the preferred ATPG method
- Approach based on **three** steps
  1. **Fault sensitization**
     - Line SAF is activated by forcing it to an opposite value from the fault value
     - Also known as **Fault activation, fault excitation**
  2. **Fault propagation**
     - Fault effect is to be propagated through one or more paths to PO
     - Also known as **path sensitization**
  3. **Line justification**
     - Justify the sensitized fault by setting the PI of the circuits

- Steps 2 and 3 may find a **conflict**
  - ATPG has to **backup** or **backtrack**
  - i.e., discard previously made signal assignment and make an alternative assignment
Example:

- **Fault sensitization**: set B to 1
Example:

- **Fault sensitization**: set B to 1
- **Fault propagation**: Try path \( f \rightarrow h \rightarrow k \rightarrow L \)
- **Line justification**: There is no way to justify the 1 on i
  - **Conflict**
  - **Backtrack**
Type of algorithms…… Path Sensitization Method

- **Example:**
  - **Fault sensitization:** set B to 1
  - **Fault propagation:**
    - Try simultaneous paths $f - h - k - L$ and $g - i - j - k - L$
    - Blocked at $k$ because $D$-frontier ($D$ or $D^*$) disappears
Type of algorithms ….. Path Sensitization Method

Example:
- **Fault sensitization**: set B to 1
- **Fault propagation**: Try path \( g \rightarrow i \rightarrow j \rightarrow k \rightarrow L \)
- **Line justification**: set A to 0

=> Test found and fault detected

![Diagram](image)
ATPG Computational Complexity

- Ibarra and Sahni analysis [1975] – **NP-Complete**
  - No polynomial expression found for compute time, presumed to be exponential

- Worst case:
  - `no_pi` inputs, $2^{no\_pi}$ input combinations
  - `no_ff` flip-flops, $4^{no\_ff}$ initial flip-flop states
    - (good machine 0 or 1 x bad machine 0 or 1)
  - Work to forward simulate or reverse simulate all logic gates $n$ rises proportional with $n$
  - Complexity: $O(n \times 2^{no\_pi} \times 4^{no\_ff})$
ATPG Computational Complexity (Cnt)

- History ATPG
  - Improve heuristic algorithms & procedures to:
    - Find all necessary signal assignments as early as possible
    - Search as little as possible of the decision space
    - Worst case complexity: \( O(n \times 2^{\text{no}_{-}\text{pi}} \times 4^{\text{no}_{-}\text{ff}}) \)

- Fault simulation
  - For combinational circuits: \( O(n^2) \)
  - For sequential circuits: estimated between \( O(n^2) \) and \( O(n^3) \)
    - Based on empirical measurements

  ⇒ whenever possible use fault simulation to avoid ATPG computations

E.g.,
- Use RPG and fault simulation to get tests
- Use ATPG for hard-to-test faults
ATPG Computational Complexity* (Cnt)

ATPG for analog:
- Analog Fault Modeling Impractical for Logic ATPG

- Problems with modeling actual defects in analog circuits
  - Huge # of different possible analog faults in digital circuit
  - Exponential complexity of ATPG algorithm – a 20 flip-flop circuit can take days of computing
    - Cannot afford to go to a lower-level model

- Most test-pattern generators for digital circuits cannot even model at the transistor switch level
  - More complex
  - (see textbook for 5 examples of switch-level ATPG)
Redundancy identification... Untestable faults (1)

- Combinational ATPG can find redundant (unnecessary) hardware.
- Untestable faults in combinational circuits indicate redundant hardware.

Fault | Test
--- | ---
\(a \text{ sa1}, \ b \text{ sa0}\) | \(A = 1\)
\(a \text{ sa0}, \ b \text{ sa1}\) | \(A = 0\)

All faults are testable.
Therefore, these faults are not redundant.
Redundancy identification... Untestable faults(2)

Test fault: s-a-0 on line d

- Set A & B to 1
- However, E will be 1
- Fault cannot be propagated to the output
  => Untestable fault
  => Redundant Hardware

- Remove the redundant hardware
  - d always 0: ground it permanently
  - Remove OR gate and replace with a wire
  - And discard AND gate and input A
Redundancy identification .... Untestable faults(3)

- Test fault q s-a-1
  - Fault sensitization: set q to 0
  - Fault propagation: set A to 1 and s to 1
  - Line justification: set B to 0 and C to 0
  ⇒ Conflict at s
  No other possible alternative assignments
  ⇒ Redundant fault
Redundancy identification …. Fault masking (1)

- **f** s-a-0 tested when fault **q** s-a-1 **not** there
  - Fault sensitization: set **f** to 1
  - Fault propagation (via felqz): set **d** to 1, **m** to 0, **s** to 1, **n** to 1
  - Line justification: set **A** to 1, **B** to 1 and **C** to 0
  - ⇒ Fault detected at **Z**

![Diagram of a digital circuit with inputs A, B, C, and corresponding outputs and gates for sensitization and propagation of faults.]
Redundancy identification .... Fault masking (2)

- $f$ s-a-0 tested when fault $q$ s-a-1 also **present**
  - Fault sensitization: set $f$ to 1
  - Fault propagation (via $felqz$): set $d$ to 1, $m$ to 0, ...
    - **Blocked** by $q$ s-a-1  $\Rightarrow$ Fault can **NOT** be detected
  - $q$ s-a-1 (redundant) **masks** $f$ s-a0

![Diagram of a digital circuit](image)
Eliminates **hazards** in circuit output

- **OUT0** = A.B + A*.C (+B.C)
  - B.C is added to prevent hazards (when A 0 → 1)
- **Redundant fault s-a-0 on line e**
  - No effect on the function
  - But **masks** any testable fault using line e to propagate to OUT
  - Additional area and power due to redundancy
Redundancy impacts:

- Performance
- Power consumption
- Area overhead
- Reliability
  - A redundant fault can **mask** the presence of other testable faults
  - Unwanted in systems critical to human safety or wealth!!

How to deal with it?

- Impossible to do that manually (e.g., Millions of gates)
- Commercial tools
  - Synthesis moderately-sized designs to irredundant hardware
  - Big designs have to be partitioned before synthesis
    - May introduce redundant hardware
    - ATPG is one of the best ways to find this redundancy
Testing as a global problem

- Testing is a **global problem**
- Combination of **fully testable** modules in a logic circuit are
  - **Not** necessary fully testable
  - May be not be testable with the same patterns that would test the modules individually

Example

- Module 1
  - Requires \( AB = 11 \) for \( d \ s-a-0 \)
  - The test cannot be used for the entire circuit
ATPG algorithms

- Definitions
- D-Algorithm (Roth) -- 1966
  - D-cubes
  - Bridging faults
  - Logic gate function change faults
- PODEM (Goel) -- 1981
  - X-Path-Check
  - Backtracing
- Summary
### ATPG algorithms... History & Speedups

<table>
<thead>
<tr>
<th>Algorithm</th>
<th>Est. speedup over D-ALG</th>
<th>Year</th>
</tr>
</thead>
<tbody>
<tr>
<td>D-ALG</td>
<td>1</td>
<td>1966</td>
</tr>
<tr>
<td>PODEM</td>
<td>7</td>
<td>1981</td>
</tr>
<tr>
<td>FAN</td>
<td>23</td>
<td>1983</td>
</tr>
<tr>
<td>TOPS</td>
<td>292</td>
<td>1987</td>
</tr>
<tr>
<td>SOCRATES</td>
<td>1574</td>
<td>ATPG System 1988</td>
</tr>
<tr>
<td>Waicukauski et al.</td>
<td>2189</td>
<td>ATPG System 1990</td>
</tr>
<tr>
<td>EST</td>
<td>8765</td>
<td>ATPG System 1991</td>
</tr>
<tr>
<td>TRAN</td>
<td>3005</td>
<td>ATPG System 1993</td>
</tr>
<tr>
<td>Recursive learning</td>
<td>485</td>
<td>1995</td>
</tr>
<tr>
<td>Tafertshofer et al.</td>
<td>25057</td>
<td>1997</td>
</tr>
</tbody>
</table>
Fault Cone and D-frontier

- **Fault Cone**: Set of hardware affected by fault
- **D-frontier**: Set of all gates with D or D* at inputs and X at the output
ATPG algorithms........Definitions(2)

Forward Implication

- Results in logic gate inputs that are significantly labeled so that output is uniquely determined.

- Example:
  - AND gate forward implication table:

<table>
<thead>
<tr>
<th>a</th>
<th>0</th>
<th>1</th>
<th>X</th>
<th>D</th>
<th>D</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
</tr>
<tr>
<td>1</td>
<td>0</td>
<td>1</td>
<td>X</td>
<td>D</td>
<td>D</td>
</tr>
<tr>
<td>X</td>
<td>0</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>X</td>
</tr>
<tr>
<td>D</td>
<td>0</td>
<td>D</td>
<td>X</td>
<td>D</td>
<td>0</td>
</tr>
<tr>
<td>D</td>
<td>0</td>
<td>D</td>
<td>X</td>
<td>0</td>
<td>D</td>
</tr>
</tbody>
</table>

VLSI Test Technology and Reliability, 2009-2010

34
Backward Implication

- **Unique** determination of all gate inputs when the gate output and some of the inputs are given

Example

\[ a \rightarrow 1 \rightarrow b \rightarrow 1 \rightarrow 0 \rightarrow z \]

\[ a \rightarrow D \rightarrow b \rightarrow 0 \rightarrow D \rightarrow z \]

\[ a \rightarrow 0 \rightarrow b \rightarrow 1 \rightarrow 0 \rightarrow z \]
Implication Stack

- It is a Push-down stack; it records:
  - Each signal set in circuit by ATPG
  - Whether alternate signal value already tried
  - Portion of binary search tree already searched

<table>
<thead>
<tr>
<th>Signal</th>
<th>Value</th>
<th>Alternative tried</th>
</tr>
</thead>
<tbody>
<tr>
<td>A</td>
<td>1</td>
<td>NO</td>
</tr>
<tr>
<td>C</td>
<td>1</td>
<td>NO</td>
</tr>
<tr>
<td>E</td>
<td>1</td>
<td>NO</td>
</tr>
<tr>
<td>B</td>
<td>0</td>
<td>YES</td>
</tr>
</tbody>
</table>
A TP G algorithms........Definitions(5)

- Implication Stack after backtrack

<table>
<thead>
<tr>
<th>Signal</th>
<th>Value</th>
<th>Alternative tried</th>
</tr>
</thead>
<tbody>
<tr>
<td>E</td>
<td>1</td>
<td>NO</td>
</tr>
<tr>
<td>B</td>
<td>0</td>
<td>YES</td>
</tr>
<tr>
<td>F</td>
<td>0</td>
<td>YES</td>
</tr>
</tbody>
</table>

Stack ptr.

Unexplored
- Present Assignment
- Searched and Infeasible
ATPG algorithms........Definitions(6)

- **ATPG Objective**
  - Desired signal value to be achieved
  - Guides it away from infeasible/hard solutions

- **Backtrace**
  - Determines which primary input and value to set to achieve objective
  - Use testability measures
ATPG algorithms........Definitions(7)

Branch-and-Bound Search

- Efficiently searches binary search tree
  - **Branching** – At each tree level, selects which input variable to set to what value
  - **Bounding** – Avoids exploring large tree portions by artificially restricting search decision choices
    - Complete exploration is impractical
    - Uses *heuristics*
Fundamental concepts invented:
- First complete ATPG algorithm
- \textit{D-Cube}
- \textit{D-Calculus}
- \textit{Implications} – forward and backward
- \textit{Implication stack}
- \textit{Backtrack}
- Test Search Space
Singular Cover:

**Minimal** set of logic signal assignments to show essential *prime implicants* of Karnaugh map

![Diagram](image)

<table>
<thead>
<tr>
<th>Gate</th>
<th>Inputs</th>
<th>Output</th>
</tr>
</thead>
<tbody>
<tr>
<td>AND</td>
<td>A</td>
<td>B</td>
</tr>
<tr>
<td>1</td>
<td>0</td>
<td>X</td>
</tr>
<tr>
<td>2</td>
<td>X</td>
<td>0</td>
</tr>
<tr>
<td>3</td>
<td>1</td>
<td>1</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Gate</th>
<th>Inputs</th>
<th>Output</th>
</tr>
</thead>
<tbody>
<tr>
<td>NOR</td>
<td>d</td>
<td>e</td>
</tr>
<tr>
<td>1</td>
<td>1</td>
<td>X</td>
</tr>
<tr>
<td>2</td>
<td>X</td>
<td>1</td>
</tr>
<tr>
<td>3</td>
<td>0</td>
<td>0</td>
</tr>
</tbody>
</table>
D-Algorithm .............. Definitions(2)

D-Cube: Collapsed truth table entry to characterize logic block

- Use Roth’s 5-valued algebra
- Can change all D’s to D*’s and D*’s to D’s (do both)

Example: AND gate propagation:

<table>
<thead>
<tr>
<th>A</th>
<th>B</th>
<th>d</th>
</tr>
</thead>
<tbody>
<tr>
<td>D</td>
<td>1</td>
<td>D</td>
</tr>
<tr>
<td>1</td>
<td>D</td>
<td>D</td>
</tr>
<tr>
<td>D</td>
<td>D</td>
<td>D</td>
</tr>
<tr>
<td>D*</td>
<td>D*</td>
<td>D*</td>
</tr>
<tr>
<td>1</td>
<td>D*</td>
<td>D*</td>
</tr>
<tr>
<td>D*</td>
<td>1</td>
<td>D*</td>
</tr>
</tbody>
</table>

Combine rows 3 (good) & 1(faulty)
Reverse inputs
And two cubes
Interchange D and D*
D-Cube Operation of D-Intersection

- **D-intersection**: $0 \cap 0 = 0 \cap X = X \cap 0 = 0$
  $1 \cap 1 = 1 \cap X = X \cap 1 = 1$
  $X \cap X = X$

- $\psi$ – undefined (same as $\phi$)
- If both $\mu$ and $\lambda$ occur, then cubes incompatible
- If only $\mu$ occur, then $D \cap D = D$ and $D^* \cap D^* = D^*$
- If only $\lambda$ occur, then inverse $D$ and $D^*$ in second cube

- **D-containment:**
  Cube $a$ contains Cube $b$ if $b$ is a subset of $a$

<table>
<thead>
<tr>
<th>$\cap$</th>
<th>0</th>
<th>1</th>
<th>X</th>
<th>D</th>
<th>D*</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>0</td>
<td>$\phi$</td>
<td>0</td>
<td>$\psi$</td>
<td>$\psi$</td>
</tr>
<tr>
<td>1</td>
<td>$\phi$</td>
<td>1</td>
<td>1</td>
<td>$\psi$</td>
<td>$\psi$</td>
</tr>
<tr>
<td>X</td>
<td>0</td>
<td>1</td>
<td>X</td>
<td>D</td>
<td>D*</td>
</tr>
<tr>
<td>D</td>
<td>$\psi$</td>
<td>$\psi$</td>
<td>D</td>
<td>$\mu$</td>
<td>$\lambda$</td>
</tr>
<tr>
<td>D*</td>
<td>$\psi$</td>
<td>$\psi$</td>
<td>D*</td>
<td>$\lambda$</td>
<td>$\mu$</td>
</tr>
</tbody>
</table>
D-Algorithm............. Definitions(4)

Primitve D-Cube of Failure
- Models circuit faults:
  - Stuck-at-0
  - Stuck-at-1
  - Bridging fault (short circuit)
  - Arbitrary change in logic function

- AND Output sa0: “1 1 D”
- AND Output sa1: “0 X D*” and “X 0 D*”
- Wire sa0: “D”; Wire sa1: “D*”;

Test cube:
- Refers to the set of PI, intermediate, and PO circuit signals that are set to get a test for the fault

Propagation D-cube: models conditions under which fault effect propagates through gate
D-Algorithm…………… Definitions(5)

Implication Procedure

1. Model fault with appropriate *primitive D-cube of failure* (PDF)

2. Select *propagation D-cubes* to propagate fault effect to a circuit output (*D-drive* procedure)

3. Select *singular cover* cubes to justify internal circuit signals (*Consistency* procedure)
   - If Cube intersection fail, back up to the last decision point and select an alternative cube

- Put signal assignments in *test cube*
- Regrettably, cubes are selected very *arbitrarily* by D-ALG
D-Algorithm ……..Top Level

1. Number all circuit lines in increasing level order from PIs to POs;

2. Select a primitive D-cube of the fault to be the test cube;
   - Put logic outputs with inputs labeled as D (D*) onto the D-frontier;
     
     // [provoke the fault and D-cube]

3. D-drive ();         // [Propagate D or D* to PO]

4. Consistency ();     // [Tracking backwards to PI]

5. return ();
D-Algorithm ........... D-drive

while (untried fault effects on D-frontier)
    select next untried D-frontier gate for propagation;
while (untried fault effect fanouts exist)
    select next untried fault effect fanout;
generate next untried propagation D-cube;
D-intersect selected cube with test cube;
if (intersection fails or is undefined) continue;
if (all propagation D-cubes tried & failed) break;
if (intersection succeeded)
    add propagation D-cube to test cube -- recreate D-frontier;
    Find all forward & backward implications of assignment;
    save D-frontier, algorithm state, test cube, fanouts, fault;
    break;
else if (intersection fails & D and D* in test cube) Backtrack ();
else if (intersection fails) break;
if (all fault effects unpropagatable) Backtrack ();
D-Algorithm......Consistency

\( g \) = coordinates of test cube with 1’s & 0’s;
if \( g \) is only PIs fault testable & stop;
for (each unjustified signal in \( g \))
    Select highest # unjustified signal \( z \) in \( g \), not a PI;
    if (inputs to gate \( z \) are both \( D \) and \( D^* \)) break;
    while (untried singular covers of gate \( z \))
        select next untried singular cover;
        if (no more singular covers)
            If (no more stack choices) fault untestable & stop;
            else if (untried alternatives in Consistency)
                pop implication stack -- try alternate assignment;
            else
                \textbf{Backtrack} ();
                \textbf{D-drive} ();
                If (singular cover D-intersects with \( z \)) delete \( z \) from \( g \), add inputs to singular cover to \( g \), find all forward and backward implications of new assignment, and break;
                If (intersection fails) mark singular cover as failed;
D-Algorithm...... *Backtrack*

if (PO exists with fault effect) *Consistency ();* else pop prior implication stack setting to try alternate assignment; if (no untried choices in implication stack) fault untestable & stop; else return;
D-Algorithm…… Example1(1)

- Circuit Example 7.1 and Truth Table

<table>
<thead>
<tr>
<th>Inputs</th>
<th>Output</th>
</tr>
</thead>
<tbody>
<tr>
<td>a b c</td>
<td>F</td>
</tr>
<tr>
<td>0 0 0</td>
<td>0</td>
</tr>
<tr>
<td>0 0 1</td>
<td>0</td>
</tr>
<tr>
<td>0 1 0</td>
<td>0</td>
</tr>
<tr>
<td>0 1 1</td>
<td>1</td>
</tr>
<tr>
<td>1 0 0</td>
<td>0</td>
</tr>
<tr>
<td>1 0 1</td>
<td>0</td>
</tr>
<tr>
<td>1 1 0</td>
<td>0</td>
</tr>
<tr>
<td>1 1 1</td>
<td>0</td>
</tr>
</tbody>
</table>
D-Algorithm …… Example 1(2)

- **Singular cover** – Used for **justifying** lines

- **Propagation D-cubes** – Conditions under which difference between good/failing machines propagates

<table>
<thead>
<tr>
<th></th>
<th>A</th>
<th>B</th>
<th>C</th>
<th>d</th>
<th>e</th>
<th>F</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>1</td>
<td>0</td>
<td>1</td>
<td>0</td>
<td>0</td>
<td>0</td>
</tr>
<tr>
<td>0</td>
<td>0</td>
<td>1</td>
<td>0</td>
<td>1</td>
<td>1</td>
<td>0</td>
</tr>
<tr>
<td></td>
<td>1</td>
<td>0</td>
<td>0</td>
<td>1</td>
<td>1</td>
<td>0</td>
</tr>
<tr>
<td>D</td>
<td>1</td>
<td>1</td>
<td>0</td>
<td>0</td>
<td>1</td>
<td>0</td>
</tr>
<tr>
<td>D</td>
<td>D</td>
<td>D</td>
<td>D</td>
<td>D*</td>
<td>D*</td>
<td>D*</td>
</tr>
<tr>
<td>D</td>
<td>D</td>
<td>D</td>
<td>D</td>
<td>D*</td>
<td>D*</td>
<td>D*</td>
</tr>
<tr>
<td>D</td>
<td>D</td>
<td>D</td>
<td>D</td>
<td>D*</td>
<td>D*</td>
<td>D*</td>
</tr>
<tr>
<td>D</td>
<td>D</td>
<td>D</td>
<td>D</td>
<td>D*</td>
<td>D*</td>
<td>D*</td>
</tr>
</tbody>
</table>
D-Algorithm…….. Example 1(3)

- Steps for Fault \(d\) sa0

<table>
<thead>
<tr>
<th>Step</th>
<th>A</th>
<th>B</th>
<th>C</th>
<th>d</th>
<th>e</th>
<th>F</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>1</td>
<td>1</td>
<td></td>
<td>D</td>
<td></td>
<td></td>
</tr>
<tr>
<td>2</td>
<td></td>
<td></td>
<td></td>
<td>D</td>
<td>0</td>
<td>D*</td>
</tr>
<tr>
<td>3</td>
<td>1</td>
<td>1</td>
<td></td>
<td></td>
<td>0</td>
<td></td>
</tr>
</tbody>
</table>

Cube type
- PDF of AND gate
- Prop. D-cube for NOR
- Singular Cover of NAND
Example 7.2 Fault A sa0

**Step 1:** Select a primitive D-cube of the fault - Set \( A = 1 \)
D-Algorithm........ Example2(2)

- Step 2: **D-Drive** – Set $f = 0$
Step 3: \textit{D-Drive} – Set $k = 1$
Step 4: **Consistency** – Set $g = 1$
D-Algorithm........ Example2(5)

- **Step 5: Consistency –** $f = 0$ Already set
**Step 6: Consistency** – Set $c=0$, Set $e=0$
D-Algorithm……. Example2(7)

- **Step 7:** *Consistency* – Set $B = 0$
- *No more line to justify:* pattern found

**Test cube:** $A, B, C, D, e, f, g, h, k, L$

$TC^7 = D00X001D1D^*$
PODEM................Motivation

- IBM introduced semiconductor DRAM memory into its mainframes – late 1970’s

- Memory had error correction and translation circuits – improved reliability
  - D-ALG unable to test these circuits
    - Search too undirected
    - Large XOR-gate trees
    - Must set all external inputs to define output
  - Needed a better ATPG tool
PODEM………Goel IBM (1981)

- **New concepts introduced:**
  - Expand binary decision tree only around primary inputs
    - Reduction from $2^n$ to $2^{|PI|}$ ($n$: # of logic gates)
  - Use *X-PATH-CHECK* to test whether *D-frontier* still there
  - **Objectives**: bring ATPG closer to propagating $D$ ($D^*$) to PO
  - **Backtracing**
    - Use # of levels or CC measures to assign PI objectives
PODEM…… High-Level Flow

1. Assign binary value to unassigned PI
2. Determine implications of all PIs
3. Test Generated? If so, done.
4. Test possible with additional assigned PIs? If maybe, go to Step 1
5. Is there untried combination of values on assigned PIs? If not, exit: untestable fault
6. Set untried combination of values on assigned PIs using objectives and backtrace. Then, go to Step 2
PODEM ................. Example(1)

- Example 7.3
- Step 1: Select path $s - Y$ for fault propagation; not $s-u-v-Z$ (level distance from PO is 1 versus 2)

(a,b) means that the line has CC0 = a and CC1 = b
Step 2:

- Initial objective: Set $r$ to 1 to sensitize fault

(a,b) means that the line has CC0 = a and CC1 = b
Example (3)

Step 3: Backtrace from \( r \)

\[ \Rightarrow \text{Intermediate objectives: Set } n \text{ to } 0, \text{ set } m \text{ to } 0, \text{ set } g \text{ to } 0 \text{ and set } d \text{ to } 0 \]

\((a,b)\) means that the line has \( CC0 = a \) and \( CC1 = b \)
PODEM .............. Example(4)

- **Step 4:** Set $A = 0$ in implication stack

(a, b) means that the line has $CC0 = a$ and $CC1 = b
PODEM............. Example(5)

- **Step 5: Forward implications: \( d = 0, \ X = 1 \)**
  - Does not define \( r \) (s-a-1)

(a, b) means that the line has CC0 = a and CC1 = b
PODEM ............... Example(6)

- **Step 6: Initial objective: set \( r \) to 1**

\[(a,b)\] means that the line has CC0 = a and CC1 = b
PODEM

Example(7)

Step 7: Backtrace from \( r \) again

(a,b) means that the line has CC0 = a and CC1 = b
PODEM

Example(8)

- **Step 8**: Set $B$ to 1.
- **Implications in stack**: $A = 0$, $B = 1$

(a,b) means that the line has $CC0 = a$ and $CC1 = b$
Step 9: Forward implications: $k=1$, $m=0$, $r=1$, $q=1$, $Y=1$, $s=D^*$, $u=D^*$, $v^*=D$, $Z=1$
**PODEM............... Example (10)**

- **Step 10:** *X-PATH-CHECK* shows paths *s–Y* and *s–u–v–Z* blocked (*D-frontier* disappeared)

(a,b) means that the line has \( \text{CC0} = a \) and \( \text{CC1} = b \)
PODEM.................. Example(11)

- **Step11**: Set $B = 0$ (alternate assignment)

\[(a,b) \text{ means that the line has } CC0 = a \text{ and } CC1 = b\]
PODEM ................ Example (12)

- **Step 12:** Forward implications: $d=0$, $X=1$, $m=1$, $r=0$, $s=1$, $q=0$, $Y=1$, $v=0$, $Z=1$. Fault not sensitized.

(a,b) means that the line has $CC0 = a$ and $CC1 = b$
PODEM…………… Example(13)

- **Step 13**: Set $A = 1$ (alternate assignment)
- No implications possible -> Backtrace $r$

(a,b) means that the line has $CC0 = a$ and $CC1 = b$
PODEM…………… Example(14)

Step 14: Backtrace from $r$ again

(a,b) means that the line has $CC0 = a$ and $CC1 = b$
PODEM

Example(15)

- **Step 15:** Set $B = 0$. Implications in stack: $A = 1$, $B = 0$

(a,b) means that the line has $CC0 = a$ and $CC1 = b$
PODEM…………… Example(16)

Step 16: Forward implications: \( d=0, \ X=1, \ m=1, \ r = 0 \). Conflict: fault not sensitized.

--> Backtrack

\[(a,b)\] means that the line has CC0 = a and CC1 = b
PODEM.............. Example(17)

- **Step 17: Set** \( B = 1 \) (alternate assignment)

\[(a,b)\) means that the line has \( CC0 = a \) and \( CC1 = b \]
Step 18: Forward implications: \( d=1, m=1, r=1, q=0, s=D^*, v=D^*, X=0, Y=D^* \)

\[(a,b)\) means that the line has \(CC0 = a\) and \(CC1 = b\)
PODEM ............ Algorithm

while (no fault effect at POs)
    if (xpathcheck (D-frontier)
        \( \langle l, v_l \rangle = \text{Objective} (\text{fault}, v_{\text{fault}}) \);
        \( \langle pi, v_{pi} \rangle = \text{Backtrace} (l, v_l) \);
        \( \text{Imply} (pi, v_{pi}) \);
        if (PODEM (\text{fault}, v_{\text{fault}}) == SUCCESS) return (SUCCESS);
        \( \langle pi, v_{pi} \rangle = \text{Backtrack} () \);
        \( \text{Imply} (pi, v_{pi}) \);
        if (PODEM (\text{fault}, v_{\text{fault}}) == SUCCESS) return (SUCCESS);
        \( \text{Imply} (pi, \text{"X"}) \);
        return (FAILURE);
    else if (implication stack exhausted)
        return (FAILURE);
    else Backtrack ()
        return (SUCCESS);
Summary classical ATPG algorithms

- D-ALG – First complete ATPG algorithm
  - \textit{D-Cube}
  - \textit{D-Calculus}
  - \textit{Implications} – forward and backward
  - Implication stack
  - Backup

- PODEM
  - Expand decision tree only around PIs
  - Use \textit{X-PATH-CHECK} to see if \textit{D-frontier} exists
  - Objectives -- bring ATPG closer to getting D (D*) to PO
  - Backtracing
Advanced combinational ATPG

- FAN – *Multiple Backtrace* (1983)
- TOPS – *Dominator* (1987)
- Legal Assignments (1990)
- EST – *Search space learning* (1991)
- BDD Test generation (1991)
Advanced comb. ATPG........FAN

- FAN(1983)
  - Further limit the ATPG space and accelerate Backtracing
  - Use **immediate implications** of uniquely-determined signals
  - **Unique sensitization**
  - Assign values to **headlines** (instead of PI)
  - **Multiple Backtrace**
Use **immediate implications** of uniquely-determined signals.
Advanced comb. ATPG........FAN

- Unique sensitization

(a) Problems with PODEM.

(b) Benefits of FAN's unique sensitization.

Path over which fault is Uniquely Sensitized
Assign values to **headlines** (instead of PI)
Advanced comb. ATPG …… FAN

- Multiple Backtrace
Summary

- Structural vs. functional test
- Three types of ATPG algorithm
- Path sensitization algorithms are the preferred ATPG
  - D-Algorithm, PODEM, FAN
  - They use ATPG algebra
- Redundancy reduces the fault coverage
  - Masking effect
  - Reliability problems
- Testing is a global problem